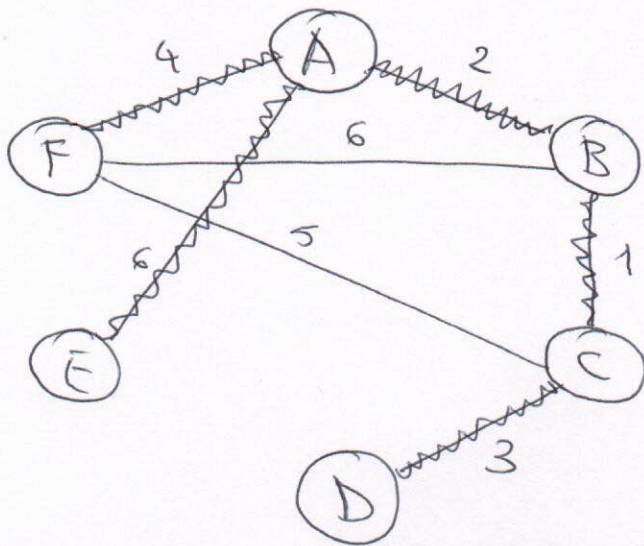


Exercise 1



min MST edge

Prim's Algorithm uses a priority queue indexed by the minimum-cost edge incident in ~~the~~ a node. Everything starts from the root A.

ops	A	B	C	D	E	F
	\ominus_A	∞	∞	∞	∞	∞
Extract A	\bullet	2_A	∞	∞	6_A	4_A
Extract B (A,B)		\bullet	1_B	∞	6_A	4_A
Extract C (B,C)			\bullet	3_C	6_A	4_A
Extract D (C,D)				\bullet	6_A	4_A
Extract F (A,F)					6_A	\bullet
Extract E (A,E)						

IF $n < M$ the PQ fits in internal memory and every node extraction takes $O(1)$, but the update of the PQ needs to scan an array (possibly) the adjacency list paying $\frac{|adj(u)|}{B}$ I/Os. If we sum over all nodes u , it is $|V| + |E|/B$ I/Os.

Exercice 2

1 amaca \$
 maca \$ a
 aca \$ am
 ea \$ ama
 a \$ amac
 \$ amaca

sort
 →

\$ amaca
 a \$ amac
 aca \$ am
 amese \$
 ca \$ ema
 mese \$ a

MTF (acmaa) with \$'s position = 4

$L = \{a, c, m\}$

a c m a a
 ↓ ↓ ↓ ↓ ↓
 1 2 3 3 1

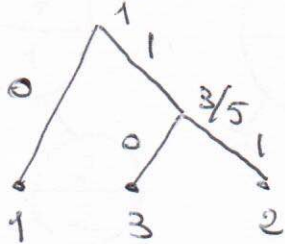
$L = \{c, a, m\}$

$L = \{m, c, a\}$

$L = \{a, m, c\}$

RLZ1 (1 2 3 3 1) → 1 2 3 3 1

Huffman (1 2 3 3 1) ⇒

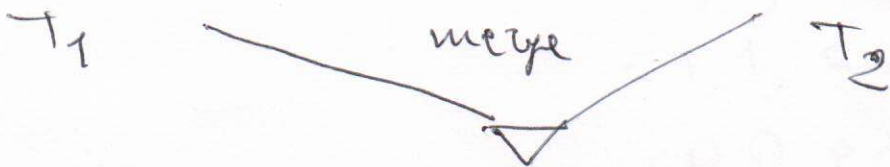
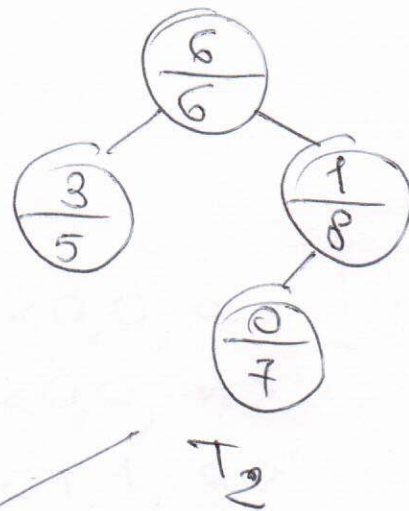
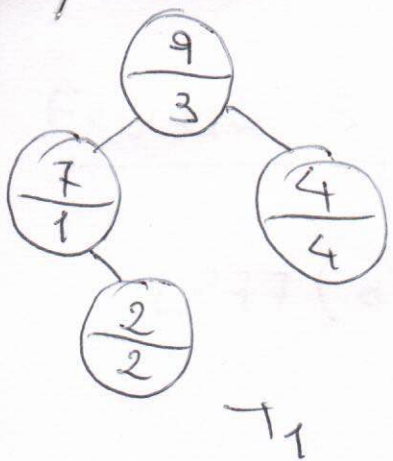


$f(1) = 2/5 \rightarrow 0$
 $f(2) = 1/5 \rightarrow 10$
 $f(3) = 2/5 \rightarrow 11$

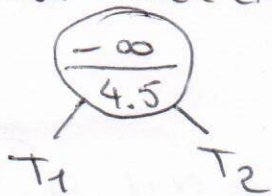
cw

Exercise 3

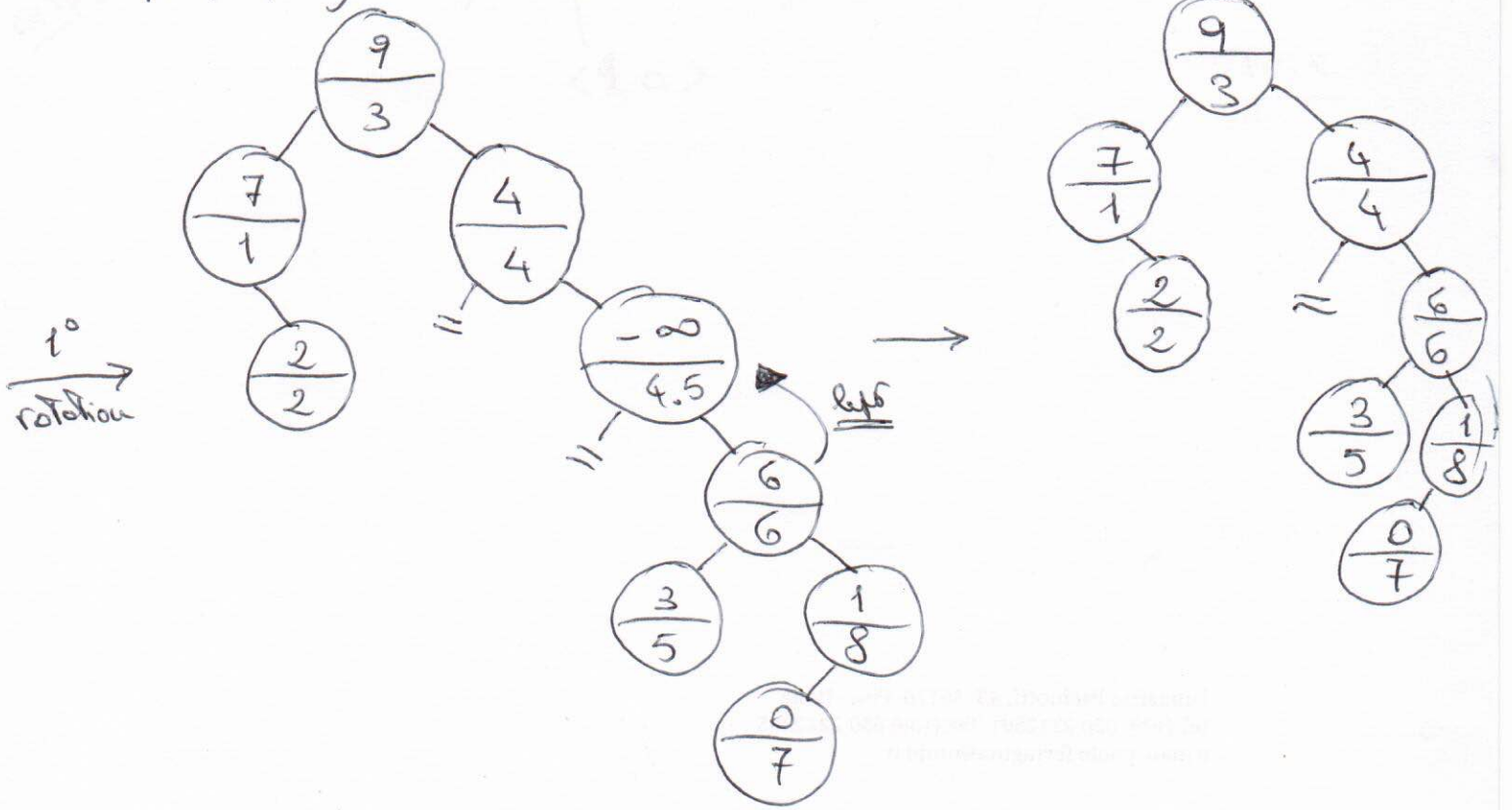
they are maximum heaps and so:



We create a dummy root with priority $-\infty$ and value between $\max(T_1) = 4$ and $\min(T_2) = 5$ so that the heap is correct in terms of key distribution.



that means it is a binary search tree. Then we make the dummy root percolate down to re-establish also the heap-property on the priorities based on rotations.



Exercice 4

see notes

Exercice 5

L277 (anne perine) =	$\langle 0, 0, a \rangle$	}	substrings
	$\langle 0, 0, n \rangle$		a
	$\langle 1, 1, a \rangle$		n
	$\langle 0, 0, p \rangle$		na
	$\langle 5, 4, EOF \rangle$		p
			anna

output

