

Tecniche di Progettazione: Design Patterns

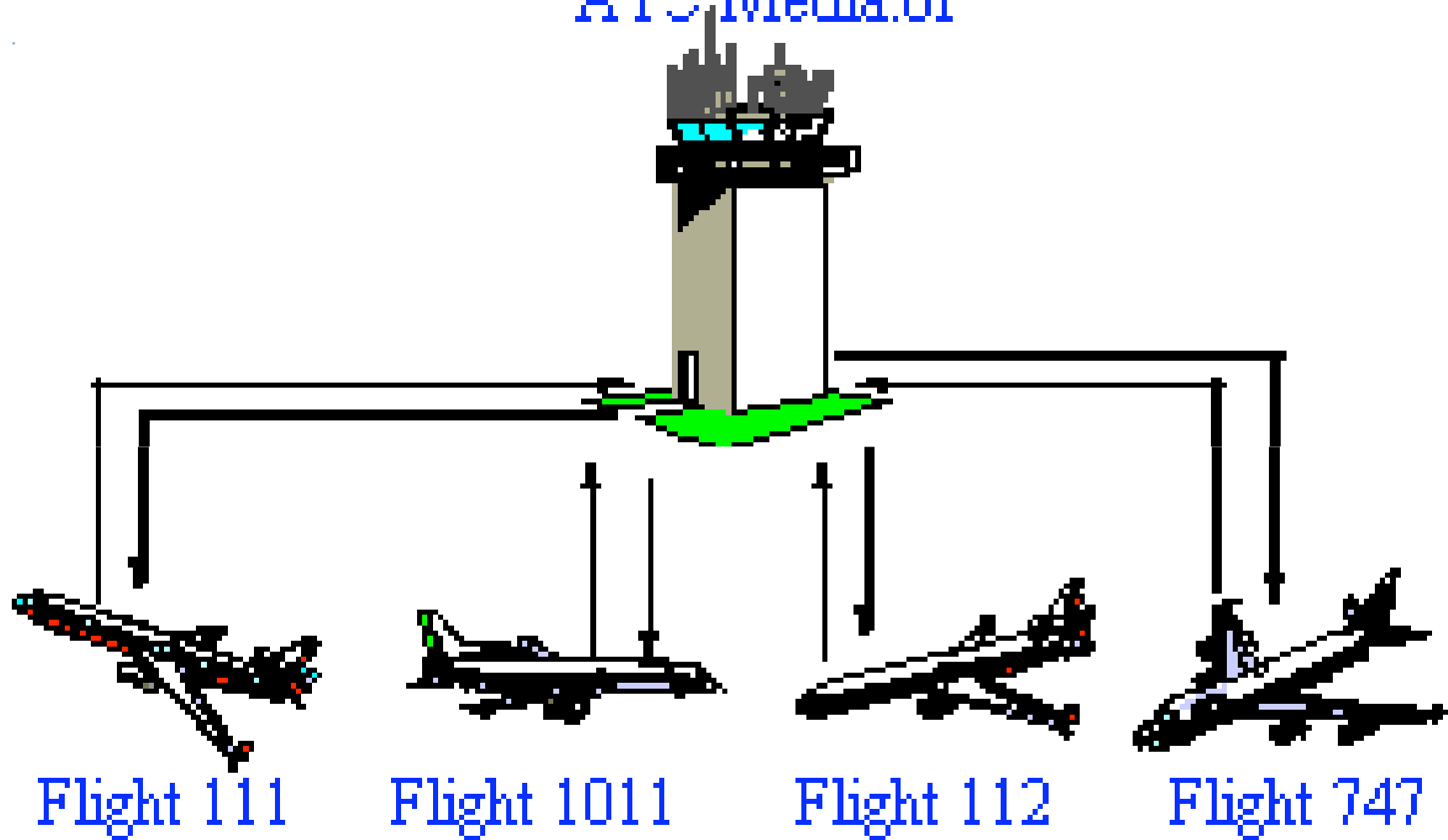
GoF: Mediator & Coordination

Applicability

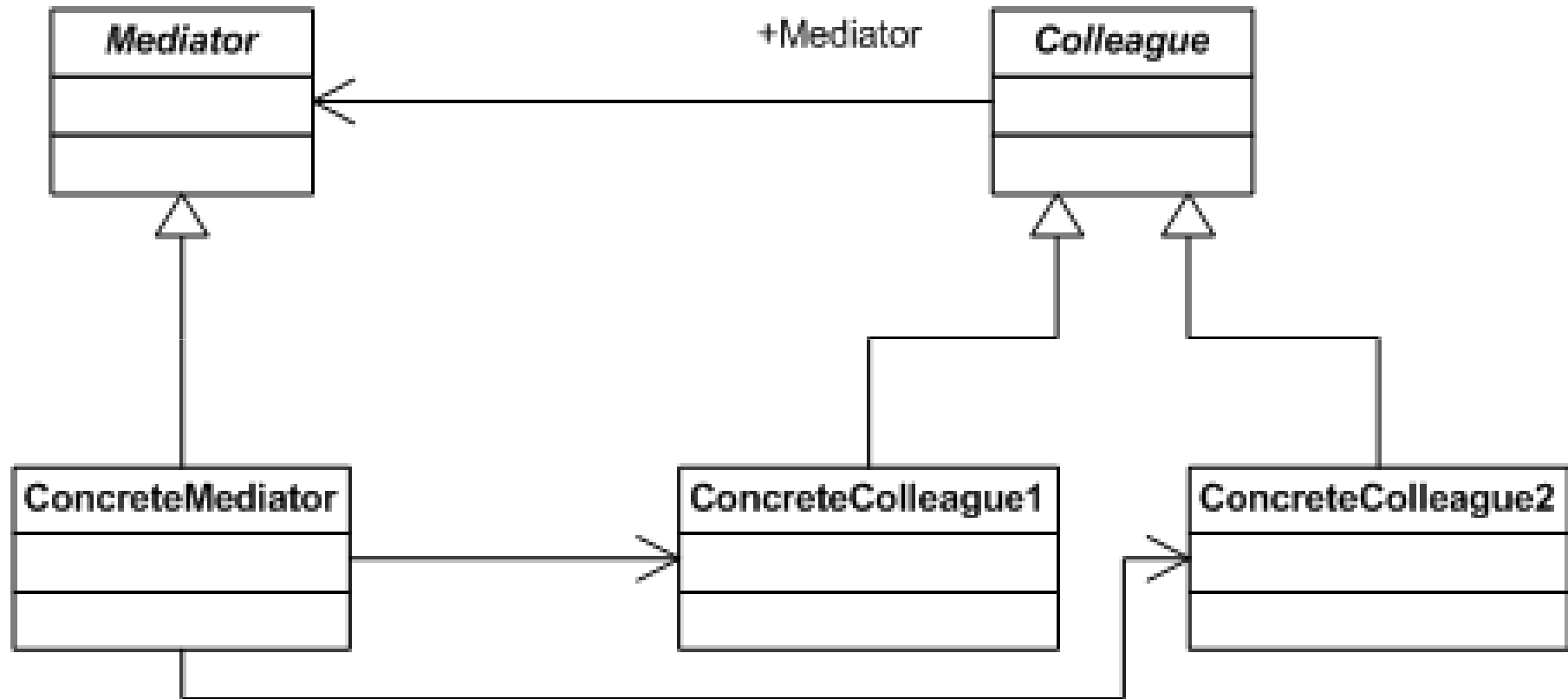
- ▶ A set of objects communicates in well-defined, but complex ways. Often with unstructured dependencies.
- ▶ Reusing objects is difficult because it refers to and communicates with many other objects.
- ▶ A behavior that's distributed between several classes should be customizable without a lot of subclassing.



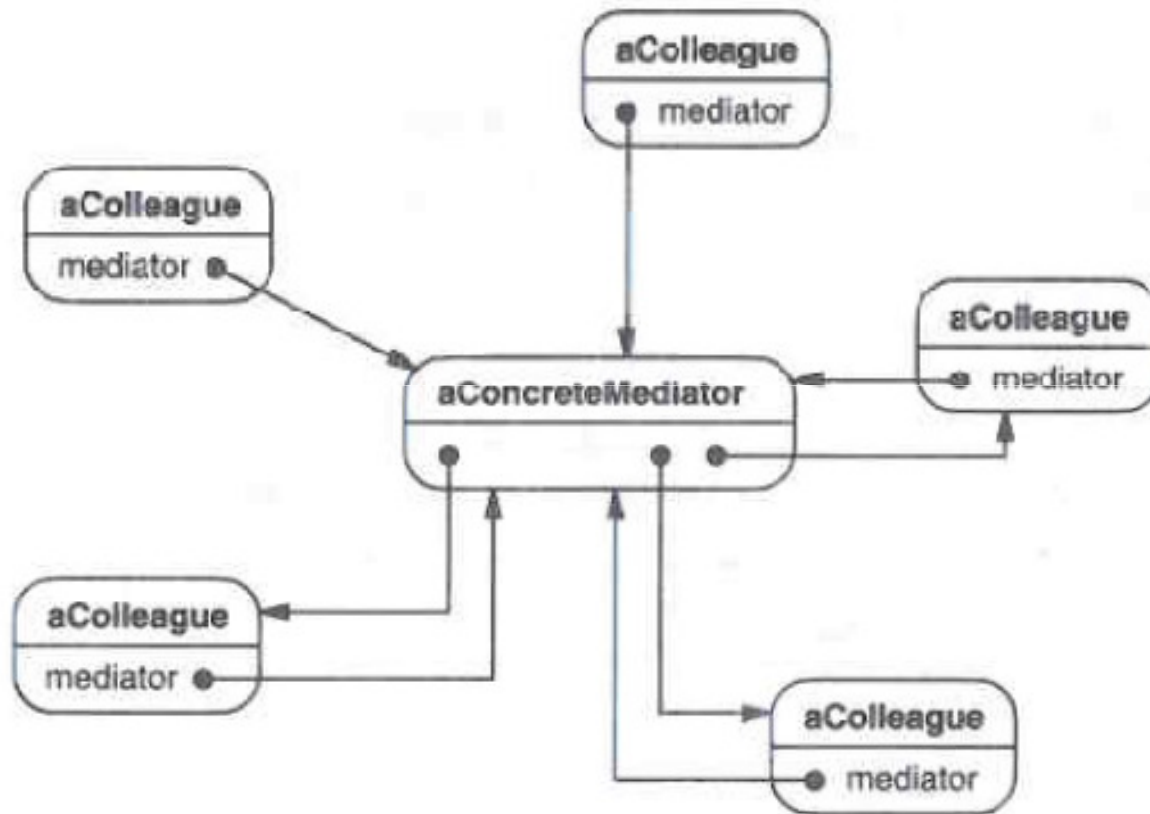
ATC Mediator



Mediator: structure



Structure



Mediator

- ▶ A mediator acts as a HUB of communication for a group of interdependent objects.
- ▶ In other words, a mediator object is the sole object that knows all other objects in a group of collaborating objects and how they should interact with each other.
 - ▶ All other objects should interact with the mediator object, instead of each other.

Consequences

- ▶ **Decouples colleagues**
 - ▶ Colleagues become more reusable.
 - ▶ You can have multiple types of interactions between colleagues, and you don't need to subclass or otherwise change the colleague class to do that.
- ▶ **Limits subclassing**
 - ▶ Localizes behavior that would be otherwise distributed among many objects
 - ▶ Changes in behavior require changing only the Mediator class



Consequences

- ▶ **Simplifies object protocols**
 - ▶ Many-to-many interactions replaced with one-to-many interactions
 - ▶ More intuitive
 - ▶ More extensible
 - ▶ Easier to maintain
- ▶ **Abstracts object cooperation**
 - ▶ Mediation becomes an object itself
 - ▶ Interaction and individual behaviors are separate concepts that are encapsulated in separate objects



Consequences

- ▶ **Centralizes control**
 - ▶ Mediator can become very complex
 - ▶ With more complex interactions, extensibility and maintenance may become more difficult
 - ▶ Using a mediator may compromise performance



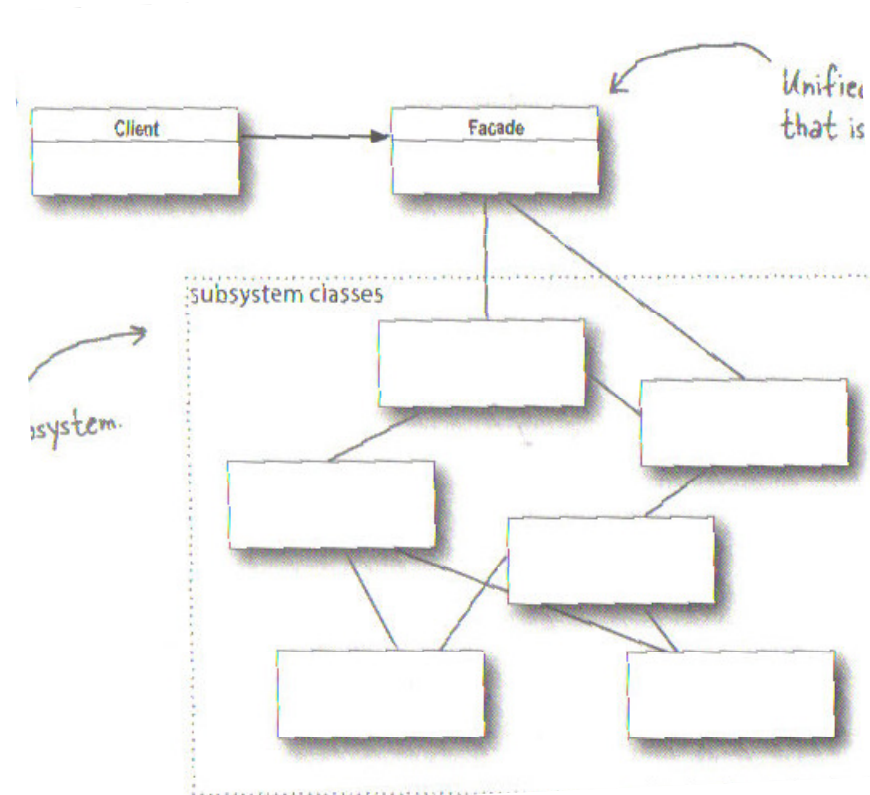
Implementation Issues

- ▶ Omitting the abstract Mediator class – possible when only one mediator exists
- ▶ Strategies for Colleague-Mediator communication
 - ▶ Observer class
 - ▶ Pointer / other identifier to “self” passed from colleague to mediator, who pass it to the other colleague(s)



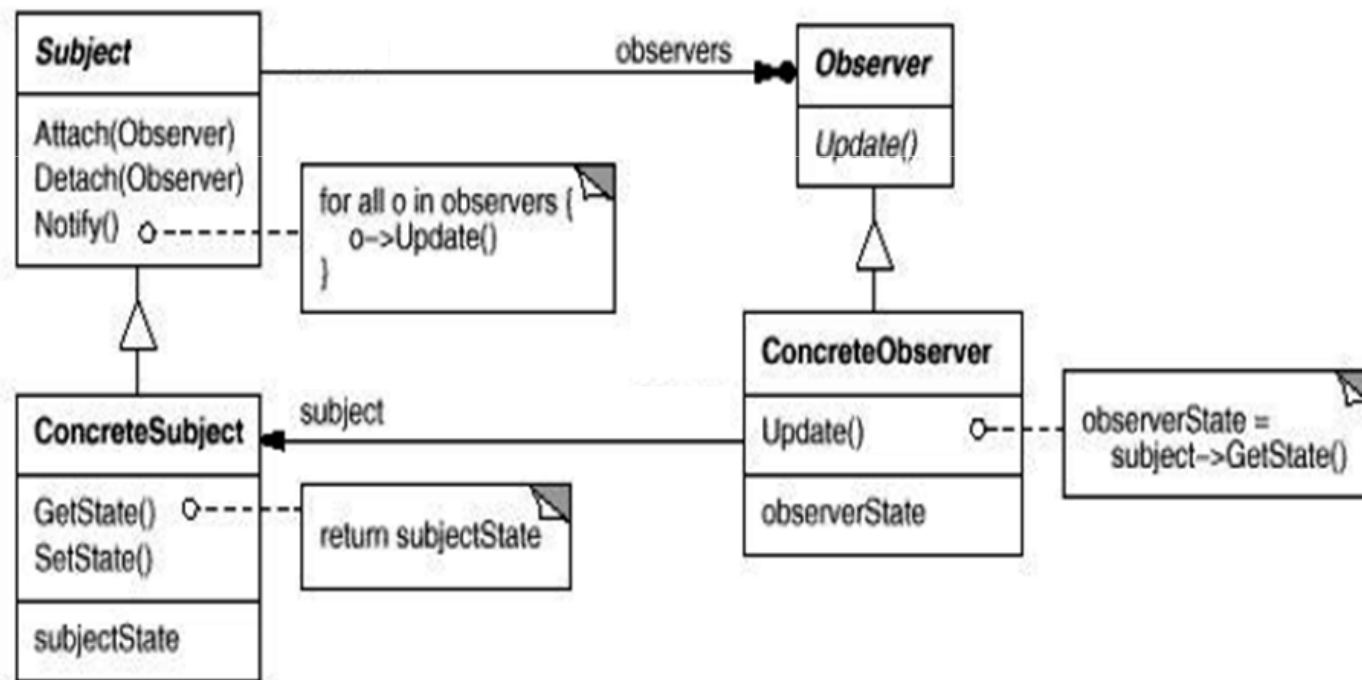
Related Patterns: Façade

- ▶ Unidirectional rather than cooperative interactions between object and subsystem
- ▶ Mediator is like a multi-way Façade pattern.
- ▶ A facade is a "unified interface" for a set of interfaces in a subsystem - for use by consumers of the subsystem - not among the components of the subsystem.



Related Patterns: Observer

- ▶ May be used as a means of communication between Colleagues and the Mediator



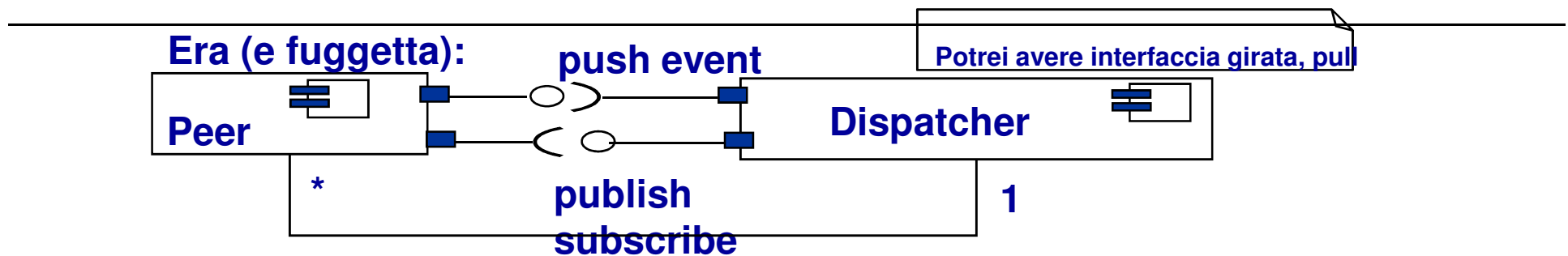
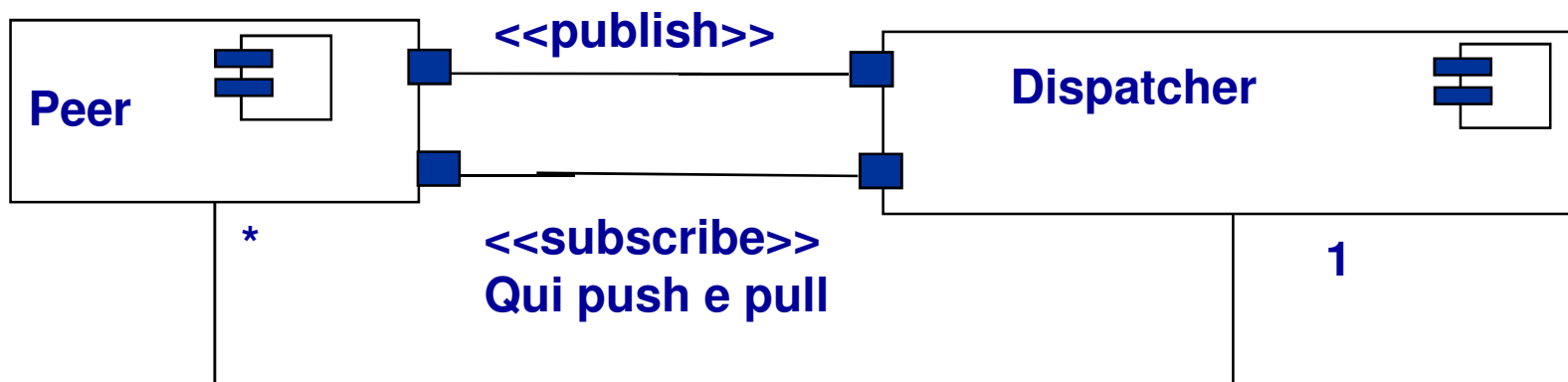
Pattern architetturale: Publish-subscribe

- ▶ Le componenti interagiscono annunciando eventi: ciascuna componente si “abbona” a classi di eventi rilevanti per il suo scopo
- ▶ Ciascuna componente, volendo, può essere sia produttore che consumatore di eventi
- ▶ Disaccoppia produttori e consumatori di eventi e favorisce le modifiche dinamiche del sistema



Pattern architetturale: Publish-subscribe

Operazioni: subscribe, unsubscribe, publish, notify (push), letMeKnow (pull)



Pattern architetturale: Publish-subscribe

- ▶ In questo stile, mittenti e destinatari di messaggi dialogano attraverso un tramite, detto dispatcher o broker.
- ▶ Il mittente di un messaggio (detto publisher) non deve essere consapevole dell'identità dei destinatari (detti subscriber); esso si limita a "pubblicare" (to publish) il proprio messaggio al dispatcher.
- ▶ I destinatari si rivolgono a loro volta al dispatcher "abbonandosi" (to subscribe) alla ricezione di messaggi.
- ▶ Il dispatcher quindi inoltra ogni messaggio inviato da un publisher a tutti i subscriber interessati a quel messaggio.



Coordination

- ▶ Coordination is managing dependencies between activities.
 - ▶ — Malone and Crowston, CACM, 26.1
- ▶ Blackboards
- ▶ Linda and tuple spaces
- ▶ BPEL (Business Process Execution Language) and web services (BPEL4WS o WS-BPEL)

Blackboard

- ▶ **Barbara Hayes-Roth**

 - A blackboard architecture for control.
Artificial Intelligence, 1985, 26, 251-321.

- ▶ **A shared data repository**

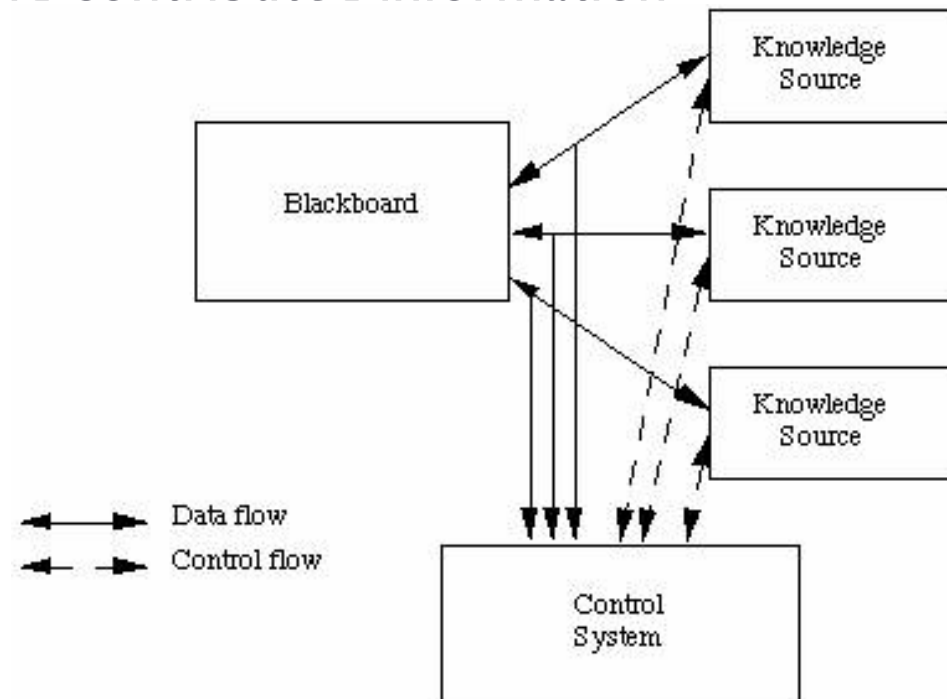
 - ▶ stores all key system data

- ▶ **Other components**

 - ▶ communicate with each other via the blackboard
 - ▶ respond to changes in the state of the blackboard

Blackboard

- ▶ **Blackboard systems consists of three major components:**
 - ▶ The software specialist modules, called knowledge sources
 - ▶ The blackboard, a shared repository of problems, partial solutions, suggestions, and contributed information
 - ▶ The control shell, which controls the flow of problem-solving activity



Blackboard

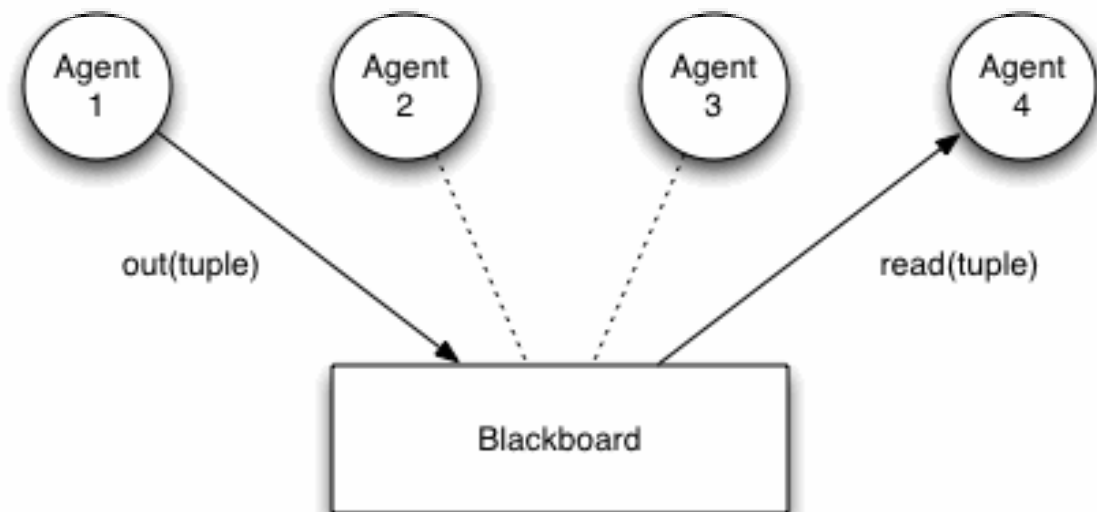
- ▶ The **blackboard** approach provides freedom from message-passing constraints. The message-passing paradigm, although modular, requires a recipient of the message as well as a sender.
- ▶ Often, the recipient is not known, or the recipient might have been deleted. In the **blackboard** approach, the “message” is placed on the **blackboard**, and the developer of the module is freed from worrying about other modules”
 - ▶ --Nii, *Blackboard Architectures and Applications*

Blackboard

- ▶ The **Blackboard** pattern uses opportunistic reasoning in solving problems for which no deterministic solution strategies are known before hand, or for problems too vast for a complete exhaustive search
- ▶ The **Blackboard** pattern provides a design in which several specialized subsystems utilize their partial knowledge bases and strategies to build an approximated solution to the original problem.
- ▶ **Blackboard** problems are often solved with solutions involving highly recursive languages such as Lisp, or backward-chaining engines such as Prolog

Form blackboards to tuple spaces

- ▶ Blackboards extended to non AI domains, e.g. as a communication media between distributed processes.
- ▶ Tuple spaces
- ▶ Tuple spaces were the theoretical underpinning of the Linda language
David Gelernter & Nicholas Carriero at Yale University.



Problem

In concurrent and distributed programming we need mechanisms for controlling interactions between the *autonomous* components

- ▶ “Come and go as they please”
- ▶ Turn on/off
- ▶ Connect/Disconnect
- ▶ Plug in/Plug out
- ▶ Fail/Crash



Components

- ▶ **Active Data Storage**
 - ▶ Database + Coordination System
 - ▶ Stores data
 - ▶ Also tracks which agents are interested in what data
 - ▶ Notifies agents when data is available
 - ▶ Pushes notifications to agents
 - ▶ Agent clients don't need to poll
- ▶ **Agent**
 - ▶ “Clients” of the tuple space
 - ▶ Different from Client/Server
 - ▶ Agents contain all intelligence in the system
 - ▶ There is no application layer server code
 - ▶ aka Knowledge Sources



Connectors

▶ Data subscription

- ▶ Agents can register interest in data
- ▶ Subscription remains in affect until data is available
 - ▶ Critical difference compared to database queries
- ▶ Provides temporal decoupling
 - ▶ Simplifying synchronization

▶ Data publication

- ▶ All communication between agents done by modifying active data storage
- ▶ Similar to traditional database updates except also trigger notification to subscribers



Example: Carnegie Mellon Uni. Robotics

- ▶ The CMU Navlab group builds robot cars, trucks, and buses, capable of autonomous driving or driver assistance
- ▶ Data Storage
 - ▶ information the Robot knows of its environment
- ▶ Agents
 - ▶ Robot parts/sensors
 - ▶ Actually software components that control robot parts
- ▶ Robot parts share data and react to changes of environment
 - ▶ ... But parts do not directly depend on each other
 - ▶ ... Some parts can fail or be destroyed without compromising the entire robot



Platforms

- ▶ Programming technology to implement blackboard-based systems
 - ▶ Linda
 - ▶ JavaSpaces
 - ▶ GigaSpaces
 - ▶ Apache River
 - ▶ Many more ...

- ▶ We'll learn about Linda
 - ▶ A specialized Blackboard programming language
 - ▶ Active storage is called a "Tuple Space"
- ▶ Nice easy to read syntax



Linda

- ▶ Linda è un linguaggio di coordinamento che estende i linguaggi tradizionali permettendone l'utilizzo nello sviluppo di applicazioni in ambiente distribuito.
- ▶ Linda è indipendente dall'architettura sottostante (a memoria condivisa o rete di calcolatori) e dal linguaggio sequenziale usato (C, C++, etc ...).
- ▶ Linda implementa una memoria associativa logicamente condivisa da tutti i processi dell'applicazione , detta Spazio delle Tuple.
 - ▶ Si parla di memoria associativa perché le tuple vengono identificate tramite matching su una chiave piuttosto che utilizzare un meccanismo tradizionale di indirizzamento.

Le tuple

- ▶ Una tupla è una sequenza di campi con tipo
(" linda ", 2, 3.14)
(sqrt(3.14), #FF3D0B,'t')
(0,'a')
- ▶ Ogni campo può contenere un valore di un qualsiasi tipo semplice o strutturato) ammesso dal linguaggio sequenziale utilizzato.

Operazioni sullo spazio delle tuple

- ▶ `out(t) eval(t)` Generano nuove tuple
- ▶ `in(s)` Elimina tuple dallo spazio delle tuple
- ▶ `rd(s)` Legge tuple dallo spazio delle tuple

- ▶ Il parametro di `out()` e di `eval ()` è una tupla `t`
- ▶ Il parametro di `in()` e di `rd()` è un' anti tupla (o template).
- ▶ Un'anti tupla è una sequenza di campi tipati che possono essere o dei campi valore (parametro attuale) o un campo indefinito (parametro formale) introdotto da “?”
- ▶ `(?radix, ``a string", ?j, 100)`

Regole di Matching

- ▶ Si ha un matching tra una tupla t ed un'anti tupla s se:
- ▶ il numero di campi di t e di s è lo stesso
- ▶ si ha una corrispondenza tra i tipi di ogni campo
- ▶ i campi valore di t sono uguali ai campi valore di s

(?radix, ``a string", ?j)

(sqrt(2), ``a string", 100)

Tuple Space Operations

- ▶ Subscription Operators

- ▶ $\text{in}(t)$

- ▶ find and remove a “matching” tuple from the tuple space; block until a matching tuple is found

- ▶ $\text{rd}(t)$

- ▶ like $\text{in}(t)$ except that the tuple is not removed



Tuple Space Operations

- ▶ Publication Operators

- ▶ out (t)

- ▶ insert a tuple t into the Tuple Space

- ▶ eval(t)

- ▶ send a live tuple into the tuple space which gets evaluated into a 'data' tuple
- ▶ spawn/fork Agents in the Tuple Space
- ▶ results of calls are collected and inserted as tuple

`eval("f_calc ", i, f(i, st))`



Example

- ▶ N-element vector stored as n tuples in the Tuple space:

(“V”, 1, element1)

(“V”, 2, element2)

...

(“V”, n, elementn)

- ▶ To read the j^{th} element and assign it to x, use

rd (“V”, j, ?x)

- ▶ To change the ith element, use

in (“V”, i, ?OldVal)

out(“V”, i, NewVal);



Example: Agent Synchronization

- ▶ Each Agent within some group must wait at a barrier until all Agents in the group have reached the barrier; then all can proceed.

- ▶ Set up barrier:

```
out("barrier", n);
```

- ▶ Each Agent does the following:

```
in("barrier", ?val);
```

```
out("barrier", val-1);
```

```
rd("barrier", 0);
```



Semafori

- ▶ Operazioni sui Semafori
- ▶ out(“semaphore”) inizializzazione del semaforo
- ▶ in(“semaphore”) acquisizione del semaforo
- ▶ out(“semaphore”) rilascio del semaforo

- ▶ es . processo che usa un semaforo
 - in(“semaphore”)
 - sezione critica
 - out(“semaphore”)

Dining Philosopher's Problem

- ▶ N agents compete for use of N data resources (tuples)
- ▶ Agents are labeled 0 to N-1
- ▶ Agents require 2 resources in order to “eat()”
 - ▶ Agent 0 shares a resource with Agent 1 and Agent N
 - ▶ for $i \neq 0$
 - ▶ Agent i needs 1 resource that is needed by Agent $i + 1 \% N$
 - ▶ Agent i needs 1 resource that is needed by Agent $i - 1$
- ▶ Agents release resources after “eat()”



Dining Philosophers Problem

```
initialize()
{
    int i;
    String[] d = {.....};
    for (i=0; i < num; i++) {
        out ("chopstick", i, d[i]);
    }
    for (i=0; i < num; i++) {
        eval ( phil(i) );
    }
}
```

```
void phil(int i) {
    while (true) {
        String d1, d2;
        in("chopstick", i, ?d1);
        in("chopstick", (i + 1) % n, ?d2);
        eat (d1, d2);
        out("chopstick", i, d1);
        out("chopstick", (i + 1) % n, d2);
    }
}
```



Tuple Spaces provide:

- ▶ Nondeterminism
- ▶ Structured naming similar to “select” in relational DB, and pattern matching in AI.
- ▶ Time uncoupling (Communication between time-disjoint processes)
- ▶ Distributed sharing (Variables shared between disjoint processes)



Implementazione

- ▶ Richiesta efficienza nel pattern matching all'interno dello spazio delle tuple
- ▶ Preprocessing
 - ▶ A: out("x", i); B: in("x", ?i)
 - ▶ x è costante durante la compilazione, quindi si mettono A e B in contatto diretto (variabile condivisa o scambio di messaggi).
 - ▶ Problemi: con compilazioni separate o linguaggi interpretati.
- ▶ Restringere la ricerca su una chiave singola
 - ▶ Se le anti tuple hanno una sola variabile al loro interno, è possibile adottare l'hashing per la memorizzazione ed il recupero delle tuple , ottenendo l'accesso ad una tupla in tempo costante

Spazi delle tuple multipli

- ▶ lo spazio delle tuple non è centralizzato , la quantità di tuple da controllare diminuisce aumentando l'efficienza.