

# A gentle introduction to the: Succinct Data Structure Library

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# Lab lectures

- **Not** part of the syllabus but useful for your future career
- Tools for working with big data: Compressed/succinct data structures
  - Complex algorithms become possible in memory, no clusters and no disks
  - The libraries we will see offer the same API of standard data structures

# From theory to practice

- Integer coding via Elias- $\gamma$  and - $\delta$  codes (§11.1 of the notes)
- Plain bitvectors with rank/select support (§15.1.1)
- Compressed bitvectors via Elias-Fano coding (§11.6 and §15.1.2)
- Pointerless programming (§15.1)

# The Succinct Data Structure Library (SDSL)

<https://github.com/simongog/sdsl-lite>

(see also the fork <https://github.com/xxsds/sdsl-lite>)

- An **easy-to-use, highly-efficient, configurable, and extensible** library of succinct data structures for researchers and practitioners
- Implements highlights of 40 research articles
- Faithful to the original proposals while using modern instruction sets
- Written in C++11, API familiar to C++ STL users
- Wrappers in other languages (e.g. Python) available online

# SDSL resources

Cheatsheet: <http://simongog.github.io/assets/data/sdsl-cheatsheet.pdf>

API Docs: <http://algo2.iti.kit.edu/gog/docs/html/index.html>

Examples: <https://github.com/simongog/sdsl-lite/examples>

# sds1 Cheat Sheet

## Data structures

The library code is in the `sds1` namespace. Either import the namespace in your program (using `namespace sds1;`) or qualify all identifiers by a `sds1::`-prefix.

Each section corresponds to a header file. The file is hyperlinked as part of the section heading.

We have two types of data structures in `sds1`. *Self-contained* and *support* structures. A support object `s` can extend a self-contained object `o` (e.g. add functionality), but requires access to `o`. Support structures contain the substring `support` in their class names.

## Integer Vectors (IV)

The core of the library is the class `int_vector<w>`. Parameter `w` corresponds to the fixed length of each element in bits. For `w = 8, 16, 32, 64, 1` the length is fixed during compile time and the vectors correspond to `std::vector<uintw_t>` resp. `std::vector<bool>`. If `w = 0` (default) the length can be set during runtime. *Constructor:* `int_vector<>(n, x, ℓ)`, with `n` equals size, `x` default integer value, `ℓ` width of integer (has no effect for `w > 0`).

*Public methods:* `operator[i]`, `size()`, `width()`, `data()`.

## Manipulating int\_vector<w> v

Method	Description
<code>v[i]=x</code>	Set entry <code>v[i]</code> to <code>x</code> .
<code>v.width(ℓ)</code>	Set width to <code>ℓ</code> , if <code>w = 0</code> .
<code>v.resize(n)</code>	Resize <code>v</code> to <code>n</code> elements.
Useful methods in namespace <code>sds1::util</code> :	
<code>set_to_value(v, k)</code>	Set <code>v[i]=k</code> for each <code>i</code> .
<code>set_to_id(v)</code>	Set <code>v[i]=i</code> for each <code>i</code> .
<code>set_random_bits(v)</code>	Set elements to random bits.
<code>mod(v, m)</code>	Set <code>v[i]=v[i] mod m</code> for each <code>i</code> .
<code>bit_compress(v)</code>	Gets <code>x=max_v[i]</code> and <code>ℓ=ceil(log(x-1))+1</code> and packs the entries in <code>ℓ</code> -bit integers.
<code>expand_width(v, ℓ)</code>	Expands the width of each integer to <code>ℓ</code> bits, if <code>ℓ ≥ v.width()</code> .

## Compressed Integer Vectors (CIV)

For a vector `v`, `enc_vector` stores the self-delimiting coded deltas (`v[i+1]-v[i]`). Fast random access is achieved by sampling values of `v` at rate `t_dens`. Available coder are `coder::elias_delta`, `coder::elias_gamma`, and `coder::fibonacci`.

Class `vlc_vector` stores each `v[i]` as self-delimiting codeword. Samples at rate `t_dens` are inserted for fast random access.

Class `dac_vector` stores for each value `x` the least ( $t_b - 1$ ) significant bits plus a bit which is set if  $x \geq 2^{b-1}$ . In the latter case, the process is repeated with  $x' = x/2^{b-1}$ .

## Bitvectors (BV)

Representations for a bitvector of length  $n$  with  $m$  set bits.

Class	Description	Space
<code>bit_vector</code>	plain bitvector	$64\lceil n/64+1\rceil$
<code>bit_vector_il</code>	interleaved bitvector	$\approx n(1+64/K)$
<code>rrr_vector</code>	$H_0$ -compressed bitvector	$\approx \log\binom{n}{m}$
<code>sd_vector</code>	sparse bitvector	$\approx m \cdot (2+\log\frac{n}{m})$
<code>hyb_vector</code>	hybrid bitvector	

`bit_vector` equals `int_vector<1>` and is therefore dynamic.

*Public Methods:* `operator[i]`, `size()`, `begin()`, `end()`

*Public Types:* `rank_1_type`, `select_1_type`, `select_0_type`<sup>1</sup>.

Each bitvector can be constructed out of a `bit_vector` object.

## Rank Supports (RS)

RSs add rank functionality to BV. Methods `rank(i)` and `operator(i)` return the number of set bits<sup>2</sup> in the prefix  $[0..i]$  of the supported BV for  $i \in [0, n]$ .

Class	Compatible BV	+Bits	Time
<code>rank_support_v</code>	<code>bit_vector</code>	$0.25n$	$\mathcal{O}(1)$
<code>rank_support_v5</code>	<code>bit_vector</code>	$0.0625n$	$\mathcal{O}(1)$
<code>rank_support_scan</code>	<code>bit_vector</code>	64	$\mathcal{O}(n)$
<code>rank_support_il</code>	<code>bit_vector_il</code>	128	$\mathcal{O}(1)$
<code>rank_support_rrr</code>	<code>rrr_vector</code>	80	$\mathcal{O}(k)$
<code>rank_support_sd</code>	<code>sd_vector</code>	64	$\mathcal{O}(\log\frac{n}{m})$
<code>rank_support_hyb</code>	<code>hyb_vector</code>	64	-

Call `util::init_support(rs, bv)` to initialize rank structure `rs` to bitvector `bv`. Call `rs(i)` to get `rank(i) = \sum_{k=0}^i bv[k]`

## Select Supports (SLS)

SLSs add select functionality to BV. Let  $m$  be the number of set bits in BV. Methods `select(i)` and `operator(i)` return the position of the  $i$ -th set bit<sup>3</sup> in BV for  $i \in [1..m]$ .

Class	Compatible BV	+Bits	Time
<code>select_support_mcl</code>	<code>bit_vector</code>	$\leq 0.2n$	$\mathcal{O}(1)$
<code>select_support_scan</code>	<code>bit_vector</code>	64	$\mathcal{O}(n)$
<code>select_support_il</code>	<code>bit_vector_il</code>	64	$\mathcal{O}(\log n)$
<code>select_support_rrr</code>	<code>rrr_vector</code>	64	$\mathcal{O}(\log n)$
<code>select_support_sd</code>	<code>sd_vector</code>	64	$\mathcal{O}(1)$

Call `util::init_support(sls, bv)` to initialize `sls` to bitvector `bv`. Call `sls(i)` to get `select(i) = min{j | rank(j+1) = i}`.

## Wavelet Trees (WT=BV+RS+SLS)

Wavelet trees represent sequences over byte or integer alphabets of size  $\sigma$  and consist of a tree of BVs. Rank and select on the sequences is reduced to rank and select on BVs, and the runtime is multiplied by a factor in  $[H_0, \log \sigma]$ .

Class	Shape	lex_ordered	Default	Traversable
<code>wt_rlmn</code>	underlying WT dependent			✗
<code>wt_gmr</code>	none	✗	integer	✗
<code>wt_ap</code>	none	✗	integer	✗
<code>wt_huff</code>	Huffman	✗	byte	✓
<code>wm_int</code>	Balanced	✗	integer	✓
<code>wt_blcd</code>	Balanced	✓	byte	✓
<code>wt_hutu</code>	Hu-Tucker	✓	byte	✓
<code>wt_int</code>	Balanced	✓	integer	✓

*Public types:* `value_type`, `size_type`, and `node_type` (if WT is

traversable). In the following let  $c$  be a symbol,  $i,j,k$ , and  $q$  integers,  $v$  a node, and  $r$  a range.

*Public methods:* `size()`, `operator[i]`, `rank(i, c)`, `select(i, c)`, `inverse_select(i)`, `begin(r)`, `end(r)`.

Traversable WTs provide also: `root()`, `is_leaf(v)`, `empty(v)`, `size(v)`, `sym(v)`, `expand(v)`, `expand(v, r)`, `expand(v, std::vector<r>)`, `bit_vec(v)`, `seq(v)`.

lex\_ordered WTs provide also: `lex_count(i, j, c)` and `lex_smaller_count(i, c)`. `wt_int` provides: `range_search_2d`.

`wt_algorithm.hpp` contains the following generic WT method (let `wt` be a WT object): `intersect(wt, vector<r>)`, `quantile_freq(wt, i, j, q)`, `interval_symbols(wt, i, j, k, ...)`, `symbol_lte(wt, c)`, `symbol_gte(wt, c)`, `restricted_unique_range_values(wt, xi, xj, yi, yj)`.

## Suffix Arrays (CSA=IV+WT)

Compressed suffix arrays use CIVs or WTs to represent the suffix arrays (SA), its inverse (ISA), BWT,  $\Psi$ , and LF. CSAs can be built over byte and integer alphabets.

Class	Description
<code>cса_bitcompressed</code>	Based on SA and ISA stored in a IV.
<code>cса_sada</code>	Based on $\Psi$ stored in a CIV.
<code>cса_wt</code>	Based on the BWT stored in a WT.

*Public methods:* `operator[i]`, `size()`, `begin()`, `end()`.

*Public members:* `isa`, `bwt`, `lf`, `psi`, `text`, `L`, `F`, `C`, `char2comp`, `comp2char`, `sigma`.

*Policy classes:* alphabet strategy (e.g. `byte_alphabet`, `succinct_byte_alphabet`, `int_alphabet`) and SA sampling strategy (e.g. `sa_order_sa_sampling`, `text_order_sa_sampling`)

## Longest Common Prefix (LCP) Arrays

Class	Description
<code>lcp_bitcompressed</code>	Values in a <code>int_vector</code> .
<code>lcp_dac</code>	Direct accessible codes used.
<code>lcp_byte</code>	Small values in a byte; 2 words per large.
<code>lcp_wt</code>	Small values in a WT; 1 word per large.
<code>lcp_vlc</code>	Values in a <code>vlc_vector</code> .
<code>lcp_support_sada</code>	Values stored permuted. CSA needed.
<code>lcp_support_tree</code>	Only depths of CST inner nodes stored.
<code>lcp_support_tree2</code>	+ large values are sampled using LF.

*Public methods:* `operator[i]`, `size()`, `begin()`, `end()`

We represent a sequence of parentheses as a `bit_vector`. An opening/closing parenthesis corresponds to 1/0.

*Class* *Description*

- `bp_support_g` Two-level pioneer structure.
- `bp_support_gg` Multi-level pioneer structure.
- `bp_support_sada` Min-max-tree over excess sequence.

*Public methods:* `find_open(i)`, `find_close(i)`, `enclose(i)`, `double_enclose(i, j)`, `excess(i)`, `rr_enclose(i, j)`, `rank(i)4`, `select(i)`.

Call `util::init_support(bps, bv)` to initialize a BPS `bps` to `bit_vector bv`.

## Suffix Trees (CST=CSA+LCP+BPS)

A CST can be parametrized by any combination of CSA , LCP, and BPS. The operation of each part can still be accessed through member variables. The additional operations are listed below. CSTs can be built for byte or integer alphabets.

Class	Description
<code>cst_sada</code>	Represents a node as position in BPS. Navigational operations are fast (they are directly translated in BPS operations on the DFS-BPS). Space: $4n+o(n)+ CSA + LCP $ bits.
<code>cst_sct3</code>	Represents nodes as intervals. Fast construction, but slower navigational operations. Space: $3n+o(n)+ CSA + LCP $

*Public types:* `node_type`. In the following let `v` and `w` be nodes and `i, d, lb, rb` integers.

Class	Description
<code>begin_bottom_up()</code>	
<code>end_bottom_up()</code>	
<code>size(v)</code>	
<code>is_leaf(v)</code>	
<code>degree(v)</code>	
<code>depth(v)</code>	
<code>node_depth(v)</code>	
<code>edge(v, d)</code>	
<code>lb(v)</code>	
<code>rb(v)</code>	
<code>id(v)</code>	
<code>inv_id(i)</code>	
<code>sn(v)</code>	
<code>select_leaf(i)</code>	
<code>node(lb, rb)</code>	
<code>parent(v)</code>	
<code>sibling(v)</code>	
<code>lca(v, w)</code>	
<code>select_child(v, i)</code>	
<code>child(v, c)</code>	
<code>children(v)</code>	
<code>sl(v)</code>	
<code>wl(v, c)</code>	
<code>leftmost_leaf(v)</code>	
<code>rightmost_leaf(v)</code>	

The `traversal` example shows how to use the DFS-iterator.

## Range Min/Max Query (RMQ)

A RMQ `rmq` can be used to determine the position of the minimum value<sup>5</sup> in an arbitrary subrange  $[i, j]$  of an preprocessed vector `v`. Operator `operator(i, j)` returns  $x = \min[r | r \in [i, j] \wedge v[r] \leq v[k] \forall k \in [i, j]]$

Class	Space	Time
<code>rmq_support_sparse_table</code>	$n \log^2 n$	$\mathcal{O}(1)$
<code>rmq_succint_sada</code>	$4n + o(n)$	$\mathcal{O}(1)$
<code>rmq_succint_sct</code>	$2n + o(n)$	$\mathcal{O}(1)$

## Constructing data structures

Let `o` be a WT-, CSA-, or CST-object. Object `o` is built with `construct(o, file, num_bytes=0)` from a sequence stored in `file`. File is interpreted dependent on the value of `num_bytes`: `Value` `File interpreted as`

<code>num_bytes=0</code>	serialized <code>int_vector</code> .
<code>num_bytes=1</code>	byte sequence of length <code>util::file_size(file)</code> .
<code>num_bytes=2</code>	16-bit word sequence.
<code>num_bytes=4</code>	32-bit word sequence.
<code>num_bytes=8</code>	64-bit word sequence.
<code>num_bytes=d</code>	Parse decimal numbers.

Note: `construct` writes/reads data to/from disk during construction. Accessing disk for small instances is a considerable overhead. `construct_im(o, data, num_bytes=0)` will build `o` using only main memory. Have a look at [this handy tool for an example](#).

## Configuring construction

The locations and names of the intermediate files can be configured by a `cache_config` object. It is constructed by `cache_config(del, tmp_dir, id, map)` where `del` is a boolean variable which specifies if the intermediate files should be deleted after construction, `tmp_dir` is a path to the directory

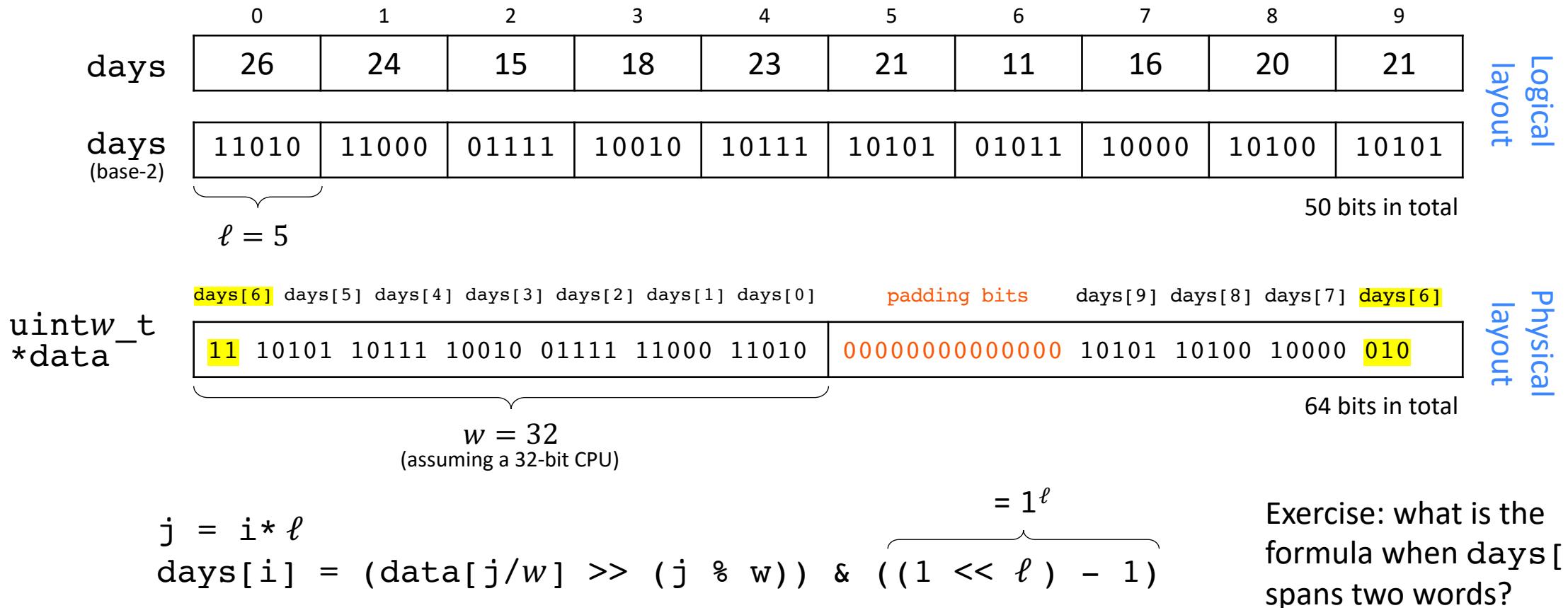
# The core of the library: `int_vector`

- Say you need to store the number of days per month worked by each employee of a company
- The `<cstdint>` header defines the (portable) types `intn_t` and `uintn_t` for  $n = 8, 16, 32, 64$
- So you would normally declare something like

```
std::vector<uint8_t> days(n_employees);
```
- What is the space? 1 byte per employee
- But only 5 bits are needed ( $31 = 11111_2$ )
- We are wasting  $8/5 - 1 = 60\%$  of extra space per employee!

# The core of the library: `int_vector`

- `int_vector<ℓ>` is a container of integers of fixed bit-width  $\ell$



# Solution to the exercise

```
j = i * ℓ

offset = j % w

if (offset + ℓ <= w) // the integer spans one word
    return (data[j/w] >> (j % w)) & ((1 << ℓ) - 1)
else                      // the integer spans two words
    return (data[j/w] >> (j % w)) |
        (data[j/w+1] & ((1 << ((offset+ℓ) % w)) - 1)) << (w-offset)
```

## Example 1: `int_vector<ℓ>`, bit-width $\ell$ set at compile-time

```
#include <iostream>
#include <sds1/vectors.hpp>

int main() {
    sds1::int_vector<5> v = {26, 24, 15, 18, 23, 21, 11, 16, 20, 21};
    std::cout << v << std::endl;
    v[3] = 10;
    v[1] = 194; // == 0b11000010
    std::cout << v << std::endl;
    std::cout << v.bit_size() << std::endl;
    std::cout << sds1::size_in_bytes(v) << std::endl;
    return 0;
}
```

26 24 15 18 23 21 11 16 20 21

26 2 15 10 23 21 11 16 20 21

50

16 ← 8 byte for the compressed v, 8 bytes for v.size()

## Example 2: `int_vector<>`, bit-width set at run-time

```
#include <iostream>
#include <sds1/vectors.hpp>

int main() {
    sds1::int_vector<> v = {26, 24, 15, 18, 23, 21, 11, 16, 20, 21};
    std::cout << sds1::size_in_bytes(v) << std::endl;
    sds1::util::bit_compress(v);
    std::cout << v << std::endl;
    std::cout << (int) v.width() << std::endl;
    std::cout << sds1::size_in_bytes(v) << std::endl;
    return 0;
}
```

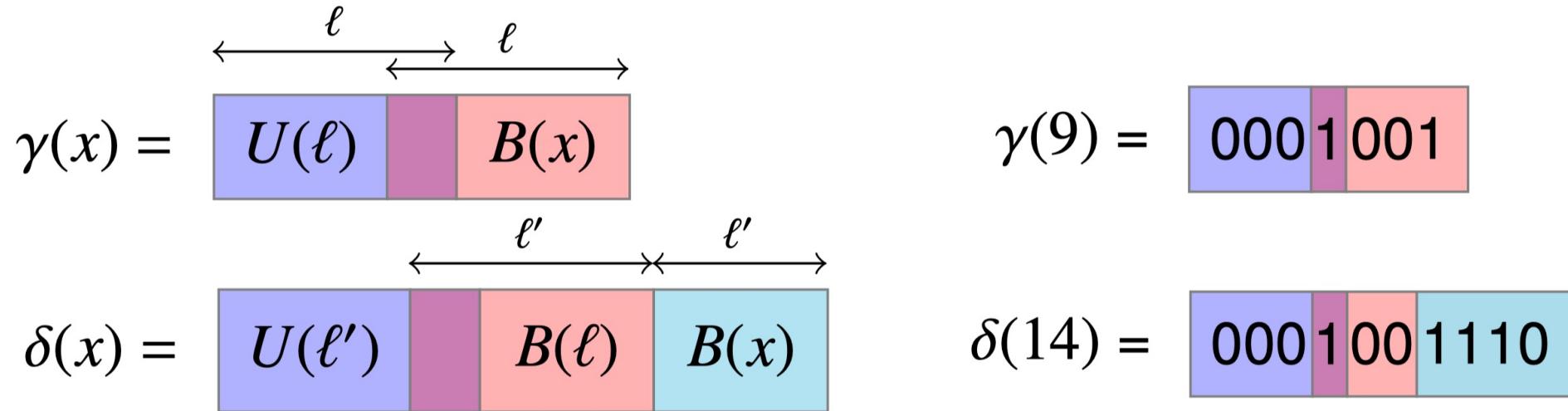
89 ← 80 bytes for the 10 ints, 8 bytes for `v.size()`, 1 byte for the bit-width

26 24 15 18 23 21 11 16 20 21

5

17 ← 8 byte for the compressed `v`, 8 bytes for `v.size()`, 1 byte for the bit-width

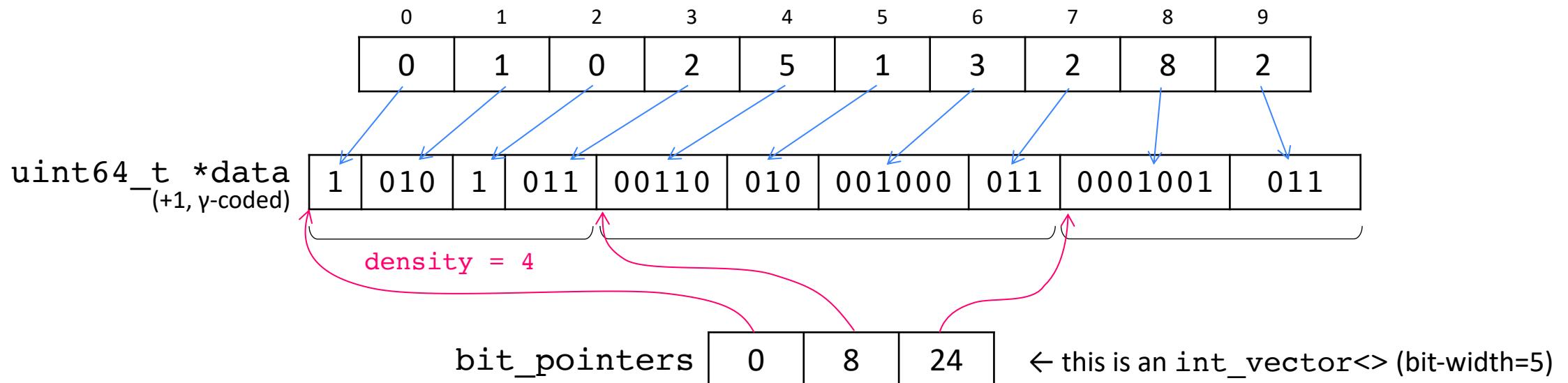
# Integer coders



- `sds1::coder::elias_gamma`
- `sds1::coder::elias_delta`

# Compressed integer vectors: `vlc_vector`

- Stores Elias-γ and -δ codes contiguously
- Zeros are permitted (internally, it uses  $x + 1$  instead of  $x$ )
- How to implement `vlc_vector::operator[ ]`?



**density** is a *space-time trade-off parameter*: decrease it for faster random access but higher space usage

# Example 3: vlc\_vector<coder, density>

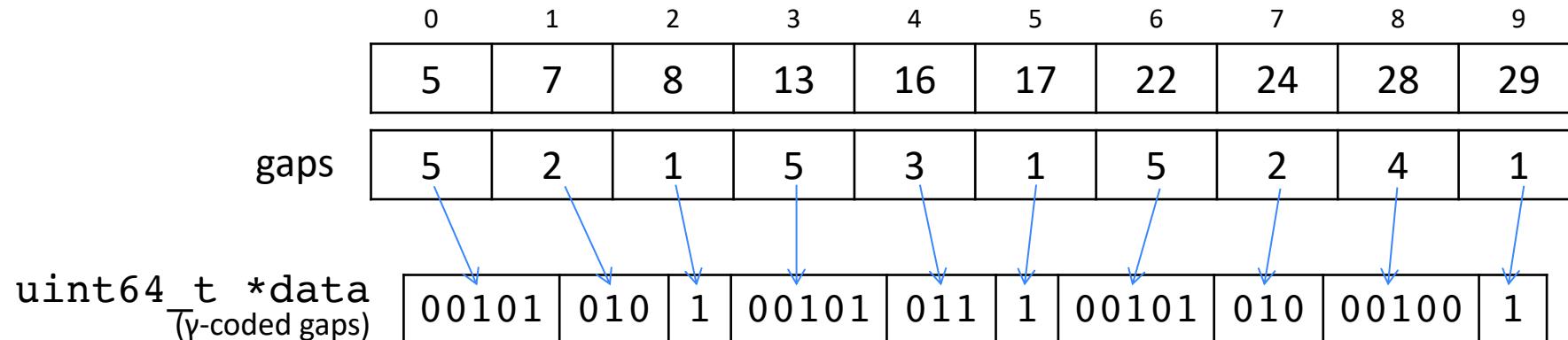
```
#include <iostream>
#include <sds1/vectors.hpp>

int main() {
    sds1::int_vector<> v(10 * (1 << 20));
    v[100] = 1ULL << 63;
    sds1::util::bit_compress(v);
    std::cout << size_in_mega_bytes(v) << std::endl;
    sds1::vlc_vector<sds1::coder::elias_delta, 128> vlc(v);
    std::cout << size_in_mega_bytes(vlc) << std::endl;
    return 0;
}
```

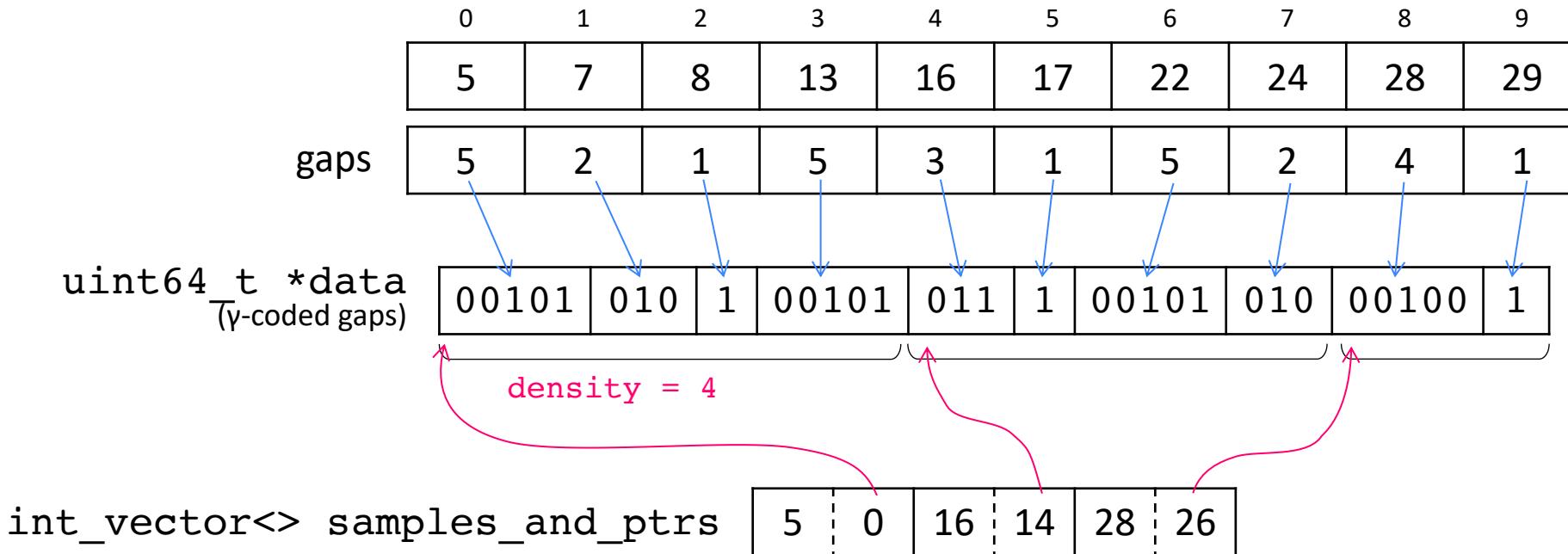
Output  
80  
1.48442

# Compressed integer vectors: `enc_vector`

- Say you have increasing integers  $x_1, x_2, x_3 \dots, x_n$   
A postings list (search engines), an adjacency list (graphs)
- $\gamma/\delta$ -code the gaps  $x_1, x_2 - x_1, x_3 - x_2, \dots, x_n - x_{n-1}$
- How to implement `vlc_vector::operator[ ]?`

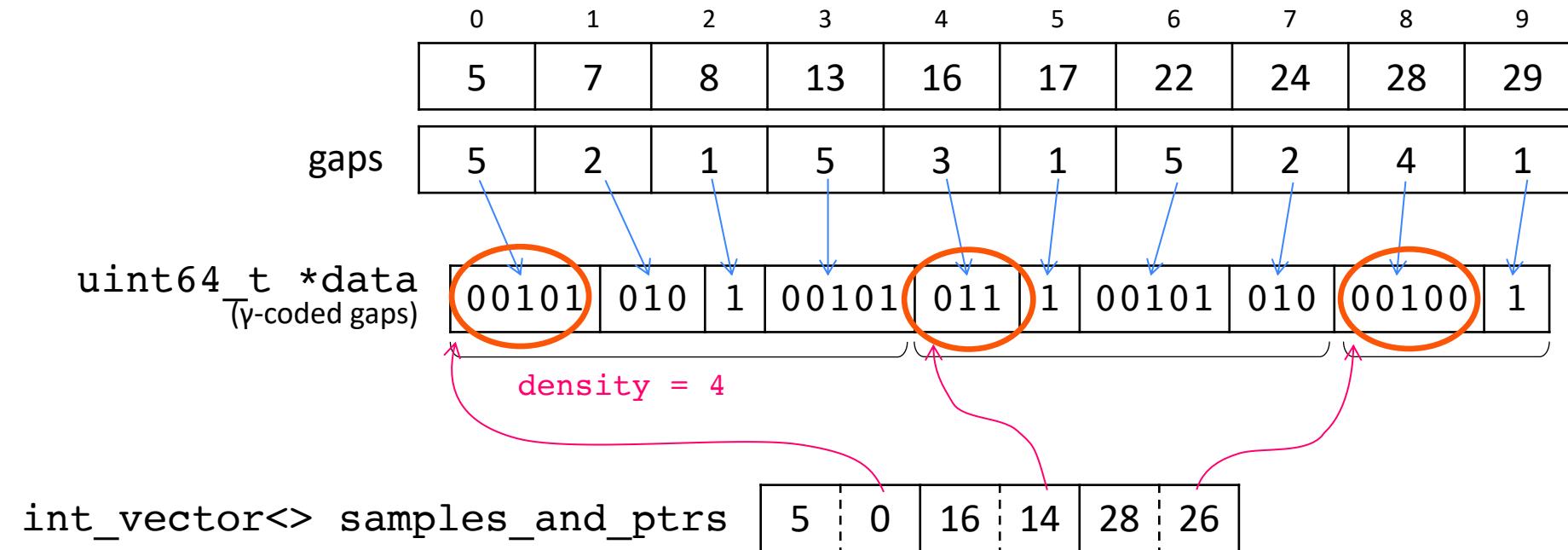


# Compressed integer vectors: `enc_vector`

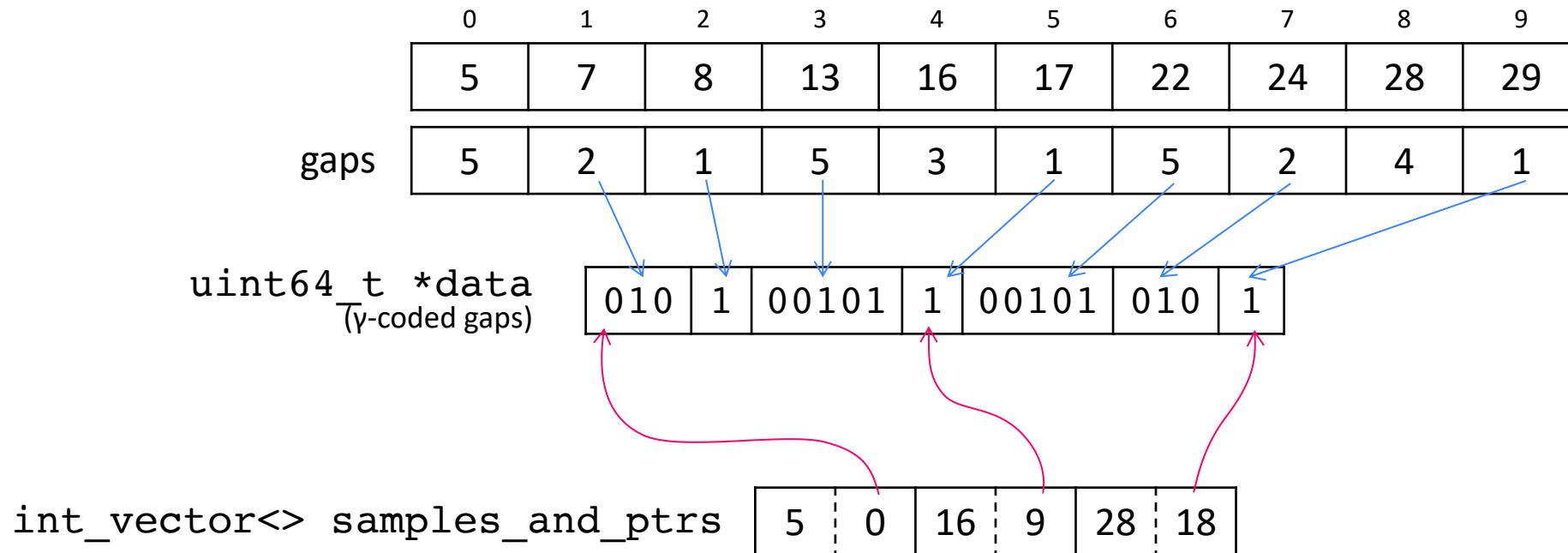


# Compressed integer vectors: `enc_vector`

- There is some **redundancy** here...



# Compressed integer vectors: `enc_vector`



**density** is a *space-time trade-off parameter*: decrease it for faster random access but higher space usage

# Example 4: `enc_vector<coder, density>`

```
#include <random>
#include <iostream>
#include <sds1/vectors.hpp>

sds1::int_vector<> random_data(size_t n, uint64_t u) { // Algorithm 3.4 of the notes
    if (n > u)
        throw std::invalid_argument("n > u");
    std::mt19937 rnd(std::random_device{}());
    sds1::int_vector<> out(n);
    for (size_t j = 1, s = 0; j <= u && s < n; ++j)
        if (rnd() % (u - j + 1) < n - s)
            out[s++] = j;
    return out;
}
...
```

# Example 4: `enc_vector<coder, density>`

```
#include <random>
#include <iostream>
#include <sdsl/vectors.hpp>

sdsl::int_vector<> random_data(size_t n, uint64_t u) { ... }

int main() {
    auto data = random_data(10 * (1 << 20), 1 << 24);
    std::cout << sdsl::size_in_mega_bytes(data) << std::endl;

    sdsl::vlc_vector<sdsl::coder::elias_delta, 128> vlc(data);
    std::cout << sdsl::size_in_mega_bytes(vlc) << std::endl;

    sdsl::enc_vector<sdsl::coder::elias_delta, 128> enc_delta(data);
    std::cout << sdsl::size_in_mega_bytes(enc_delta) << std::endl;

    return 0;
}
```

What if we use `sdsl::coder::elias_gamma`?

80  
39.0282  
3.19339

# Example 4(bis): `enc_vector<coder, density>`

```
#include <random>
#include <iostream>
#include <sdsl/vectors.hpp>

sdsl::int_vector<> random_data(size_t n, uint64_t u) {...}

int main() {
    auto data = random_data(10 * (1 << 20), 1 << 24);
    std::cout << sdsl::size_in_mega_bytes(data) << std::endl;

    sdsl::vlc_vector<sdsl::coder::elias_delta, 128> vlc(data);
    std::cout << sdsl::size_in_mega_bytes(vlc) << std::endl;

    sdsl::enc_vector<sdsl::coder::elias_delta, 128> enc_delta(data);
    std::cout << sdsl::size_in_mega_bytes(enc_delta) << std::endl;

    sdsl::enc_vector<sdsl::coder::elias_gamma, 128> enc_gamma(data);
    std::cout << sdsl::size_in_mega_bytes(enc_gamma) << std::endl;
    return 0;
}
```

80  
39.0282  
3.19339  
2.79236

# Plain bitvectors (`bit_vector`)

- A specialised version of `int_vector<1>`
- Mutable
  - `b[i] = 1`
  - `b.flip()`
- Bitwise operations between bitvectors `b1 |= b2` (also `&=`, `^=`)
- Auxiliary data structures extend a (static) bitvector functionality

Class	+Bits	Time
<code>rank_support_v</code>	$0.25n$	$O(1)$
<code>rank_support_scan</code>	64	$O(n)$
<code>select_support_mcl</code>	$\leq 0.2n$	$O(1)$
<code>select_support_scan</code>	64	$O(n)$

# Example 5: bit\_vector

```
#include <iostream>
#include <sds1/bit_vectors.hpp>

int main() {
    sds1::bit_vector b1 = {1, 1, 0, 1, 0, 0, 0, 1};
    sds1::bit_vector b2 = {0, 0, 1, 1, 0, 1, 0};
    b1 |= b2;
    b1.flip();
    std::cout << b1 << std::endl;

    sds1::bit_vector b(80 * (1 << 20), 0);
    for (size_t i = 0; i < b.size(); i += 100)
        b[i] = 1;
    std::cout << sds1::size_in_mega_bytes(b) << std::endl;
}
```

0000100

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# Recap

- **Integer vectors**

- `int_vector<>`, bit-width set at run-time, starts with 64 bits by default
- `int_vector<ℓ>`, bit-width  $\ell$  fixed at compile-time
- Implementation of random access

- **Compressed integer vectors**

- `vlc_vector<coder, density>`, vector of  $\gamma/\delta$ -coded integers
  - Implementation of random access via bit pointers
- `enc_vector<coder, density>`, vector of  $\gamma/\delta$ -coded gaps between ints
  - Implementation of random access via bit pointers and samples

- **Plain bitvectors**

- `bit_vector = int_vector<1>`
- Bitwise AND, OR, XOR
- Set individual bit, flip all bits

# Example 6: bit\_vector with rank/select

```
#include <iostream>
#include <sds1/bit_vectors.hpp>

int main() {
    sds1::bit_vector b = {0, 1, 0, 1, 1, 1, 0, 0, 0, 1, 1};
    sds1::rank_support_v<1> b_rank(&b);
    for (size_t i = 1; i <= b.size(); ++i)
        std::cout << b_rank(i) << " ";           counts #1s in b[0, i)
    std::cout << std::endl;

    sds1::select_support_mcl<1> b_select(&b);
    size_t ones = b_rank(b.size());
    for (size_t i = 1; i <= ones; ++i)
        std::cout << b_select(i) << " ";
}
```

Output

```
b = 0 1 0 1 1 1 0 0 0 1 1
0 1 1 2 3 4 4 4 4 5 6
1 3 4 5 9 10
```

# Elias-Fano compressed bitvectors (**sd\_vector**)

- Construct from a bitvector

```
sds1::bit_vector b = {0,1,0,1,1,1,0,0,0,1,1};  
sds1::sd_vector<> ef1(b);
```

- Construct from a vector of increasing integers

```
sds1::int_vector<> v = {1, 3, 4, 5, 9, 10};  
sds1::sd_vector<> ef2(v.begin(), v.end());
```

- **ef1 == ef2**

# Example 7: Elias-Fano

```
#include <iostream>
#include <sds1/bit_vectors.hpp>

sds1::int_vector<> random_data(size_t n, uint64_t u) {...}

int main() {
    auto data = random_data(10 * (1 << 20), 1 << 25);
    sds1::sd_vector<> ef(data.begin(), data.end());
    std::cout << sds1::size_in_mega_bytes(data) << std::endl;
    std::cout << sds1::size_in_mega_bytes(ef) << std::endl;
    return 0;
}
```

# Example 8: Elias-Fano internals

```
#include <iostream>
#include <sdsdl/bit_vectors.hpp>

int main() {
    sdsdl::int_vector<> v = {1, 4, 7, 18, 24, 26, 30, 31};
    sdsdl::sd_vector<> ef(v.begin(), v.end());
    uint64_t u = v[v.size() - 1] + 1;
    uint64_t bpi = std::ceil(std::log2(u));
    std::cout << "Data = " << v << std::endl
        << "Bitvector = " << ef << std::endl
        << "u = " << u << std::endl
        << "Bits per integer = " << bpi << std::endl
        << "Bits in the high part = " << bpi - ef.wl << std::endl
        << "Bits in the low part = " << (int) ef.wl << std::endl
        << "L = " << ef.low << " (in decimal)" << std::endl
        << "H = " << ef.high << std::endl;
    return 0;
}
```

```
Data = 1 4 7 18 24 26 30 31
Bitvector = 0100100100000000010000010100011
u = 32
Bits per integer = 5
Bits in the high part = 3
Bits in the low part = 2
L = 1 0 3 2 0 2 2 3 (in decimal)
H = 101100010011011000000000
```

1 =	000	01
4 =	001	00
7 =	001	11
18 =	100	10
24 =	110	00
26 =	110	10
30 =	111	10
31 =	111	11

Technical note: wrt the lecture notes, SDSL uses a different splitting point for the low/high parts (a difference of  $\pm 1$ ); but in this example the encodings match

$$h = \lceil \log n \rceil = 3$$
$$b = \lceil \log u \rceil = 5$$
$$\ell = b - h = 2$$

$$L = 0100111000101011$$

$$\begin{array}{c|c|c|c|c|c|c} \text{bucket} & 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 \\ \hline H = & 1011000 & 100 & 110 & 110 & \end{array}$$

# Example 9: Elias-Fano access and query

```
#include <iostream>
#include <sds1/bit_vectors.hpp>

int main() {
    sds1::int_vector<> v = {1, 4, 7, 18, 24, 26, 30, 31};
    sds1::sd_vector<> ef(v.begin(), v.end());
    std::cout << ef << std::endl;

    // Use as a bitvector with rank/select/operator[]
    sds1::rank_support_sd<> ef_rank(&ef);
    sds1::select_support_sd<> ef_select(&ef);
    std::cout << ef_rank(5) << std::endl;
    std::cout << ef_select(4) << std::endl;
    std::cout << ef[24] << std::endl; // access to the bit in pos 24

    // Use as an integer vector with access/nextGEQ given select/rank
    auto access = [&] (size_t i) { return ef_select(i + 1); };
    auto nextGEQ = [&] (uint64_t x) { return ef_select(ef_rank(x) + 1); };
    std::cout << access(5) << std::endl;
    std::cout << nextGEQ(4) << std::endl;
    std::cout << nextGEQ(27) << std::endl;
    return 0;
}
```

Output

```
0100100100000000010000010100011
2
18
1
26
4
30
```