

**Models of computation (MOD) 2016/17**  
 Appello straordinario – April 6, 2017

[Ex. 1] Let us extend IMP with the command

$$x := c$$

whose denotational semantics is

$$\mathcal{C} \llbracket x := c \rrbracket \sigma \stackrel{\text{def}}{=} (\lambda \sigma_1. \sigma[\sigma_1(x)/x])^* (\mathcal{C} \llbracket c \rrbracket \sigma)$$

1. Define the operational semantics for the new construct.
2. Extend the proofs of correctness and completeness between the operational and the denotational semantics to account for the new construct.

[Ex. 2] Let  $(S, \prec)$  be a set  $S$  with a binary relation  $\prec \subseteq S \times S$  such that for all  $s \in S$  the set  $[s] \stackrel{\text{def}}{=} \{ x \mid x \prec s \}$  is finite, and let  $f : (\wp(S), \subseteq) \rightarrow (\wp(S), \subseteq)$  be the function over the  $\text{CPO}_\perp (\wp(S), \subseteq)$  such that, for any  $X \in \wp(S)$

$$f(X) \stackrel{\text{def}}{=} \{ y \mid [y] \subseteq X \}$$

1. Is  $(S, \prec)$  always well founded? (If not, exhibit a counterexample)
2. Prove that  $f$  is monotone.
3. Prove that  $f$  is continuous.

[Ex. 3] Let us consider the HOFL term

$$t \stackrel{\text{def}}{=} \mathbf{rec} f. \lambda x. (\mathbf{fst}(x), \mathbf{snd}(f x))$$

1. Find the principal type of  $t$ .
2. Compute the denotational semantics of  $t$ .

[Ex. 4] Let us consider the CCS process

$$p \stackrel{\text{def}}{=} \mathbf{rec} x. ( (\alpha.x \mid \bar{\alpha}.\mathbf{nil}) \backslash \alpha + \beta.x )$$

and let  $p \backslash_\alpha^0 \stackrel{\text{def}}{=} p$  and  $p \backslash_\alpha^{n+1} \stackrel{\text{def}}{=} (p \backslash_\alpha^n \mathbf{nil}) \backslash \alpha$  for  $n \geq 0$ .

1. Draw, at least in part, the LTS for the process  $p$ .
2. What are the transitions leaving from  $p \backslash_\alpha^n$ ?
3. Prove that  $p \backslash_\alpha^n$  is strongly bisimilar to  $p \backslash_\alpha^m$  for any  $n, m \geq 0$ .