

Models of computation (MOD) 2014/15
Exam – July 22, 2015

[Ex. 1] Add to IMP the arithmetic expression

eval a after c

that returns the value of a in the store obtained by the execution of c (if it terminates).

1. Define the operational semantics for the new expression.
2. Show a simple counterexample to the property

$$\forall a \in Aexp. \forall \sigma \in \Sigma. \exists n \in \mathbb{N}. \langle a, \sigma \rangle \rightarrow n$$

3. Redesign the denotational semantics of IMP expressions and commands to take into account the fact that the evaluation of expressions may not terminate: illustrate all needed changes.

[Ex. 2] Let $\mathcal{D} = (D, \sqsubseteq_D)$ be a CPO_\perp . Let us consider the function $\text{min} : \mathcal{D} \times \mathcal{D} \rightarrow \mathcal{D}$ defined by letting

$$\text{min}(x, y) = \begin{cases} x & \text{if } x \sqsubseteq_D y \\ y & \text{if } y \sqsubseteq_D x \\ \perp_D & \text{if } x \not\sqsubseteq_D y \text{ and } y \not\sqsubseteq_D x \end{cases}$$

1. Exhibit a $\text{CPO}_\perp \mathcal{D}$ such that min is not monotone.¹
2. Prove that if \sqsubseteq_D is a total order then min is monotone.
3. Prove that if \sqsubseteq_D is a total order then min is continuous.

[Ex. 3] Let us consider the HOFL term

$$t \stackrel{\text{def}}{=} \text{rec } f. \lambda x. \text{if fst}(x) - \text{snd}(x) \text{ then } \text{snd}(x) \text{ else } f((\text{fst}(x) - \text{snd}(x), \text{snd}(x)))$$

1. Find the principal type of t .
2. Compute the (lazy) canonical form of $(t (6, 3))$.

[Ex. 4] Consider the CCS processes

$$\begin{aligned} p &\stackrel{\text{def}}{=} \text{rec } X. (\tau.X + \alpha.\text{nil}) \\ q &\stackrel{\text{def}}{=} \text{rec } Y. \tau.(Y + \alpha.\text{nil}) \\ r &\stackrel{\text{def}}{=} \text{rec } Z. (\tau.(Z + \alpha.\text{nil}) + \alpha.\text{nil}) \end{aligned}$$

1. Draw the labelled transition systems for p , q and r .
2. Which of the above processes are strong bisimilar?
3. Which of the above processes are weak bisimilar?

¹Assume the usual order over the product domain $\mathcal{D} \times \mathcal{D}$.