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MPP 2025/26 (0077A, 9CFU)

Models for Programming Paradigms

Roberto Bruni Filippo Bonchi http://www.di.unipi.it/~bruni/

Ol - Introduction

English vs Italian



Classes

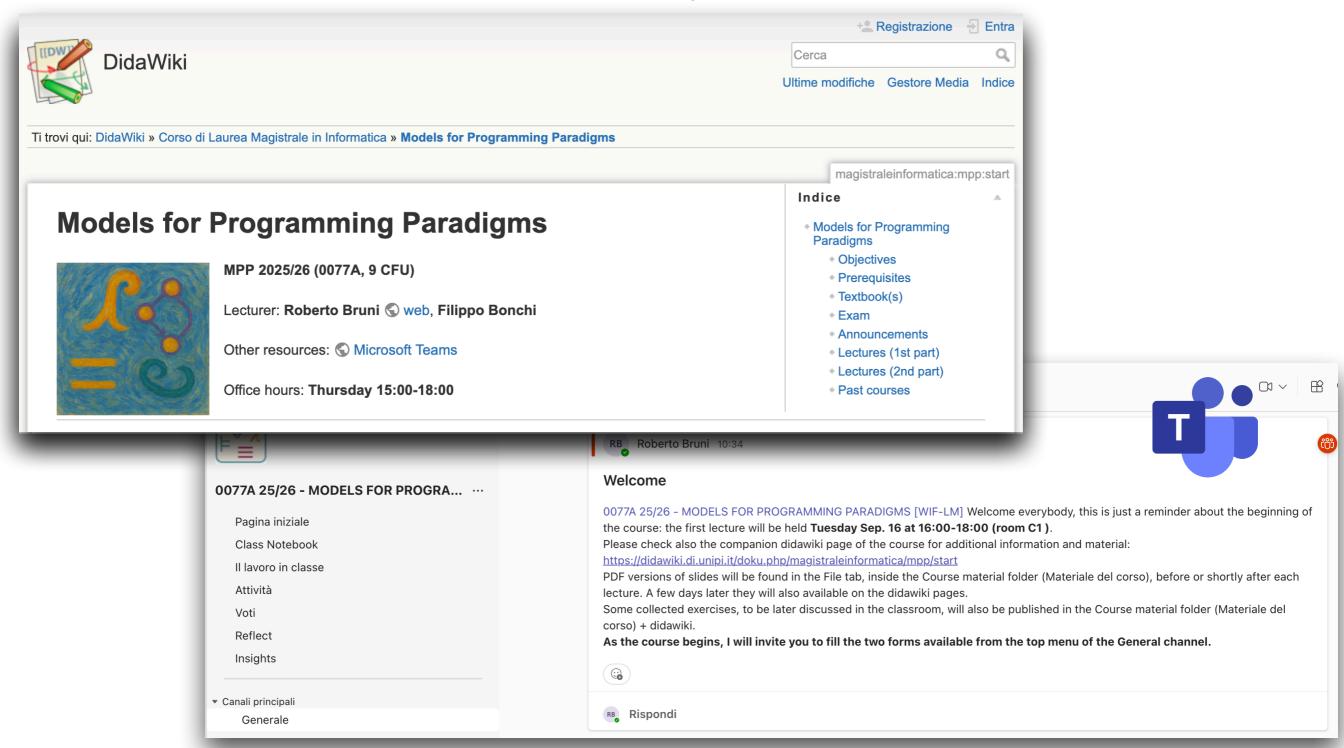
In presence, every

Tuesday: 16:00-18:00, C1

Thursday: 09:00-11:00, L1

Friday: 16:00-18:00, L1

Course material



Who am I?



http://www.di.unipi.it/~bruni





bruni@di.unipi.it

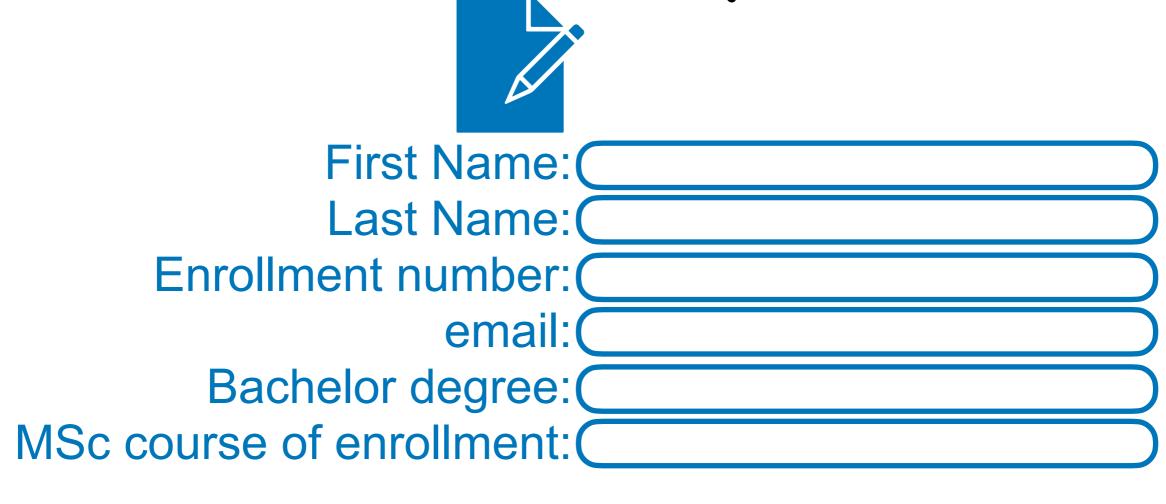
Office hours: by appointment preferably
Thursday 15:00-18:00

Research topics (theses?)

- Program logics and Abstract Interpretation
- Formal approaches to code obfuscation
- Quantum Computation and concurrency models
- Modelling and analysis of biological systems
- Graphical specification languages
- Algebraic approaches to structured graphs
- Rewrite rules for reversible languages



Who are you?



Please fill the form!



Who are you?

First Name: John

Last Name: Smith

Enrollment number: 123456

email: john.smith@email.com

Bachelor degree: Comp. Sci., Pisa, IT

MSc course of enrollment: Comp. Sci. - SW





The Course

Some quotes

Computer science is no more about computers than astronomy is about telescopes

- Edsger W. Dijkstra

Studying programming languages without formal semantics would be like studying physics without math

- from the web

All models are wrong, but some are useful

- George Box

Subjects are divided in two categories:

- 1) too difficult matters, that CANNOT be studied
- 2) easy matters, that DO NOT NEED to be studied
- back of a t-shirt

Objectives

Programming paradigms (imperative, declarative, higher order, concurrent, mobile, stochastic)

Google Go!









Mathematical frameworks (concrete & abstract)

(domains, inference rules, transition systems, λ-calculus, process algebras)



Understand

(recursion, semantics, compositionality)





Reason

(induction, modal and temporal logics, behavioural and logical equivalences) Explain (correctness, compliance, performance)



The approach

(in their simplest form, still Turing equivalent)

programming paradigms



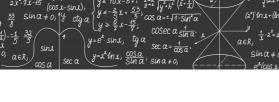
mathematical frameworks



meta-properties + proof techniques

(for all programs or just some classes of programs)





models



specifications

Key question

Given two programs *p* and *q*:

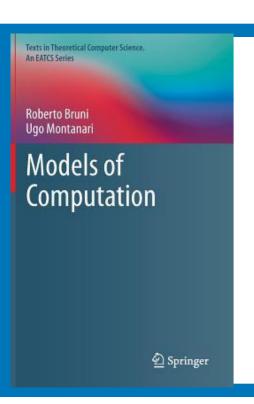
are they "equivalent"?

are they "congruent"?

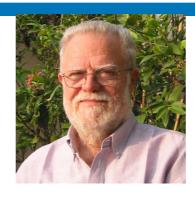
Do they behave the same?

Is it safe to replace one with the other in any context?

Textbooks

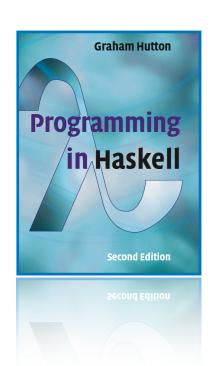


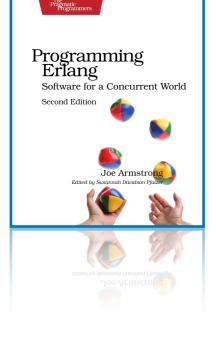


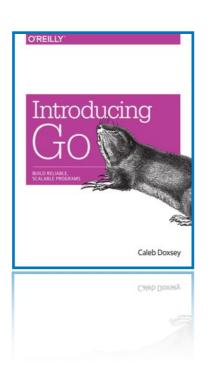


Roberto Bruni and Ugo Montanari Models of Computation Texts in Theoretical Computer Science (an EATCS series)

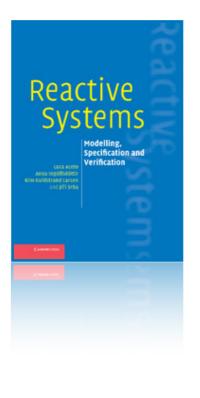
https://www.springer.com/book/9783319428987

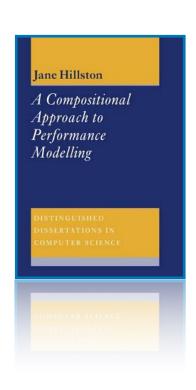










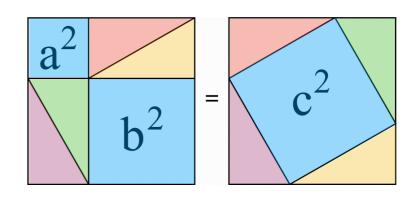


Course activities



attend classrooms: ask questions! (sleep quietly)

learn theorems: (drink many coffees)





do some thinking: solve ALL your homeworks (at least try to)

give the exam: time for a party!



Be proactive!

Let's spell out definitions together

```
% find the least (non-unitary) divisor d of n>1
% example: for n=21 the result is d=3
d := 0;
x := 2;
while (.....) do {
  if (n%x == 0) then {
     d := x;
   } else {
    x := x + 1;
```

Be proactive!

Correct me if I'm wrong

```
% find the index i of the last occurrence of n in a
% example: for n=5, a=[3,5,7,5,9] the result is i=3
i := length(a)-1;
while ( i>0 && n!=a[i] ) do {
  i := i-1;
}
```

Be proactive!

Sometimes tricky questions!

Can you find the the mistake? 1 2 3 4 5 6 7 8 9

Interaction protocol

when I will ask questions such as:

"does the program c satisfy the property p?"

there are the 3 possible answers

O yes

O no

O don't know

you are welcome to take your time... but then please pick one option whenever I ask questions in these classes

Exam

In past years, the evaluation was based on written and oral exams.

Since the covid-19 emergency, and for the current period, the evaluation will be solely based on a final oral exam.



Registration to exams is mandatory:

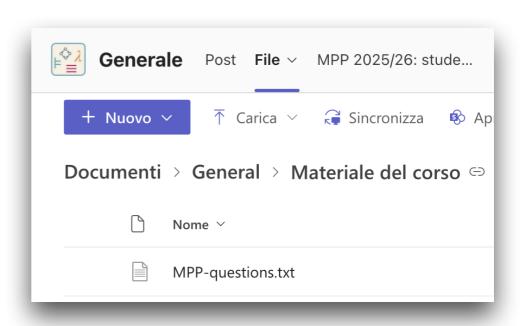
https://esami.unipi.it/esami

The exam will typically consist of:

- 1. three to four preliminary questions
- 2. one exercise (analogous to past written exams)
- 3. redoing one of the proofs seen in the course
- 4. some additional questions

The list of preliminary questions is available on Microsoft Teams, in the File tab

```
(MPP-questions.txt)
```



A sample exam

What is a complete partial order?

What are the rules of the type system of HOFL?

How is iteration achieved in CCS?

Why only positive normal forms are considered in the mu-calculus?

Consider the HOFL term

$$t \stackrel{\text{def}}{=} \mathbf{rec} \ f. \ \lambda x. \ \mathbf{if} \ x \ \mathbf{then} \ (x, \mathbf{fst}(f \ x)) \ \mathbf{else} \ (\mathbf{snd}(f \ x), x)$$

- 1. Find the principal type of t.
- 2. Find the denotational semantics of t.

Prove the Switch Lemma

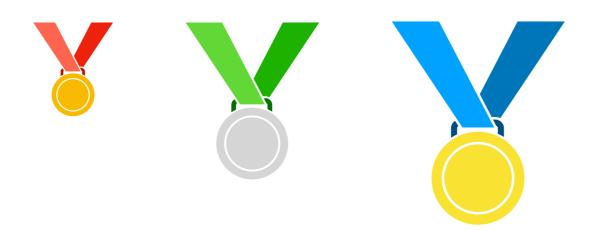
Given the initial state distribution and a DTMC, how do we compute the state distribution at time 3?

Badges?

No mid-terms
No self-evaluation tests
During the course: some "badge" exercises



Submit your solutions by email to earn bronze / silver / gold badges (no extra scores, but be proud of yourselves)



Basic set theory

$$A \cap B$$

$$A \cup B$$

$$A \setminus B$$

$$\overline{A}$$

$$a \in A$$

$$A \subset B$$

$$A \subseteq B$$

$$A \times B$$

$$a \not\in A$$

$$A \not\subseteq B$$

$$A \cap B = \emptyset$$

N

 \mathbb{R}

$$N \subset \mathbb{N}$$

$$N \subseteq \mathbb{N}$$
 $N \in \wp(\mathbb{N})$

$$S \subseteq \wp(\mathbb{N})$$

Basic set theory: functions, relations

$$f:A\to B$$

$$R \subseteq A \times B$$

functions as relations

$$R_f \stackrel{\triangle}{=} \{(a, f(a)) \mid a \in A\}$$

sets as functions (characteristic function)

$$f_N:\mathbb{N}\to\mathbb{B}$$

$$f_N(n) \stackrel{\triangle}{=} \left\{ \begin{array}{ll} 1 & n \in N \\ 0 & \text{otherwise} \end{array} \right.$$

$$N = \{ n \mid f_N(n) = 1 \}$$

First order logic

meaning of implication!

$$P \Rightarrow Q$$

$$Q \lor \neg P$$

$$\neg Q \Rightarrow \neg P$$

order of quantifiers matters!

$$\forall n \in \mathbb{N}. \ \exists m \in \mathbb{N}. \ n < m$$

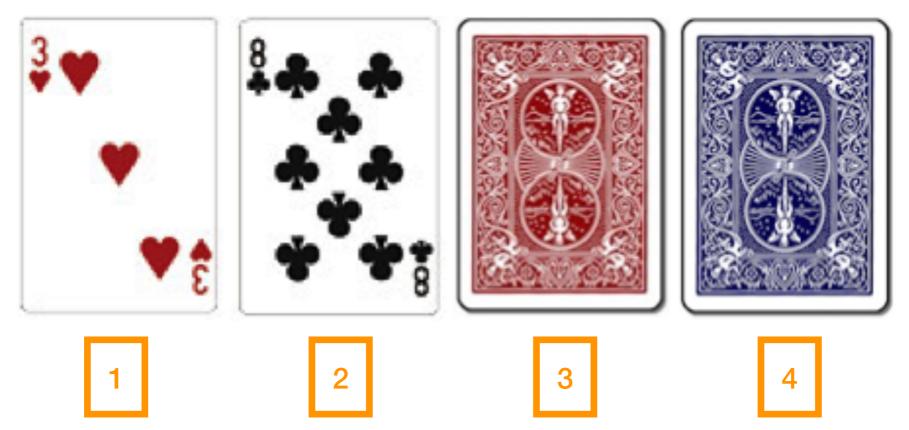
$$\exists m \in \mathbb{N}. \ \forall n \in \mathbb{N}. \ n < m$$

Implication is a tricky concept

under which circumstances is the following promise broken?

"If Rob Bery wins the election, I promise to leave the country"

(Peter Cathart) Wason Selection Task (4 cards game)



Which cards must be turned over to make sure the following statement is true?

"If the front face of a card bears an even number, then its back face is red."

(Peter Cathart) Wason Selection Task (with beers)

A policeman enters a local in Florida where a large sign reminds customers that:

"to drink beer you must be over 16 years old"

There are four customers in the local:

- 1. a boy who is drinking water,
 - a girl who is drinking beer,













- 3. an elderly lady who is drinking at a table, and
- 4. a 14-year-old teenager who is drinking at the counter.

Which customers should the policeman check to verify that the rule is respected?

Strings and context-free grammars

 $A^* \stackrel{\triangle}{=} \bigcup A^n$

$$\mathbb{B} = \{0, 1\}$$
 $\mathbb{B}^0 = \{\epsilon\}$
 $\mathbb{B}^1 = \{0, 1\}$
 $\mathbb{B}^2 = \{00, 01, 10, 11\}$
 $\mathbb{B}^3 = \{000, 001, 010, 011, 100, 101, 110, 111\}$
....
 $\mathbb{B}^* = \{\epsilon, 0, 1, 00, 01, 10, 11, 000, \ldots\}$

Alphabet A $A^n \stackrel{\triangle}{=} \underbrace{A \times \cdots \times A}$

Strings and context-free grammars

Alphabet
$$A$$

Alphabet
$$A$$
 $A^n \stackrel{\triangle}{=} \underbrace{A \times \cdots \times A}_n$

$$A^* \stackrel{\triangle}{=} \bigcup_{n \in \mathbb{N}} A^n$$

$$\mathbb{B}^* = \{\epsilon, 0, 1, 00, 01, 10, 11, 000, \ldots\}$$

$$A ::= \epsilon \mid 0 A \mid 1 B$$

$$\underbrace{A \rightarrow 0 \ A} \rightarrow 0 \ 1 \ \underbrace{B} \rightarrow 0 \ 1 \ 1 \ A \rightarrow 0 \ 1 \ 1 \ \epsilon = 0 \ 1 \ 1$$

$$\mathcal{L}(\mathsf{A}) = ?$$

$$\mathcal{L}(\mathsf{B}) = ?$$

Inductive and recursive definitions

$$0! \stackrel{\triangle}{=} 1$$
$$(n+1)! \stackrel{\triangle}{=} n! \cdot (n+1)$$

$$A^{0} \stackrel{\triangle}{=} \{\epsilon\}$$

$$A^{(n+1)} \stackrel{\triangle}{=} A \times A^{n}$$

$$f(n) \stackrel{\triangle}{=} \begin{cases} 1 & \text{if } n \leq 1\\ f(n/2) & \text{if } n > 1 \land n\%2 = 0\\ f(3n+1) & \text{otherwise} \end{cases}$$

$$f(12) = f(6) = f(3) = f(10) = f(5) = f(16) = f(8) = f(4) = f(2) = f(1) = 1$$

Conjectures vs theorems

a natural number *p* is **prime** if it cannot be written as the product of two smaller numbers

n	Is <i>n</i> prime?	$2^{n}-1$	Is $2^n - 1$ prime?
2	yes	3	yes
3	yes	7	yes
4	no: $4 = 2 \cdot 2$	15	no: $15 = 3 \cdot 5$
5	yes	31	yes
6	no: $6 = 2 \cdot 3$	63	no: $63 = 7 \cdot 9$
7	yes	127	yes
8	no: $8 = 2 \cdot 4$	255	no: $255 = 15 \cdot 17$
9	no: $9 = 3 \cdot 3$	511	no: $511 = 7 \cdot 73$
10	no: $10 = 2 \cdot 5$	1023	no: $1023 = 31 \cdot 33$

Conjectures vs theorems

if p is prime then $2^p - 1$ is prime

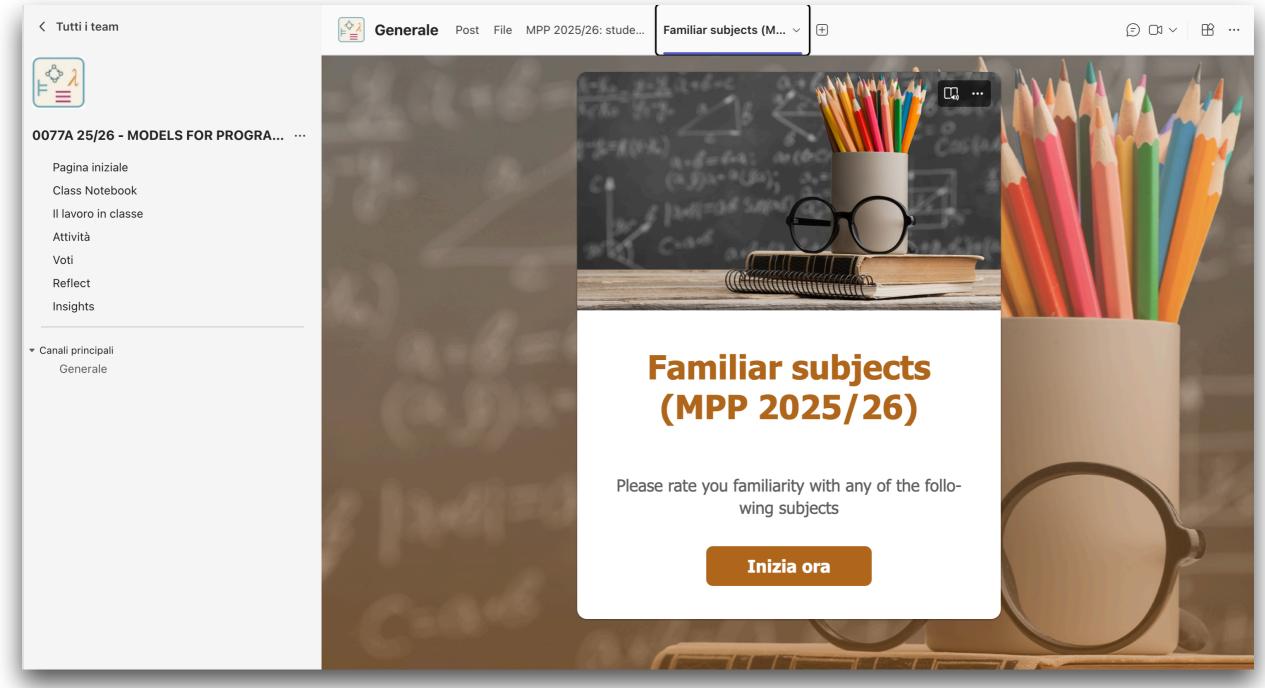
if n > 1 is not prime then $2^n - 1$ is not prime



Use any mean to prove or disprove the above conjectures

Your background?

Please fill the form about "Familiar subjects"



What's your method?



Reading textbooks



Reading lecture notes/handouts



Reviewing slides



Watching/listening to recorded lectures



Doing exercises in class



Doing exercises at home



Studying with peers / group work



Asking questions during office hours

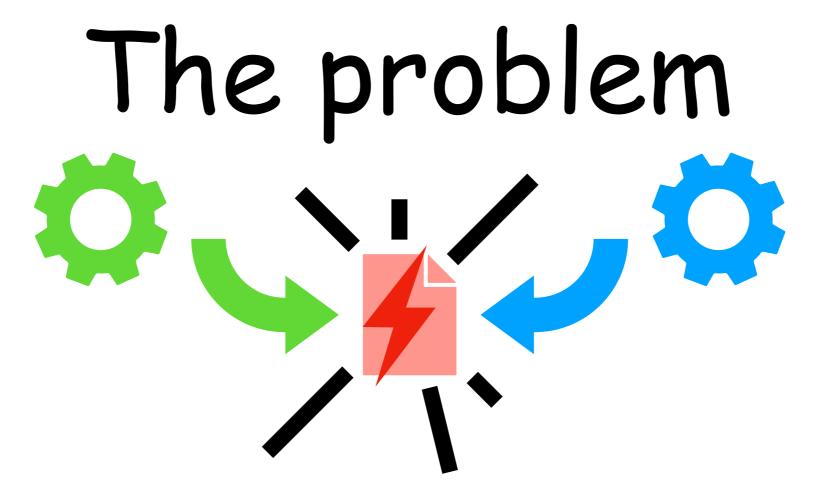


Bring laptop in class / annotate pdf



Take notes on paper

An Appetiser



Two concurrent processes share a single-use resource

They can communicate using shared memory

We want to guarantee that there are no conflicts when the processes access the resource

No strict alternation of naive turn taking is imposed

Peterson's mutual exclusion algorithm (1981)

```
% Two processes P1, P2
% Two boolean variables b1, b2 (both initially false)
% when Pi wants to enter the critical section, then it sets bi to true
% An integer variable k, taking values in {1,2}
% (initial value is arbitrary)
% the process Pk has priority over the other process
% Process P1 in pseudocode
while (true) {
                              % non critical section
                              % P1 wants to enter the critical section
    b1 := true ;
    k := 2;
                              % P1 gives priority to the other process
    while (b2 && k==2) skip; % P1 waits its turn
                              % P1 enters the critical section
    b1 := false
                              % P1 leaves the critical section
% Process P2 is analogous to P1
```

Which question?

Does Peterson's algorithm work?

What does it mean that "it works"? What do we expect?

(Progress)

If the resource is available, no process is forced to wait

(Bounded Waiting)

No process will wait forever for the resource (otherwise the easiest solution is no one gets in)

(Mutual Exclusion)

P1 and P2 are never in the critical section at the same time

Hyman's mutual exclusion algorithm (1966)

```
% Two processes H1, H2
% Two boolean variables b1, b2 (both initially false)
% when Hi wants to enter the critical section, then it sets bi to true
% An integer variable k, taking values in {1,2}
% (initial value is arbitrary)
% the process Hk has priority over the other process
용
% Process H1 in pseudocode
while (true) {
                              % non critical section
    b1 := true ;
                              % H1 wants to enter the critical section
    while (k=2) {
                             % while H2 has priority
        while (b2) skip; % H1 waits
        k := 1;
                              % H1 sets priority to itself
                              % H1 enters the critical section
    b1 := false
                              % H1 leaves the critical section
 Process H2 is analogous to H1
```

The question

Does Peterson's algorithm satisfy mutual exclusion?

Does Hyman's algorithm satisfy mutual exclusion?

For the answers be patient and wait November lectures