

Branch and Bound

①

(Wolsey : Chapter 7 (7.1, 7.2))

Consider

$$(P) \quad z = \max_{x \in S} c(x)$$

How can we break (P) into a series of smaller (and easier) problems, solve the smaller problems, and then put the information together to solve (P)?

Divide and conquer approach

Proposition: Let $S = S_1 \cup S_2 \cup \dots \cup S_K$ be a decomposition of S , and let

$$z^k = \max \{ c(x) : x \in S_k \}, \quad k = 1, \dots, K.$$

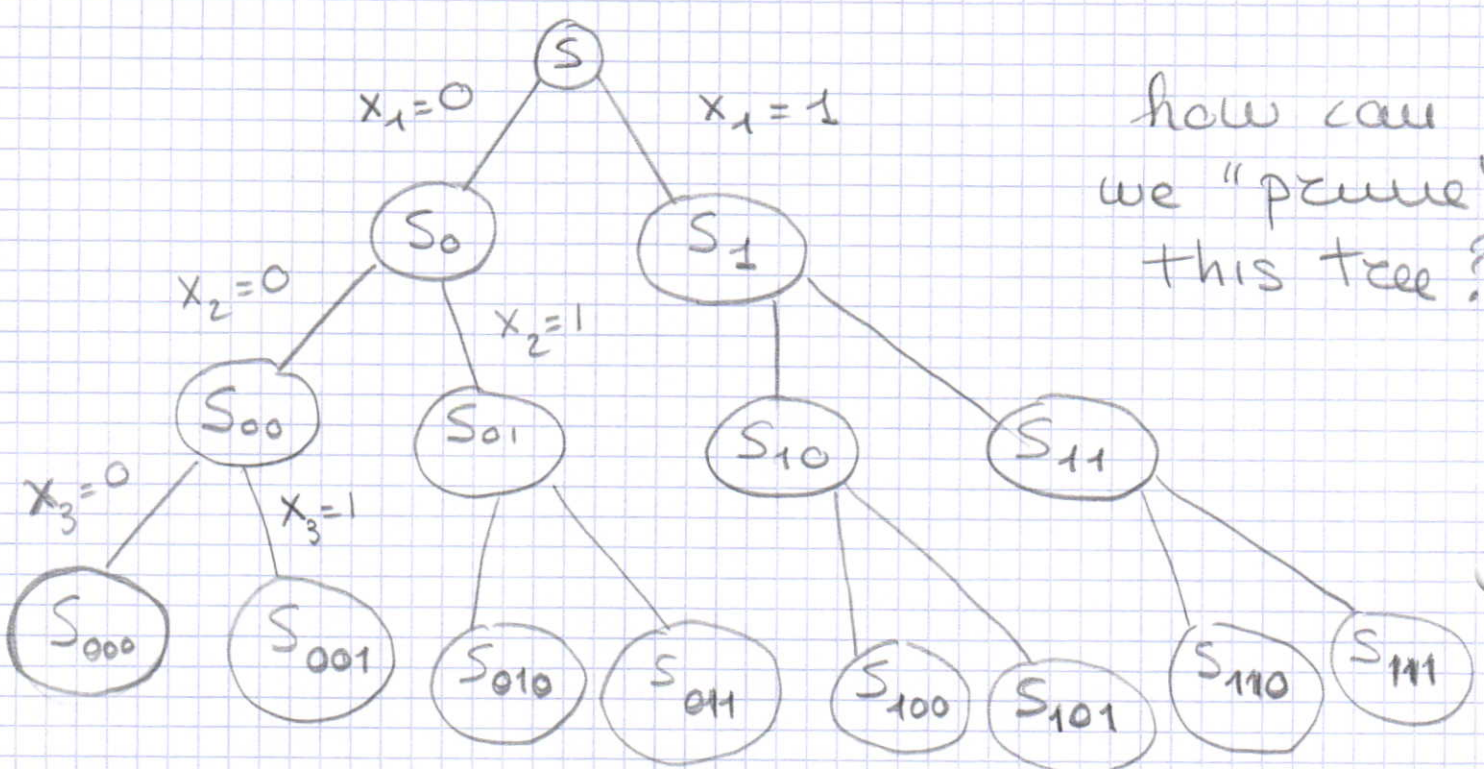
Then $z = \max_k z^k$.

- A typical way is to decompose S via an enumeration tree
- However: a complete enumeration is usually impossible (we can not divide indefinitely)

So:

- how can we use some bounds on $f(x^k)$ intelligently?
- how can we put together bound information?

example: binary enumeration tree for $S \subseteq \{0, 1\}^3$



how can we "prune" this tree?

Implicit enumeration

(3)

Proposition: Let $S = S_1 \cup \dots \cup S_K$ be a decomposition of S , and let $z^k = \max\{c(x) : x \in S_k\}$, \bar{z}^k be an upper bound on z^k , and \underline{z}^k be a lower bound on z^k , $k = 1, \dots, K$.

Then:

$$\bar{z} = \max_k \bar{z}^k \text{ is an upper bound on } z$$

and

$$\underline{z} = \max_k \underline{z}^k \text{ is a lower bound on } z$$

- So, bound information (partial information) about subproblems can be put together to derive bounds on z !
- What can be deduced from these bounds, and which sets need further examination to compute z ?