

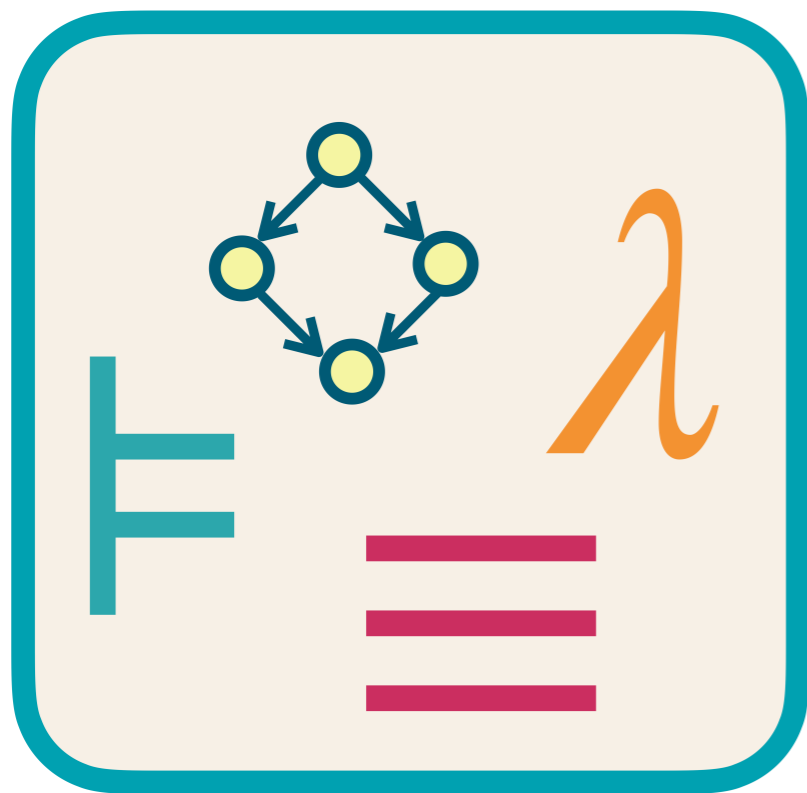
**MPP 2025/26** (0077A, 9CFU)

Models for Programming Paradigms

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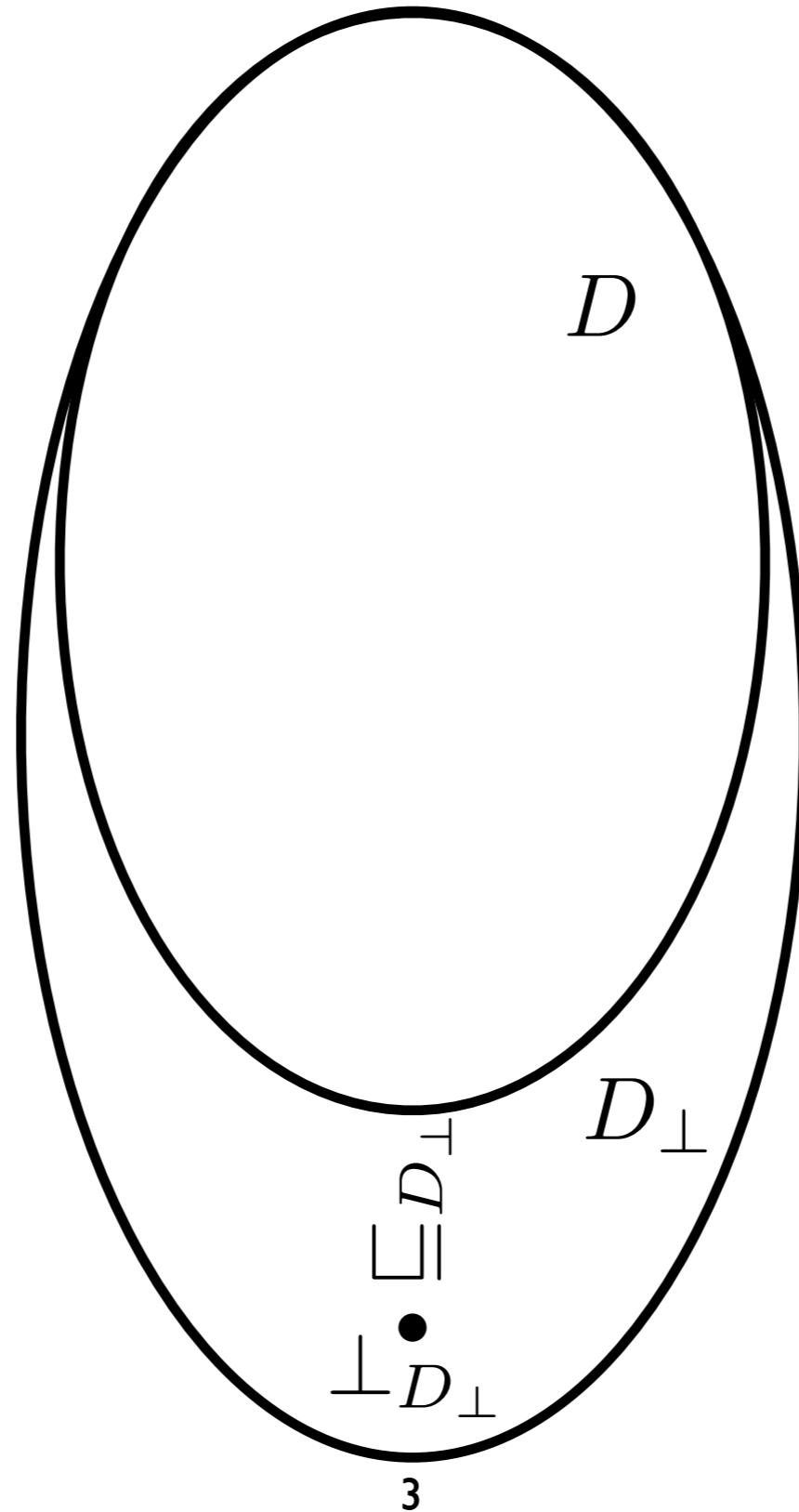


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13b - Continuity theorems

# Lifted Domains

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# Lifted Domains

$$\mathcal{D} = (D, \sqsubseteq_D) \text{ CPO} \quad \Rightarrow \quad \mathcal{D}_\perp = (D_\perp, \sqsubseteq_{D_\perp})$$

$$D_\perp \triangleq \{\perp\} \uplus D$$

$$= \{(0, \perp)\} \cup (\{1\} \times D) = \{(0, \perp)\} \cup \{(1, d) \mid d \in D\}$$

$$\perp_{D_\perp} \triangleq (0, \perp)$$

$$[\cdot] : D \rightarrow D_\perp$$

lifting function

$$[d] \triangleq (1, d)$$

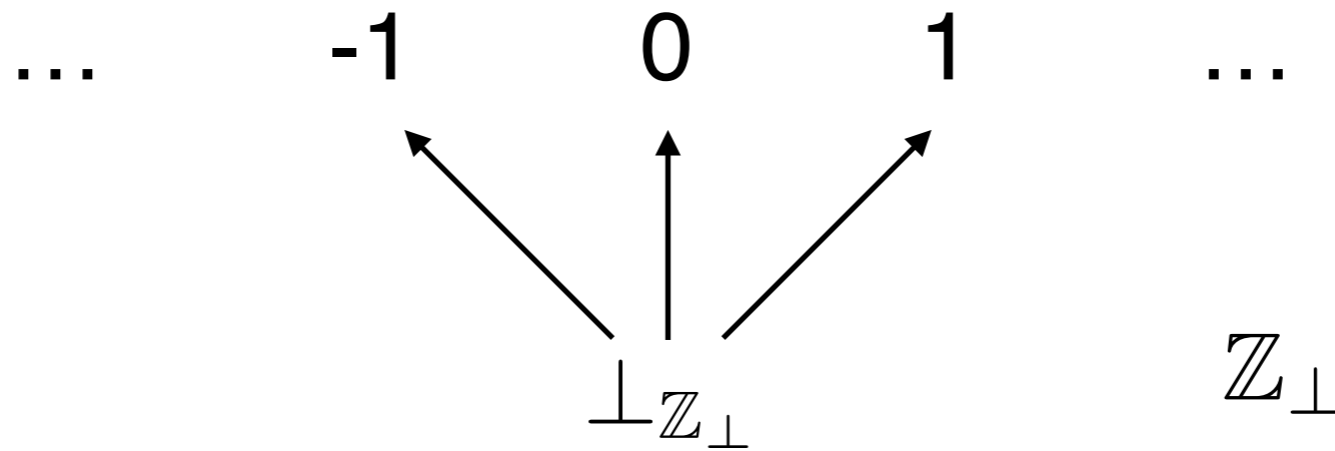
how to order lifted elements?

$$\forall x \in D_\perp. \perp_{D_\perp} \sqsubseteq_{D_\perp} x$$

$$\forall d_1, d_2 \in D. [d_1] \sqsubseteq_{D_\perp} [d_2] \Leftrightarrow d_1 \sqsubseteq_D d_2$$

# Example

$(\mathbb{Z}, =)$



# Lifted Domains

**TH.**  $\mathcal{D}_\perp = ( D_\perp , \sqsubseteq_{D_\perp} ) \text{ CPO}_\perp$

try on your own to prove:

PO,

bottom element,

complete

observe that:  $\bigsqcup_{i \in \mathbb{N}} \lfloor d_i \rfloor = \left\lfloor \bigsqcup_{i \in \mathbb{N}} d_i \right\rfloor$

it is an upper bound

it is the least upper bound

# Lifting operator

$(D, \sqsubseteq_D)$  CPO

$(E, \sqsubseteq_E)$  CPO $_{\perp}$

$$(\cdot)^* : [D \rightarrow E] \rightarrow [D_{\perp} \rightarrow E]$$

$$\forall f \in [D \rightarrow E]. f^*(x) \triangleq \begin{cases} \perp_E & \text{if } x = \perp_{D_{\perp}} \\ f(d) & \text{if } x = [d] \end{cases}$$

for the definition to be well-given  
we need to prove:

$$f \in [D \rightarrow E] \quad \Rightarrow \quad f^* \in [D_{\perp} \rightarrow E]$$

$f$  continuous implies  $f^*$  continuous

**TH.** the lifting operator is well-defined

*proof.* assume  $f$  continuous, take a chain  $\{x_n\}_{n \in \mathbb{N}}$  in  $D_\perp$

we need to prove  $f^* \left( \bigsqcup_{n \in \mathbb{N}} x_n \right) = \bigsqcup_{n \in \mathbb{N}} f^*(x_n)$

if  $\forall n \in \mathbb{N}. x_n = \perp_{D_\perp}$  then it is obvious

otherwise, let  $k = \min\{i \mid x_i \neq \perp_{D_\perp}\}$

then  $\forall m \geq k. \exists d_m \in D. x_m = \lfloor d_m \rfloor$

and by prefix independence of lub

$$\bigsqcup_{n \in \mathbb{N}} x_n = \bigsqcup_{n \in \mathbb{N}} x_{n+k} \quad \bigsqcup_{n \in \mathbb{N}} f^*(x_n) = \bigsqcup_{n \in \mathbb{N}} f^*(x_{n+k})$$

we can just prove  $f^* \left( \bigsqcup_{n \in \mathbb{N}} x_{n+k} \right) = \bigsqcup_{n \in \mathbb{N}} f^*(x_{n+k})$

(see next slide)



(continue)

$$f^* \left( \bigsqcup_{n \in \mathbb{N}} x_{n+k} \right) = \bigsqcup_{n \in \mathbb{N}} f^*(x_{n+k})$$

$$f^* \left( \bigsqcup_{n \in \mathbb{N}} x_{n+k} \right) = f^* \left( \bigsqcup_{n \in \mathbb{N}} \lfloor d_{n+k} \rfloor \right)$$

by def of  $k$

$$= f^* \left( \left[ \bigsqcup_{n \in \mathbb{N}} d_{n+k} \right] \right)$$

by lub in a lifted domain

$$= f \left( \bigsqcup_{n \in \mathbb{N}} d_{n+k} \right)$$

by def of lifting

$$= \bigsqcup_{n \in \mathbb{N}} f(d_{n+k})$$

by continuity of  $f$

$$= \bigsqcup_{n \in \mathbb{N}} f^*(\lfloor d_{n+k} \rfloor)$$

by def of lifting

$$= \bigsqcup_{n \in \mathbb{N}} f^*(x_{n+k})$$

by def of  $k$

**TH.**  $(\cdot)^*$  is monotone

(try to prove on your own)

**TH.**  $(\cdot)^*$  is continuous

*proof.* take a chain of continuous functions  $\{f_i : D \rightarrow E\}_{i \in \mathbb{N}}$

we need to prove 
$$\left( \bigsqcup_{i \in \mathbb{N}} f_i \right)^* = \bigsqcup_{i \in \mathbb{N}} f_i^*$$

take a generic  $x \in D_{\perp}$

we need to prove 
$$\left( \bigsqcup_{i \in \mathbb{N}} f_i \right)^* (x) = \left( \bigsqcup_{i \in \mathbb{N}} f_i^* \right) (x)$$

if  $x = \perp_{D_{\perp}}$  it is obvious

if  $x = \lfloor d \rfloor$  we have...

(see next slide)

**(continue)**

$$\left( \bigsqcup_{i \in \mathbb{N}} f_i \right)^* (\llbracket d \rrbracket) = \left( \bigsqcup_{i \in \mathbb{N}} f_i^* \right) (\llbracket d \rrbracket)$$

$$\left( \bigsqcup_{i \in \mathbb{N}} f_i \right)^* (\llbracket d \rrbracket) = \left( \bigsqcup_{i \in \mathbb{N}} f_i \right) (d) \quad \text{by def of lifting}$$

$$= \bigsqcup_{i \in \mathbb{N}} f_i(d) \quad \text{by lub in a functional domain}$$

$$= \bigsqcup_{i \in \mathbb{N}} f_i^*(\llbracket d \rrbracket) \quad \text{by def of lifting}$$

$$= \left( \bigsqcup_{i \in \mathbb{N}} f_i^* \right) (\llbracket d \rrbracket) \quad \text{by lub in a functional domain}$$

# Let notation (de-lifting)

$(E, \sqsubseteq_E)$  CPO $_{\perp}$       $\lambda x. e \in [D \rightarrow E]$       $t \in D_{\perp}$

$$\text{let } x \leftarrow t. e \quad \triangleq \quad \underbrace{\underbrace{\underbrace{(\lambda x. e)^*}_{[D \rightarrow E]} \underbrace{t}_{D_{\perp}}}_{[D_{\perp} \rightarrow E]}}_E = \begin{cases} \perp_E & \text{if } t = \perp_{D_{\perp}} \\ e^{[d/x]} & \text{if } t = [d] \end{cases}$$

intuitively:

if  $t$  is a lifted value  $[d]$  then we de-lift the value and  
assign it to  $x$  in  $e$

otherwise returns  $\perp_E$

# Disjoint Union



$$\mathcal{D} = (D, \sqsubseteq_D)$$

$$\mathcal{E} = (E, \sqsubseteq_E) \quad \text{CPO}_\perp \quad \Rightarrow \quad \mathcal{D} + \mathcal{E} = (D \uplus E, \sqsubseteq_{D \uplus E})$$

$$D \uplus E \triangleq \{(0, d) \mid d \in D\} \cup \{(1, e) \mid e \in E\}$$

how to order elements?

is there a bottom element?

is it a complete order?

how to define (continuous) injections?

$$\iota_D : D \rightarrow D \uplus E$$

$$\iota_E : E \rightarrow D \uplus E$$