

MPP 2026/27 (0077A, 9CFU)

Models for Programming Paradigms

Roberto Bruni

Filippo Bonchi

<http://www.di.unipi.it/~bruni/>



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03 - Unification

Some solutions

Conjectures vs theorems

Conjectures vs theorems

a natural number **p** is **prime**

if it cannot be written as the product of two smaller numbers

n	Is n prime?	$2^n - 1$	Is $2^n - 1$ prime?
2	yes	3	yes
3	yes	7	yes
4	no: $4 = 2 \cdot 2$	15	no: $15 = 3 \cdot 5$
5	yes	31	yes
6	no: $6 = 2 \cdot 3$	63	no: $63 = 7 \cdot 9$
7	yes	127	yes
8	no: $8 = 2 \cdot 4$	255	no: $255 = 15 \cdot 17$
9	no: $9 = 3 \cdot 3$	511	no: $511 = 7 \cdot 73$
10	no: $10 = 2 \cdot 5$	1023	no: $1023 = 31 \cdot 33$

Conjectures vs theorems

if p is prime
then $2^p - 1$ is prime



$p = 11$ is prime

$2^{11} - 1 = 2047$ is not prime

$$2047 = 23 \cdot 89$$

Use any mean to prove or disprove the above conjectures

Conjectures vs theorems

if $n > 1$ is not prime
then $2^n - 1$ is not prime



$n = a \cdot b$ with $a, b > 1$

$$2^a - 1 > 1$$

$$2^n - 1 = 2^{a \cdot b} - 1 = (2^a)^b - 1 =$$

$$(2^a)^b - (2^a)^0 = (2^a)^b - 1 =$$

$$(2^a)^b - 1 = (2^a)^b - (2^a)^0 =$$

$$2^a \left((2^a)^{b-1} + \dots + (2^a)^2 + 2^a + 1 \right) -$$

$$\left((2^a)^{b-1} + \dots + (2^a)^2 + 2^a + 1 \right) =$$

$$(2^a - 1) \left((2^a)^{b-1} + \dots + (2^a)^2 + 2^a + 1 \right)$$

Use any mean to prove or disprove the above conjectures

Exercise

Expressions with variables $E ::= x \mid N \mid E \oplus E \mid E \otimes E$

How to evaluate expressions such as $(x \oplus 4) \otimes y$?

Need some memories $\mathbb{M} \triangleq \{\sigma \mid \sigma : X \rightarrow \mathbb{N}\}$

machine states $\langle E, \sigma \rangle$

interpretation function $\mathcal{E}[\![\cdot]\!] : Exp \rightarrow (\mathbb{M} \rightarrow \mathbb{N})$

Let's redefine the various semantics and properties

Exercise

$$E ::= x \mid N \mid E \oplus E \mid E \otimes E$$

$$\mathbb{M} \triangleq \{\sigma \mid \sigma : X \rightarrow \mathbb{N}\}$$

$$(num) \frac{}{\langle N, \sigma \rangle \longrightarrow n}$$

$$(var) \frac{}{\langle x, \sigma \rangle \longrightarrow \sigma(x)}$$

$$(sum) \frac{\langle E_0, \sigma \rangle \longrightarrow n_0 \quad \langle E_1, \sigma \rangle \longrightarrow n_1}{\langle E_0 \oplus E_1, \sigma \rangle \longrightarrow n} \quad n = n_0 + n_1$$

$$(prod) \frac{\langle E_0, \sigma \rangle \longrightarrow n_0 \quad \langle E_1, \sigma \rangle \longrightarrow n_1}{\langle E_0 \otimes E_1, \sigma \rangle \longrightarrow n} \quad n = n_0 \times n_1$$

Exercise

$$E ::= x \mid N \mid E \oplus E \mid E \otimes E \qquad \mathbb{M} \triangleq \{\sigma \mid \sigma : X \rightarrow \mathbb{N}\}$$

$$\mathcal{E}[\![\cdot]\!] : Exp \rightarrow (\mathbb{M} \rightarrow \mathbb{N})$$

$$\mathcal{E}[\![N]\!]\sigma = n$$

$$\mathcal{E}[\![E_0 \oplus E_1]\!]\sigma = \mathcal{E}[\![E_0]\!]\sigma + \mathcal{E}[\![E_1]\!]\sigma$$

$$\mathcal{E}[\![E_0 \otimes E_1]\!]\sigma = \mathcal{E}[\![E_0]\!]\sigma \times \mathcal{E}[\![E_1]\!]\sigma$$

$$\mathcal{E}[\![x]\!]\sigma = \sigma(x)$$

Exercise

$$E ::= x \mid N \mid E \oplus E \mid E \otimes E$$

$$\mathbb{M} \triangleq \{\sigma \mid \sigma : X \rightarrow \mathbb{N}\}$$

termination? $\forall E. \forall \sigma. \exists n. \langle E, \sigma \rangle \longrightarrow n$

$$\forall E. \forall \sigma. \exists n. \mathcal{E}[[E]]\sigma = n$$

determinacy?

$$\forall E, \sigma, n, m. \left(\begin{array}{c} \langle E, \sigma \rangle \longrightarrow n \\ \wedge \\ \langle E, \sigma \rangle \longrightarrow m \end{array} \right) \Rightarrow n = m$$

$$\forall E, \sigma, n, m. \left(\begin{array}{c} \mathcal{E}[[E]]\sigma = n \\ \wedge \\ \mathcal{E}[[E]]\sigma = m \end{array} \right) \Rightarrow n = m$$

Inference

SOS rule application?

1. a goal

$$(1 \oplus 2) \otimes (3 \oplus 4) \longrightarrow m$$

SOS rule application?

$$(prod) \frac{E_0 \longrightarrow n_0 \quad E_1 \longrightarrow n_1}{E_0 \otimes E_1 \longrightarrow n} \quad n = n_0 \cdot n_1$$

2. take a rule

$$(1 \oplus 2) \otimes (3 \oplus 4) \longrightarrow m$$

SOS rule application?

$$(prod) \frac{E_0 \longrightarrow n_0 \quad E_1 \longrightarrow n_1}{E_0 \otimes E_1 \longrightarrow n} \quad n = n_0 \cdot n_1$$

3. unify
(if possible)

$$\begin{aligned} E_0 &= 1 \oplus 2 \\ E_1 &= 3 \oplus 4 \\ n &= m \end{aligned}$$

$$(1 \oplus 2) \otimes (3 \oplus 4) \longrightarrow m$$

SOS rule application?

$$(prod) \frac{1 \oplus 2 \longrightarrow n_0 \quad 3 \oplus 4 \longrightarrow n_1}{(1 \oplus 2) \otimes (3 \oplus 4) \longrightarrow m} \quad m = n_0 \cdot n_1$$

4. instantiate

$$(1 \oplus 2) \otimes (3 \oplus 4) \longrightarrow m$$

SOS rule application?

$$(prod) \frac{\boxed{1 \oplus 2 \longrightarrow n_0} \quad \boxed{3 \oplus 4 \longrightarrow n_1}}{(1 \oplus 2) \otimes (3 \oplus 4) \longrightarrow m} \quad m = n_0 \cdot n_1$$

5. recursively solve subgoals

$$(1 \oplus 2) \otimes (3 \oplus 4) \longrightarrow m$$

SOS rule application?

$$(\text{prod}) \frac{1 \oplus 2 \longrightarrow \boxed{3} \quad 3 \oplus 4 \longrightarrow \boxed{7}}{(1 \oplus 2) \otimes (3 \oplus 4) \longrightarrow m} \quad m = \boxed{3 \cdot 7}$$

6. combine results

$$(1 \oplus 2) \otimes (3 \oplus 4) \longrightarrow m$$

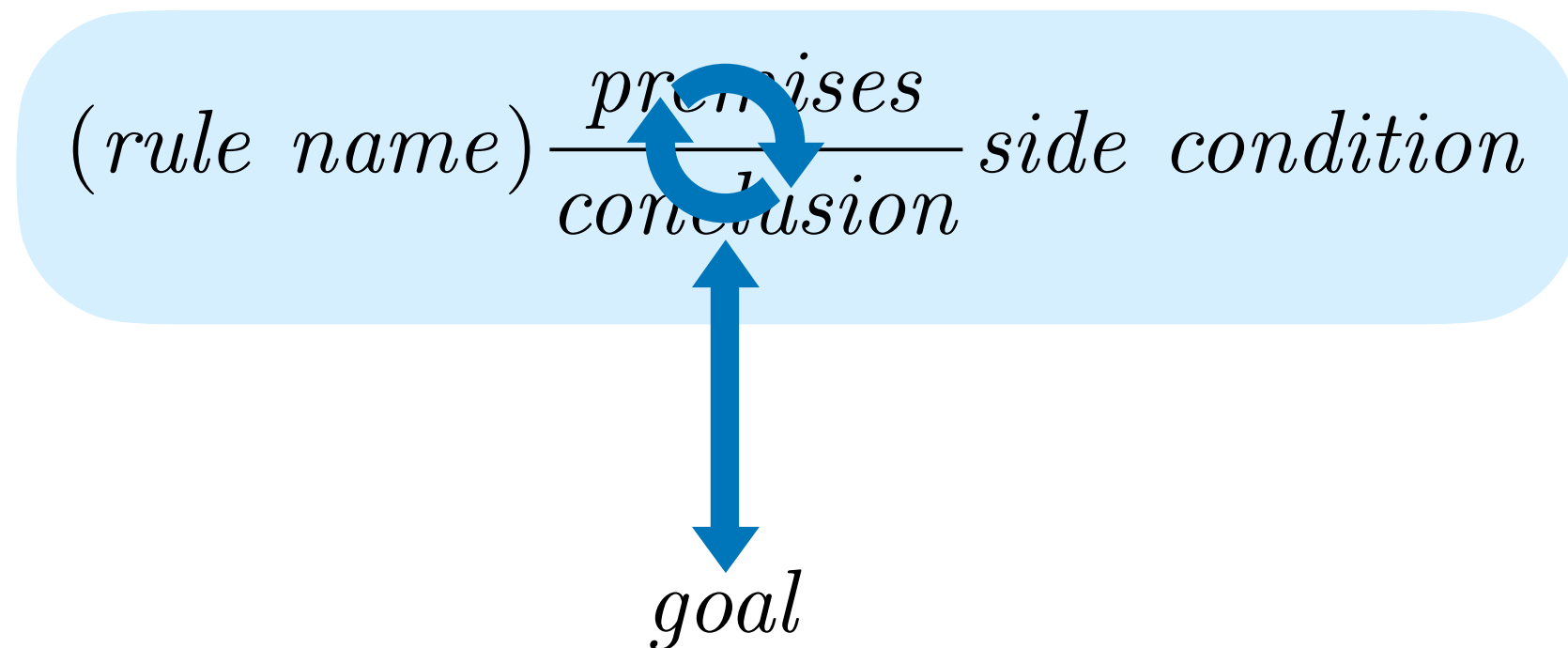
SOS rule application?

$$(\textit{prod}) \frac{1 \oplus 2 \longrightarrow 3 \quad 3 \oplus 4 \longrightarrow 7}{(1 \oplus 2) \otimes (3 \oplus 4) \longrightarrow \boxed{21}}$$

7. return results

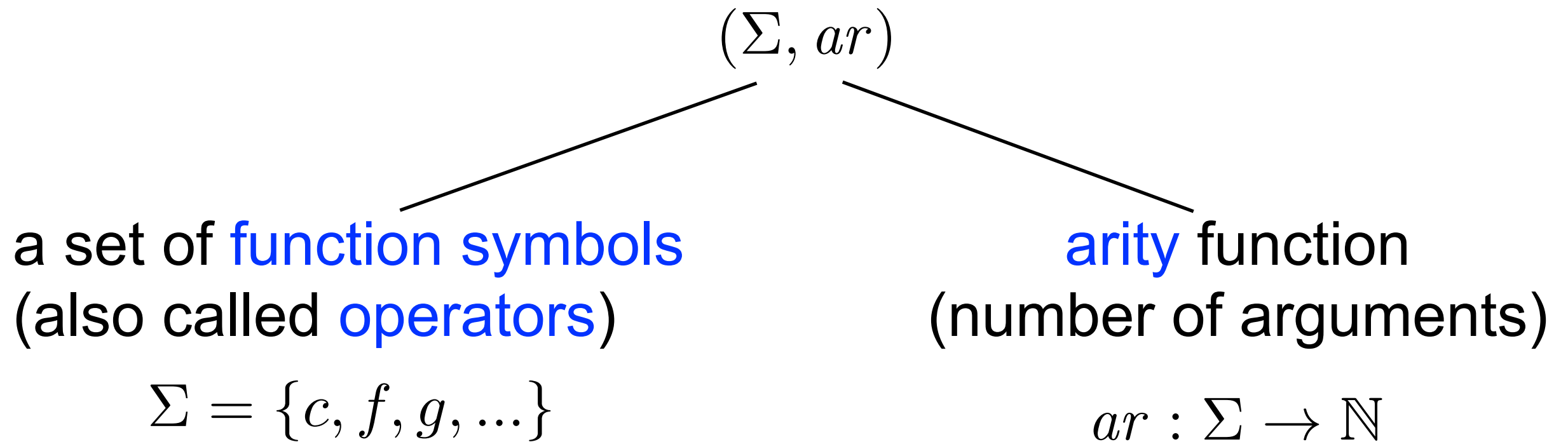
$$(1 \oplus 2) \otimes (3 \oplus 4) \longrightarrow \boxed{21}$$

Deduction process



Signatures

Unsorted signature



each function symbol has an arity

Running example

$\Sigma = \{0, \text{succ}, \text{plus}\}$

$ar(0) = 0$ constant

$ar(\text{succ}) = 1$ unary

$ar(\text{plus}) = 2$ binary

Equivalently

$$(\Sigma, ar)$$

$$\begin{aligned} \text{let } \Sigma_n &\triangleq ar^{-1}(n) \\ &= \{f \in \Sigma \mid ar(f) = n\} \end{aligned}$$

a signature is an arity-indexed family of sets of operators

$$\Sigma = \{\Sigma_n\}_{n \in \mathbb{N}}$$

Running example

$$\Sigma_0 = \{0\}$$

$$\Sigma_1 = \{\text{succ}\}$$

$$\Sigma_2 = \{\text{plus}\}$$

$$\Sigma_n = \emptyset \quad \text{if } n > 2$$

Terms over a signature

$\Sigma = \{\Sigma_n\}_{n \in \mathbb{N}}$ a signature

$X = \{x, y, z, \dots\}$ an infinite set of **variables**

$T_{\Sigma, X}$ denotes the set of all **terms** over Σ, X

$T_{\Sigma, X}$ is the least set such that:

- if $x \in X$, then $x \in T_{\Sigma, X}$
- if $c \in \Sigma_0$, then $c \in T_{\Sigma, X}$
- if $f \in \Sigma_n$ and $t_1, \dots, t_n \in T_{\Sigma, X}$, then $f(t_1, \dots, t_n) \in T_{\Sigma, X}$

i.e. $T_{\Sigma, X} \ni t ::= x \mid c \mid f(t_1, \dots, t_n)$

$x \in X \quad c \in \Sigma_0 \quad f \in \Sigma_n$

Vars

$$\Sigma = \{\Sigma_n\}_{n \in \mathbb{N}} \quad X = \{x, y, z, \dots\} \quad t \in T_{\Sigma, X}$$

$vars(t)$ set of variables that appears in the term t

$$vars : T_{\Sigma, X} \rightarrow \wp(X)$$

$$vars(x) \triangleq \{x\}$$

$$vars(c) \triangleq \emptyset$$

$$vars(f(t_1, \dots, t_n)) \triangleq \bigcup_{i=1}^n vars(t_i)$$

Closed terms

a term with no variables is called **closed**

$$T_{\Sigma} \triangleq \text{vars}^{-1}(\emptyset) = \{t \in T_{\Sigma, X} \mid \text{vars}(t) = \emptyset\}$$

(obviously $T_{\Sigma} \subseteq T_{\Sigma, X}$)

Skill levels



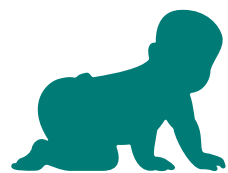
advanced
multitasking



intermediate
does many things



novice
does one thing



beginner
does nothing

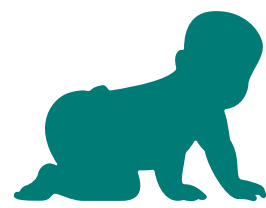
Running example

$$\Sigma_0 = \{0\}$$

$$\Sigma_1 = \{\text{succ}\}$$

$$\Sigma_2 = \{\text{plus}\}$$

$$\Sigma_n = \emptyset \quad \text{if } n > 2$$



Exercise

let's complete the schema
(obviously $T_\Sigma \subseteq T_{\Sigma, X}$)



	t	$t \in ?$	$vars(t)$
	0	$\begin{array}{ c c } \hline \bullet & \circ \\ \hline T_\Sigma & T_{\Sigma, X} \\ \hline \end{array}$	$\emptyset$
	x	$\begin{array}{ c c } \hline \circ & \bullet \\ \hline T_\Sigma & T_{\Sigma, X} \\ \hline \end{array}$	$\{x\}$
	$\text{succ}(0)$	$\begin{array}{ c c } \hline \circ & \circ \\ \hline T_\Sigma & T_{\Sigma, X} \\ \hline \end{array}$	
	$\text{succ}(x)$	$\begin{array}{ c c } \hline \circ & \circ \\ \hline T_\Sigma & T_{\Sigma, X} \\ \hline \end{array}$	
	$\text{succ}(\text{plus}(0), x)$	$\begin{array}{ c c } \hline \circ & \circ \\ \hline T_\Sigma & T_{\Sigma, X} \\ \hline \end{array}$	
	$\text{plus}(\text{succ}(x), 0)$	$\begin{array}{ c c } \hline \circ & \circ \\ \hline T_\Sigma & T_{\Sigma, X} \\ \hline \end{array}$	
	$\text{succ}(\text{succ}(0), \text{plus}(x))$	$\begin{array}{ c c } \hline \circ & \circ \\ \hline T_\Sigma & T_{\Sigma, X} \\ \hline \end{array}$	
	$\text{succ}(\text{plus}(w, z))$	$\begin{array}{ c c } \hline \circ & \circ \\ \hline T_\Sigma & T_{\Sigma, X} \\ \hline \end{array}$	
	$\text{plus}(\text{plus}(x, \text{succ}(y)), \text{plus}(0, \text{succ}(x)))$	$\begin{array}{ c c } \hline \circ & \circ \\ \hline T_\Sigma & T_{\Sigma, X} \\ \hline \end{array}$	

Substitutions

$\rho : X \rightarrow T_{\Sigma, X}$ a **substitution** assigns terms to variables

we only consider substitutions that are identity everywhere, except for a finite number of cases, written

$$\rho = [x_1 = t_1, \dots, x_n = t_n] \qquad \rho(x) = \begin{cases} t_i & \text{if } x = x_i \\ x & \text{otherwise} \end{cases}$$

all different

Notation

$\rho : X \rightarrow T_{\Sigma, X}$ a **substitution** assigns terms to variables

overloaded notation for the lifted function $\rho : T_{\Sigma, X} \rightarrow T_{\Sigma, X}$

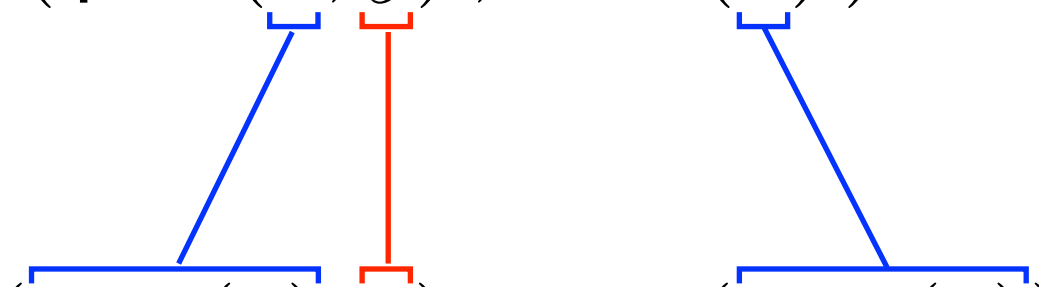
$\rho(t)$ denotes the term obtained by simultaneous application of the substitution to all variable occurrences in t

$t\rho$ alternative notation

Example

$$\rho \triangleq [x = \text{succ}(y), y = 0]$$

$$t \triangleq \text{plus}(\text{plus}(x, y), \text{succ}(x))$$



$$t\rho = \text{plus}(\text{plus}(\text{succ}(y), 0), \text{succ}(\text{succ}(y)))$$

mgt relation

t is more general than t' if $\exists \rho. t' = t\rho$

in which case, we also say t' is an instance of t

$\text{plus}(x, \text{succ}(y))$	mgt	$\text{plus}(0, \text{succ}(\text{succ}(z)))$
$\text{plus}(0, x)$	mgt	$\text{plus}(y, 0)$
$\text{plus}(y, 0)$	mgt	$\text{plus}(0, x)$
$\text{plus}(0, x)$ $\text{plus}(y, 0)$	mgt	$\text{plus}(0, 0)$

mgt relation

mgt is transitive and reflexive

$$\begin{array}{c} \diagup \\ t \text{ mgt } t \end{array}$$

if $(t_1 \text{ mgt } t_2)$ and $(t_2 \text{ mgt } t_3)$, then $(t_1 \text{ mgt } t_3)$

there are terms $t \neq t'$ such that $(t \text{ mgt } t')$ and $(t' \text{ mgt } t)$

$\text{succ}(x) \quad \text{succ}(y)$

mgt can be extended to substitutions pointwise

$\rho \text{ mgt } \rho'$ if $\exists \rho''. \forall x. \rho'(x) = \rho''(\rho(x))$

Unification

(in its simplest form: syntactic, first-order)

Unification problem

Given a set of potential equalities

$$\mathcal{G} = \{\ell_1 \stackrel{?}{=} r_1, \dots, \ell_n \stackrel{?}{=} r_n\}$$

where $\ell_1, \dots, \ell_n, r_1, \dots, r_n \in T_{\Sigma, X}$

can we find a substitution ρ such that

$$\forall i \in [1, n]. \rho(\ell_i) = \rho(r_i) \quad ?$$

we call such a ρ a solution of $\mathcal{G}$

$$\text{sols}(\mathcal{G}) \triangleq \{\rho \mid \forall i \in [1, n]. \rho(\ell_i) = \rho(r_i)\}$$

Intuitively



Intuitively



Examples

$$\mathcal{G} = \{\text{plus}(x, 0) \stackrel{?}{=} \text{plus}(0, y)\}$$

$$\text{sols}(\mathcal{G}) \triangleq \{[x = 0, y = 0]\}$$

$$\mathcal{G} = \{\text{succ}(x) \stackrel{?}{=} \text{succ}(\text{succ}(y))\}$$

$$\text{sols}(\mathcal{G}) \triangleq \{[x = \text{succ}(y)], [x = \text{succ}(0), y = 0], \dots\}$$

Unification problem

More interestingly, can we solve the following problem?

Given a set of potential equalities

$$\mathcal{G} = \{\ell_1 \stackrel{?}{=} r_1, \dots, \ell_n \stackrel{?}{=} r_n\}$$

where $\ell_1, \dots, \ell_n, r_1, \dots, r_n \in T_{\Sigma, X}$

can we find a most general solution ρ ?

$$\rho \in \text{sols}(\mathcal{G}) \quad \text{and}$$

$$\forall \rho' \in \text{sols}(\mathcal{G}). \quad \rho \text{ mgt } \rho'$$

Unification algorithm

Idea: we iteratively reduce the set $\mathcal{G}$
by solution-preserving transformations until
either a solution is found
or we can prove there is no solution

$$\mathcal{G} \cdots \mathcal{G}_1 \cdots \mathcal{G}_n \cdots \{x_1 \stackrel{?}{=} t_1, \dots, x_k \stackrel{?}{=} t_k\}$$



Solutions may not exist
and even if they exist may not be unique

Termination

$$\mathcal{G} = \{\ell_1 \stackrel{?}{=} r_1, \dots, \ell_n \stackrel{?}{=} r_n\}$$

$\mathcal{G}$ and $\mathcal{G}'$ are equivalent if $\text{sols}(\mathcal{G}) = \text{sols}(\mathcal{G}')$

the algorithm terminates successfully when we reach

$$\mathcal{G}' = \{x_1 \stackrel{?}{=} t_1, \dots, x_k \stackrel{?}{=} t_k\} \text{ equivalent to } \mathcal{G}$$

all different
 k

$$\{x_1, \dots, x_k\} \cap \bigcup_{i=1}^k \text{vars}(t_i) = \emptyset$$

any such $\mathcal{G}'$ defines a straightforward solution $[x_1 = t_1, \dots, x_k = t_k]$

Notation

$$\mathcal{G} = \{\ell_1 \stackrel{?}{=} r_1, \dots, \ell_n \stackrel{?}{=} r_n\}$$

$$vars(\mathcal{G}) \stackrel{\Delta}{=} \bigcup_{i=1}^n (vars(\ell_i) \cup vars(r_i))$$

$$\mathcal{G}\rho \stackrel{\Delta}{=} \{\ell_1\rho \stackrel{?}{=} r_1\rho, \dots, \ell_n\rho \stackrel{?}{=} r_n\rho\}$$

Unification algorithm

delete

$\mathcal{G} \cup \{t \stackrel{?}{=} t\}$
becomes
 $\mathcal{G}$

eliminate

$\mathcal{G} \cup \{x \stackrel{?}{=} t\}$
becomes if $x \in \text{vars}(\mathcal{G}) \setminus \text{vars}(t)$
 $\mathcal{G}[x = t] \cup \{x \stackrel{?}{=} t\}$

swap

$\mathcal{G} \cup \{f(t_1, \dots, t_m) \stackrel{?}{=} x\}$
becomes
 $\mathcal{G} \cup \{x \stackrel{?}{=} f(t_1, \dots, t_m)\}$

decompose

$\mathcal{G} \cup \{f(t_1, \dots, t_m) \stackrel{?}{=} f(u_1, \dots, u_m)\}$
becomes
 $\mathcal{G} \cup \{t_1 \stackrel{?}{=} u_1, \dots, t_m \stackrel{?}{=} u_m\}$

occur-check

$\mathcal{G} \cup \{x \stackrel{?}{=} f(t_1, \dots, t_m)\}$
fails if $x \in \text{vars}(f(t_1, \dots, t_m))$

conflict

$\mathcal{G} \cup \{f(t_1, \dots, t_m) \stackrel{?}{=} g(u_1, \dots, u_h)\}$
fails if $f \neq g$ or $m \neq h$

Example

$$\{\text{plus}(\text{succ}(x), x) \stackrel{?}{=} \text{plus}(y, 0)\}$$

decompose

$$\mathcal{G} \cup \{f(t_1, \dots, t_m) \stackrel{?}{=} f(u_1, \dots, u_m)\}$$

becomes

$$\mathcal{G} \cup \{t_1 \stackrel{?}{=} u_1, \dots, t_m \stackrel{?}{=} u_m\}$$

$$\{\text{succ}(x) \stackrel{?}{=} y, x \stackrel{?}{=} 0\}$$

eliminate

$$\mathcal{G} \cup \{x \stackrel{?}{=} t\}$$

becomes if $x \in \text{vars}(\mathcal{G}) \setminus \text{vars}(t)$

$$\mathcal{G}[x = t] \cup \{x \stackrel{?}{=} t\}$$

$$\{\text{succ}(0) \stackrel{?}{=} y, x \stackrel{?}{=} 0\}$$

swap

$$\mathcal{G} \cup \{f(t_1, \dots, t_m) \stackrel{?}{=} x\}$$

becomes

$$\mathcal{G} \cup \{x \stackrel{?}{=} f(t_1, \dots, t_m)\}$$

$$\{y \stackrel{?}{=} \text{succ}(0), x \stackrel{?}{=} 0\}$$

success!

$$\rho = [y = \text{succ}(0), x = 0]$$

Example

$$\{\text{plus}(0, x) \stackrel{?}{=} \text{succ}(y)\} \quad \text{plus} \neq \text{succ}$$

conflict

$$\mathcal{G} \cup \{f(t_1, \dots, t_m) \stackrel{?}{=} g(u_1, \dots, u_h)\}$$

fails if $f \neq g$ or $m \neq h$

failure!

Example

$$\{\text{succ}(x) \stackrel{?}{=} y, \text{succ}(y) \stackrel{?}{=} x\}$$

swap

$$\mathcal{G} \cup \{f(t_1, \dots, t_m) \stackrel{?}{=} x\}$$

becomes

$$\mathcal{G} \cup \{x \stackrel{?}{=} f(t_1, \dots, t_m)\}$$

$$\{\text{succ}(x) \stackrel{?}{=} y, x \stackrel{?}{=} \text{succ}(y)\}$$

eliminate

$$\mathcal{G} \cup \{x \stackrel{?}{=} t\}$$

becomes if $x \in \text{vars}(\mathcal{G}) \setminus \text{vars}(t)$

$$\mathcal{G}[x = t] \cup \{x \stackrel{?}{=} t\}$$

$$\{\text{succ}(\text{succ}(y)) \stackrel{?}{=} y, x \stackrel{?}{=} \text{succ}(y)\}$$

swap

$$\mathcal{G} \cup \{f(t_1, \dots, t_m) \stackrel{?}{=} x\}$$

becomes

$$\mathcal{G} \cup \{x \stackrel{?}{=} f(t_1, \dots, t_m)\}$$

$$\{y \stackrel{?}{=} \text{succ}(\text{succ}(y)), x \stackrel{?}{=} \text{succ}(y)\}$$

$$y \in \text{vars}(\text{succ}(\text{succ}(y)))$$

occur-check

$$\mathcal{G} \cup \{x \stackrel{?}{=} f(t_1, \dots, t_m)\}$$

fails if $x \in \text{vars}(f(t_1, \dots, t_m))$

failure!



Exercise

$$\{\text{plus}(x, \text{succ}(x)) \stackrel{?}{=} \text{plus}(0, y), \text{plus}(y, z) \stackrel{?}{=} \text{plus}(z, w)\}$$

delete

$$\mathcal{G} \cup \{t \stackrel{?}{=} t\} \\ \text{becomes} \\ \mathcal{G}$$

eliminate

$$\mathcal{G} \cup \{x \stackrel{?}{=} t\} \\ \text{becomes} \quad \text{if } x \in \text{vars}(\mathcal{G}) \setminus \text{vars}(t) \\ \mathcal{G}[x = t] \cup \{x \stackrel{?}{=} t\}$$

swap

$$\mathcal{G} \cup \{f(t_1, \dots, t_m) \stackrel{?}{=} x\} \\ \text{becomes} \\ \mathcal{G} \cup \{x \stackrel{?}{=} f(t_1, \dots, t_m)\}$$

decompose

$$\mathcal{G} \cup \{f(t_1, \dots, t_m) \stackrel{?}{=} f(u_1, \dots, u_m)\} \\ \text{becomes} \\ \mathcal{G} \cup \{t_1 \stackrel{?}{=} u_1, \dots, t_m \stackrel{?}{=} u_m\}$$

occur-check

$$\mathcal{G} \cup \{x \stackrel{?}{=} f(t_1, \dots, t_m)\} \\ \text{fails if } x \in \text{vars}(f(t_1, \dots, t_m))$$

conflict

$$\mathcal{G} \cup \{f(t_1, \dots, t_m) \stackrel{?}{=} g(u_1, \dots, u_h)\} \\ \text{fails if } f \neq g \text{ or } m \neq h$$