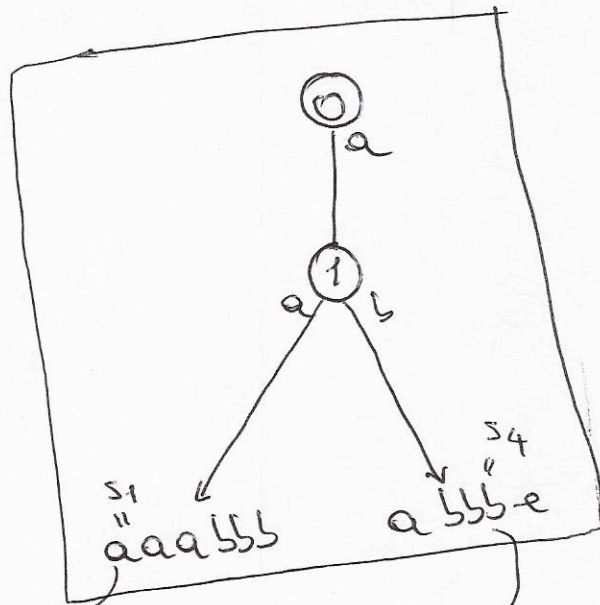


Exercise 1

FC =  $\langle 0, baci \rangle \langle 3, 0 \rangle \langle 1, oss \rangle \langle 1, uouuo \rangle \langle 2, u \rangle \langle 0, costi \rangle \langle 4, ola \rangle$

LPC =  $\langle 0, baaci \rangle \langle 0, bacoo \rangle \langle 0, bonn \rangle \langle 1, uouuo \rangle \langle 0, suu \rangle$   
 $\langle 0, costi \rangle \langle 4, ola \rangle$

Exercise 2



Two-level indexing is described in the notes: 9-5, 9-12

$\langle 0, \cancel{aaaabb} \rangle \langle 4, bd \rangle \langle 1, bbbd \rangle$

$\langle 0, \cancel{abbbe} \rangle \langle 3, c \rangle \langle 2, c \rangle$

The search for the string  $P = aaabcb$  consists of 3 phases:

1- downward: It matches the first two "a" of  $P$  and reaches the first leaf.

2- check: It computes  $LCP(P, aaaabb) = 4$  &  $P > s_1$

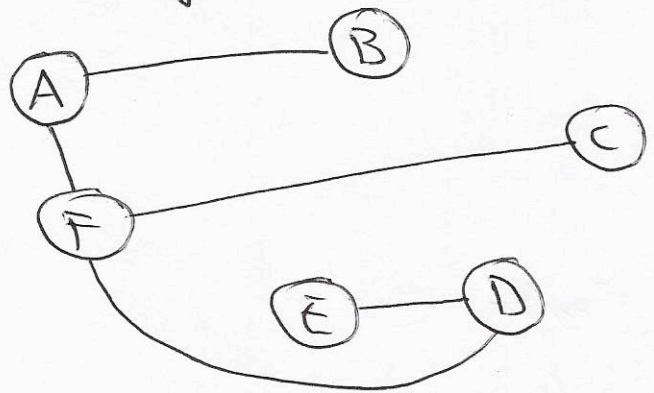
3- upward: It finds the partition of the mismatch, which is on the leftmost downward edge and since  $P > s_1$  it declares the lexicographic partition of  $P$  between  $s_1$  and  $s_4$ .

Thus it accesses the first block on disk, and finds that  $P$  occurs between  $s_1$  and  $s_2$ , by decompressing the front-coded strings of that block in 1 I/O.

### Exercise 3

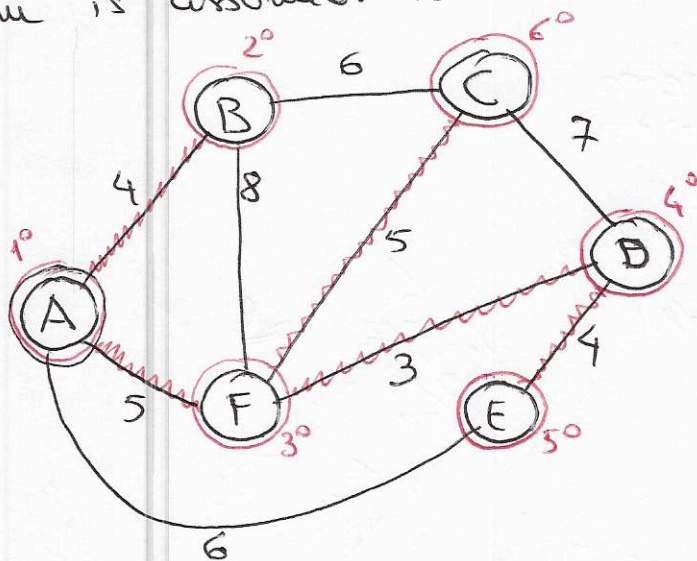
• Kruskal sorts edges by their increasing weight

- (D, F, 3) ✓
- (A, B, 4) ✓
- (D, E, 4) ✓
- (A, F, 5) ✓
- (C, F, 5) ✓
- (A, E, 6) ✗
- (B, C, 6) ✗
- (C, D, 7) ✗
- (B, F, 8) ✗



MST built by inserting edges in increasing weight and preventing that no cycle is formed, hence no edge internal to a component.

• Prim is assumed to start from node A



	A	B	C	D	E	F
1 <sup>o</sup>	0	∞	∞	∞	∞	∞
2 <sup>o</sup>	✓	4 <sub>A</sub>	∞	∞	6 <sub>A</sub>	5 <sub>A</sub>
3 <sup>o</sup>		✓	6 <sub>B</sub>	∞	6 <sub>A</sub>	5 <sub>A</sub>
4 <sup>o</sup>			5 <sub>F</sub>	3 <sub>F</sub>	6 <sub>A</sub>	✓
5 <sup>o</sup>			5 <sub>F</sub>	✓	4 <sub>D</sub>	
6 <sup>o</sup>			✓			

The MST is the same one computed by Kruskal's algorithm.