

<http://didawiki.di.unipi.it/doku.php/magistraleinformatica/psc/>

PSC 2020/21 (375AA, 9CFU)

Principles for Software Composition

Roberto Bruni

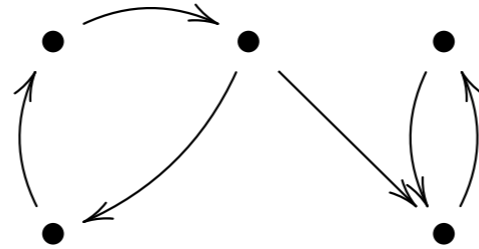
<http://www.di.unipi.it/~bruni/>

23 - Mu-calculus

μ -calculus

mu-calculus

models



syntax

ψ	$::=$	tt ff $\psi_0 \wedge \psi_1$ $\psi_0 \vee \psi_1$	classical ops
		p $\neg p$	atomic propositions
		$\diamond\psi$	there is a next state where ψ holds
		$\square\psi$	ψ holds at every next state
		x	predicate variable, for recursive def
		$\mu x. \psi$	LEAST FIXPOINT of $x =_{\min} \psi$
		$\nu x. \psi$	GREATEST FIXPOINT of $x =_{\max} \psi$

mu-calculus: semantics

$$G = (V, \rightarrow)$$

$[[\psi]]_\rho$ set of nodes where ψ holds


$$\rho : P \cup X \rightarrow \wp(V)$$

assignment

$(\wp(V), \subseteq)$ is a complete lattice

for monotone functions: least / greatest fixpoint exist
for this reason negation is not present in the syntax
the formulas we consider: *positive normal form*

otherwise:

even number of negations before each variable occurrence

Positive normal form

$$\neg \diamond \psi \equiv \square \neg \psi$$

$$\neg \square \psi \equiv \diamond \neg \psi$$

$$\neg \mu x. \psi \equiv \nu x. \neg \psi[\neg x / x]$$

$$\neg \nu x. \psi \equiv \mu x. \neg \psi[\neg x / x]$$

$$\neg \neg \psi \equiv \psi$$

mu-calculus: semantics

$$\llbracket \cdot \rrbracket : \mathcal{F} \rightarrow (P \cup X \rightarrow \wp(V)) \rightarrow \wp(V)$$

defined by structural induction

$$\llbracket \mathbf{tt} \rrbracket \rho \triangleq V$$

$$\llbracket \psi_0 \wedge \psi_1 \rrbracket \rho \triangleq \llbracket \psi_0 \rrbracket \rho \cap \llbracket \psi_1 \rrbracket \rho$$

$$\llbracket \mathbf{ff} \rrbracket \rho \triangleq \emptyset$$

$$\llbracket \psi_0 \vee \psi_1 \rrbracket \rho \triangleq \llbracket \psi_0 \rrbracket \rho \cup \llbracket \psi_1 \rrbracket \rho$$

$$\llbracket p \rrbracket \rho \triangleq \rho(p)$$

$$\llbracket \diamond \psi \rrbracket \rho \triangleq \{v \mid \exists w \in \llbracket \psi \rrbracket \rho. v \rightarrow w\}$$

$$\llbracket \neg p \rrbracket \rho \triangleq V \setminus \rho(p)$$

$$\llbracket \square \psi \rrbracket \rho \triangleq \{v \mid \forall w. v \rightarrow w \Rightarrow w \in \llbracket \psi \rrbracket \rho\}$$

$$\llbracket x \rrbracket \rho \triangleq \rho(x)$$

$$\llbracket \mu x. \psi \rrbracket \rho \triangleq \mathit{fix} \ \lambda S. \llbracket \psi \rrbracket \rho [S/x]$$

$$\llbracket \nu x. \psi \rrbracket \rho \triangleq \mathit{FIX} \ \lambda S. \llbracket \psi \rrbracket \rho [S/x]$$

mu-calculus: fixpoint

V finite

$f : \wp(V) \rightarrow \wp(V)$ monotone (hence continuous)

we can compute the least fixpoint by

$$\text{fix } f = \bigcup_{n \in \mathbb{N}} f^n(\emptyset)$$

we can compute the greatest fixpoint by

$$\text{FIX } f = \bigcap_{n \in \mathbb{N}} f^n(V)$$

Examples

$$\begin{aligned} \llbracket \diamond \mathbf{tt} \rrbracket \rho &\triangleq \{v \mid \exists w \in \llbracket \mathbf{tt} \rrbracket \rho. v \rightarrow w\} \\ &= \{v \mid \exists w \in V. v \rightarrow w\} \end{aligned}$$

$\diamond \mathbf{tt}$
non deadlocked states

$$\begin{aligned} \llbracket \square \mathbf{ff} \rrbracket \rho &\triangleq \{v \mid \forall w. v \rightarrow w \Rightarrow w \in \llbracket \mathbf{ff} \rrbracket \rho\} \\ &= \{v \mid \forall w. v \rightarrow w \Rightarrow w \in \emptyset\} \\ &= \{v \mid v \nrightarrow\} \end{aligned}$$

$\square \mathbf{ff}$
deadlocks

Examples

$$\begin{aligned} \llbracket \diamond \mathbf{ff} \rrbracket \rho &\triangleq \{v \mid \exists w \in \llbracket \mathbf{ff} \rrbracket \rho. v \rightarrow w\} \\ &= \{v \mid \exists w \in \emptyset. v \rightarrow w\} \\ &= \emptyset \end{aligned}$$

$\diamond \mathbf{ff}$
false

$$\begin{aligned} \llbracket \square \mathbf{tt} \rrbracket \rho &\triangleq \{v \mid \forall w. v \rightarrow w \Rightarrow w \in \llbracket \mathbf{tt} \rrbracket \rho\} \\ &= \{v \mid \forall w. v \rightarrow w \Rightarrow w \in V\} \\ &= V \end{aligned}$$

$\square \mathbf{tt}$
true

Examples

$$\begin{aligned} \llbracket \mu x. x \rrbracket \rho &\triangleq \text{fix } \lambda S. \llbracket x \rrbracket \rho [S / x] \\ &= \text{fix } \lambda S. S \\ &= \emptyset \end{aligned}$$

$\mu x. x$

false

$$\begin{aligned} \llbracket \nu x. x \rrbracket \rho &\triangleq \text{FIX } \lambda S. \llbracket x \rrbracket \rho [S / x] \\ &= \text{FIX } \lambda S. S \\ &= V \end{aligned}$$

$\nu x. x$

true

Examples

$$\begin{aligned} \llbracket \mu x. \diamond x \rrbracket \rho &\triangleq \text{fix } \lambda S. \llbracket \diamond x \rrbracket \rho[S/x] \\ &= \text{fix } \lambda S. \{v \mid \exists w \in \llbracket x \rrbracket \rho[S/x]. v \rightarrow w\} \\ &= \text{fix } \lambda S. \{v \mid \exists w \in S. v \rightarrow w\} \end{aligned}$$

$$S_0 = \emptyset$$

$$S_1 = \{v \mid \exists w \in S_0. v \rightarrow w\}$$

$$= \{v \mid \exists w \in \emptyset. v \rightarrow w\}$$

$$= \emptyset$$

$\mu x. \diamond x$

false

$\nu x. \square x$

true

Examples

how to represent?

EO ψ $\diamond\psi$

AO ψ $\square\psi$

EF p

$p \vee \dots$

$p \vee \diamond(p \vee \dots)$

$p \vee \diamond(p \vee \diamond(p \vee \dots))$

$x = p \vee \diamond x$

$\mu x. p \vee \diamond x$?

$\nu x. p \vee \diamond x$

EF p ?

Example

$$\begin{aligned} \llbracket \nu x. p \vee \diamond x \rrbracket \rho &\triangleq \text{FIX } \lambda S. \llbracket p \vee \diamond x \rrbracket \rho[S/x] \\ &= \text{FIX } \lambda S. \llbracket p \rrbracket \rho[S/x] \cup \llbracket \diamond x \rrbracket \rho[S/x] \\ &= \text{FIX } \lambda S. \rho(p) \cup \{v \mid \exists w \in \llbracket x \rrbracket \rho[S/x]. v \rightarrow w\} \\ &= \text{FIX } \lambda S. \rho(p) \cup \{v \mid \exists w \in S. v \rightarrow w\} \end{aligned}$$

$$S_0 = V$$

$$\begin{aligned} S_1 &= \rho(p) \cup \{v \mid \exists w \in S_0. v \rightarrow w\} \\ &= \rho(p) \cup \{v \mid \exists w \in V. v \rightarrow w\} \\ &= \rho(p) \cup \{v \text{ can move}\} \end{aligned}$$

EF p ?

Example

$$\begin{aligned} \llbracket \nu x. p \vee \diamond x \rrbracket \rho &\triangleq \text{FIX } \lambda S. \llbracket p \vee \diamond x \rrbracket \rho[S/x] \\ &= \text{FIX } \lambda S. \llbracket p \rrbracket \rho[S/x] \cup \llbracket \diamond x \rrbracket \rho[S/x] \\ &= \text{FIX } \lambda S. \rho(p) \cup \{v \mid \exists w \in \llbracket x \rrbracket \rho[S/x]. v \rightarrow w\} \\ &= \text{FIX } \lambda S. \rho(p) \cup \{v \mid \exists w \in S. v \rightarrow w\} \end{aligned}$$

$$S_0 = V$$

$$S_1 = \rho(p) \cup \{v \text{ can move}\}$$

$$S_2 = \rho(p) \cup \{v \mid \exists w \in S_1. v \rightarrow w\}$$

$$= \rho(p) \cup \{v \mid \exists w \in \rho(p). v \rightarrow w\} \cup \{v \text{ can make 2 moves}\}$$

EF p ?

Examples

$$\begin{aligned} \llbracket \nu x. p \vee \diamond x \rrbracket \rho &\triangleq \text{FIX } \lambda S. \llbracket p \vee \diamond x \rrbracket \rho[S/x] \\ &= \text{FIX } \lambda S. \llbracket p \rrbracket \rho[S/x] \cup \llbracket \diamond x \rrbracket \rho[S/x] \\ &= \text{FIX } \lambda S. \rho(p) \cup \{v \mid \exists w \in \llbracket x \rrbracket \rho[S/x]. v \rightarrow w\} \\ &= \text{FIX } \lambda S. \rho(p) \cup \{v \mid \exists w \in S. v \rightarrow w\} \end{aligned}$$

$$S_0 = V$$

$$S_1 = \rho(p) \cup \{v \text{ can move}\}$$

$$S_2 = \rho(p) \cup \{v \mid \exists w \in \rho(p). v \rightarrow w\} \cup \{v \text{ can make 2 moves}\}$$

$$\begin{aligned} S_n &= \{v \text{ can reach a state in } \rho(p) \text{ in less than } n \text{ moves}\} \\ &\cup \{v \text{ can make } n \text{ moves}\} \end{aligned}$$

EF p ?

Example

$$\begin{aligned} \llbracket \nu x. p \vee \diamond x \rrbracket \rho &\triangleq \text{FIX } \lambda S. \llbracket p \vee \diamond x \rrbracket \rho[S/x] \\ &= \text{FIX } \lambda S. \llbracket p \rrbracket \rho[S/x] \cup \llbracket \diamond x \rrbracket \rho[S/x] \\ &= \text{FIX } \lambda S. \rho(p) \cup \{v \mid \exists w \in \llbracket x \rrbracket \rho[S/x]. v \rightarrow w\} \\ &= \text{FIX } \lambda S. \rho(p) \cup \{v \mid \exists w \in S. v \rightarrow w\} \end{aligned}$$

$$S_n = \{v \text{ can reach a state in } \rho(p) \text{ in less than } n \text{ moves}\} \\ \cup \{v \text{ can make } n \text{ moves}\}$$

$$\bigcap_{n \in \mathbb{N}} S_n = \{v \text{ can reach a state in } \rho(p) \text{ or has an infinite path}\}$$

EF p ?

Example

$$\begin{aligned} \llbracket \mu x. p \vee \diamond x \rrbracket \rho &\triangleq \text{fix } \lambda S. \llbracket p \vee \diamond x \rrbracket \rho [S/x] \\ &= \text{fix } \lambda S. \rho(p) \cup \{v \mid \exists w \in S. v \rightarrow w\} \end{aligned}$$

$$S_0 = \emptyset$$

$$\begin{aligned} S_1 &= \rho(p) \cup \{v \mid \exists w \in S_0. v \rightarrow w\} \\ &= \rho(p) \cup \{v \mid \exists w \in \emptyset. v \rightarrow w\} \\ &= \rho(p) \end{aligned}$$

EF p ?

Example

$$\begin{aligned} \llbracket \mu x. p \vee \diamond x \rrbracket \rho &\triangleq \text{fix } \lambda S. \llbracket p \vee \diamond x \rrbracket \rho[S/x] \\ &= \text{fix } \lambda S. \rho(p) \cup \{v \mid \exists w \in S. v \rightarrow w\} \end{aligned}$$

$$S_0 = \emptyset$$

$$S_1 = \rho(p)$$

$$S_2 = \rho(p) \cup \{v \mid \exists w \in S_1. v \rightarrow w\}$$

$$= \rho(p) \cup \{v \mid \exists w \in \rho(p). v \rightarrow w\}$$

$$= \{v \text{ can reach a state in } \rho(p) \text{ in less than 2 moves}\}$$

EF p ?

Example

$$\begin{aligned} \llbracket \mu x. p \vee \diamond x \rrbracket \rho &\triangleq \text{fix } \lambda S. \llbracket p \vee \diamond x \rrbracket \rho[S/x] \\ &= \text{fix } \lambda S. \rho(p) \cup \{v \mid \exists w \in S. v \rightarrow w\} \end{aligned}$$

$$S_0 = \emptyset$$

$$S_1 = \rho(p)$$

$$S_2 = \{v \text{ can reach a state in } \rho(p) \text{ in less than 2 moves}\}$$

$$S_n = \{v \text{ can reach a state in } \rho(p) \text{ in less than } n \text{ moves}\}$$

$$\bigcup_{n \in \mathbb{N}} S_n = \{v \text{ can reach a state in } \rho(p)\}$$

$\mu x. p \vee \diamond x$

EF p

Example

which formula for
“*some deadlock is reachable*”?

$\Box \mathbf{ff}$
deadlocks

$$\mu x. \Box \mathbf{ff} \vee \Diamond x$$

$\mu x. p \vee \Diamond x$
EF p

which formula for
“*deadlock free*”?

$$\begin{aligned} \neg(\mu x. \Box \mathbf{ff} \vee \Diamond x) &= \nu x. \neg(\Box \mathbf{ff} \vee \Diamond \neg x) \\ &= \nu x. \neg(\Box \mathbf{ff}) \wedge \neg(\Diamond \neg x) \\ &= \nu x. \Diamond \mathbf{tt} \wedge \Box x \end{aligned}$$

Example

$$\begin{aligned} \llbracket \nu x. p \wedge \Box x \rrbracket \rho &\triangleq \text{FIX } \lambda S. \llbracket p \wedge \Box x \rrbracket \rho[S/x] \\ &= \text{FIX } \lambda S. \llbracket p \rrbracket \rho[S/x] \cap \llbracket \Box x \rrbracket \rho[S/x] \\ &= \text{FIX } \lambda S. \rho(p) \cap \{v \mid \forall w. v \rightarrow w \Rightarrow w \in \llbracket x \rrbracket \rho[S/x]\} \\ &= \text{FIX } \lambda S. \rho(p) \cap \{v \mid \forall w. v \rightarrow w \Rightarrow w \in S\} \end{aligned}$$

$$S_0 = V$$

$$\begin{aligned} S_1 &= \rho(p) \cap \{v \mid \forall w. v \rightarrow w \Rightarrow w \in S_0\} \\ &= \rho(p) \cap \{v \mid \forall w. v \rightarrow w \Rightarrow w \in V\} \\ &= \rho(p) \cap V \\ &= \rho(p) \end{aligned}$$

Example

$$\begin{aligned} \llbracket \nu x. p \wedge \Box x \rrbracket \rho &\triangleq \text{FIX } \lambda S. \llbracket p \wedge \Box x \rrbracket \rho[S/x] \\ &= \text{FIX } \lambda S. \llbracket p \rrbracket \rho[S/x] \cap \llbracket \Box x \rrbracket \rho[S/x] \\ &= \text{FIX } \lambda S. \rho(p) \cap \{v \mid \forall w. v \rightarrow w \Rightarrow w \in \llbracket x \rrbracket \rho[S/x]\} \\ &= \text{FIX } \lambda S. \rho(p) \cap \{v \mid \forall w. v \rightarrow w \Rightarrow w \in S\} \end{aligned}$$

$$S_0 = V$$

$$S_1 = \rho(p)$$

$$S_2 = \rho(p) \cap \{v \mid \forall w. v \rightarrow w \Rightarrow w \in S_1\}$$

$$= \rho(p) \cap \{v \mid \forall w. v \rightarrow w \Rightarrow w \in \rho(p)\}$$

$$= \{v \text{ s.t. all nodes reachable in less than 2 moves are in } \rho(p)\}$$

Example

$$\begin{aligned} \llbracket \nu x. p \wedge \Box x \rrbracket \rho &\triangleq \text{FIX } \lambda S. \llbracket p \wedge \Box x \rrbracket \rho[S/x] \\ &= \text{FIX } \lambda S. \llbracket p \rrbracket \rho[S/x] \cap \llbracket \Box x \rrbracket \rho[S/x] \\ &= \text{FIX } \lambda S. \rho(p) \cap \{v \mid \forall w. v \rightarrow w \Rightarrow w \in \llbracket x \rrbracket \rho[S/x]\} \\ &= \text{FIX } \lambda S. \rho(p) \cap \{v \mid \forall w. v \rightarrow w \Rightarrow w \in S\} \end{aligned}$$

$$S_0 = V$$

$$S_1 = \rho(p)$$

$$\nu x. p \wedge \Box x$$

$$\text{AG } p$$

$$S_2 = \{v \text{ s.t. all nodes reachable in less than 2 moves are in } \rho(p)\}$$

$$S_n = \{v \text{ s.t. all nodes reachable in less than } n \text{ moves are in } \rho(p)\}$$

$$\bigcap_{n \in \mathbb{N}} S_n = \{v \text{ can only reach states in } \rho(p)\}$$

Invariants & possibly

Invariants $Inv(\psi) \triangleq \nu x. \psi \wedge \square x$

Possibly $Pos(\psi) \triangleq \mu x. \psi \vee \diamond x$

Example

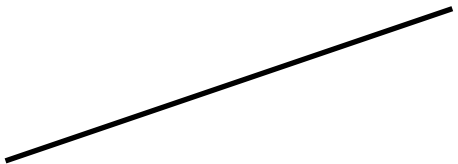
which temporal formula?

$$\mu x. q \vee (p \wedge \diamond x)$$

$$E(p \text{ U } q) \quad (\text{CTL})$$

$$\mu x. q \vee (p \wedge \diamond x \wedge \square x)$$

$$A(p \text{ U } q) \quad (\text{LTL / CTL})$$

$$\nu x. \underbrace{\mu y. (p \wedge \diamond x) \vee \diamond y}$$


$$E \text{ G } F \ p \quad (\text{CTL}^*)$$

after a finite number of steps you reach a state where

1) p holds

2) there is a next step where the property holds recursively

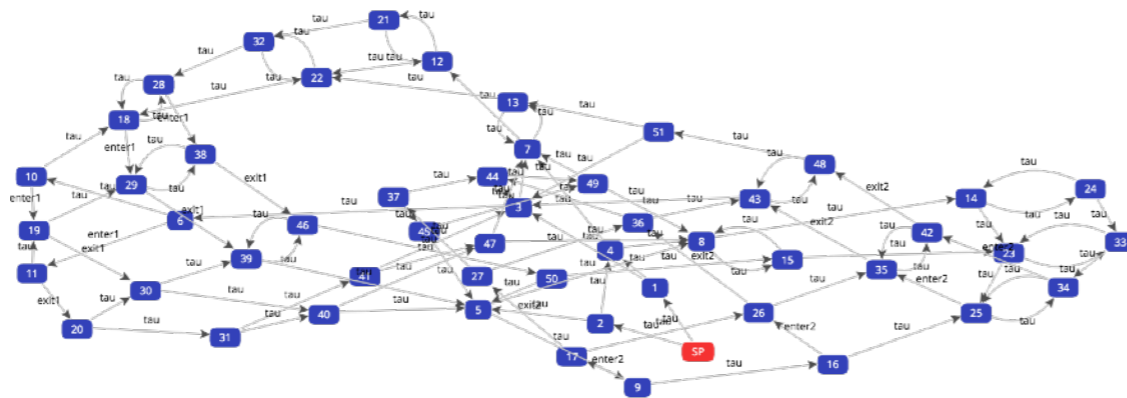
mu-calculus with labels

$\psi ::= \dots$ set of labels
| $\diamond_L \psi$
| $\square_L \psi$

$$\llbracket \diamond_L \psi \rrbracket \rho \triangleq \{v \mid \exists \mu \in L. \exists w \in \llbracket \psi \rrbracket \rho. v \xrightarrow{\mu} w\}$$

$$\llbracket \square_L \psi \rrbracket \rho \triangleq \{v \mid \forall \mu \in L. \forall w. v \xrightarrow{\mu} w \Rightarrow w \in \llbracket \psi \rrbracket \rho\}$$

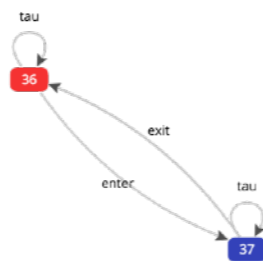
Space reduction



?
 $\models \psi$

identify
bisimilar
states

\Leftrightarrow



?
 $\models \psi$