

# Time Series - Shapelet/Motif Discovery

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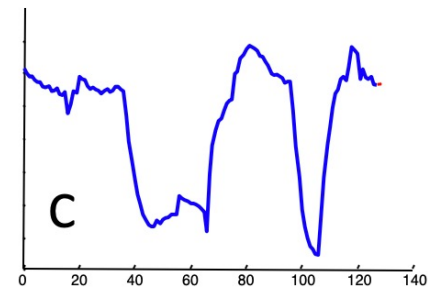
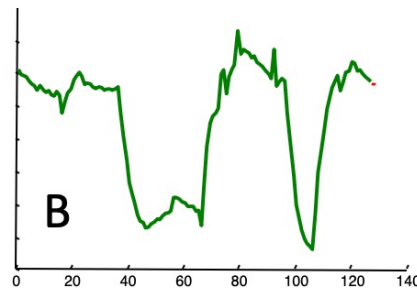
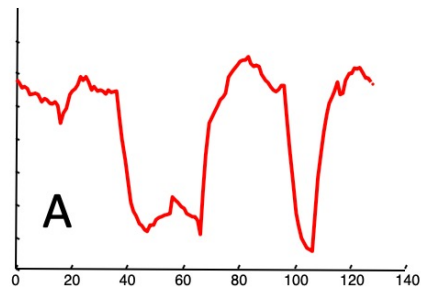
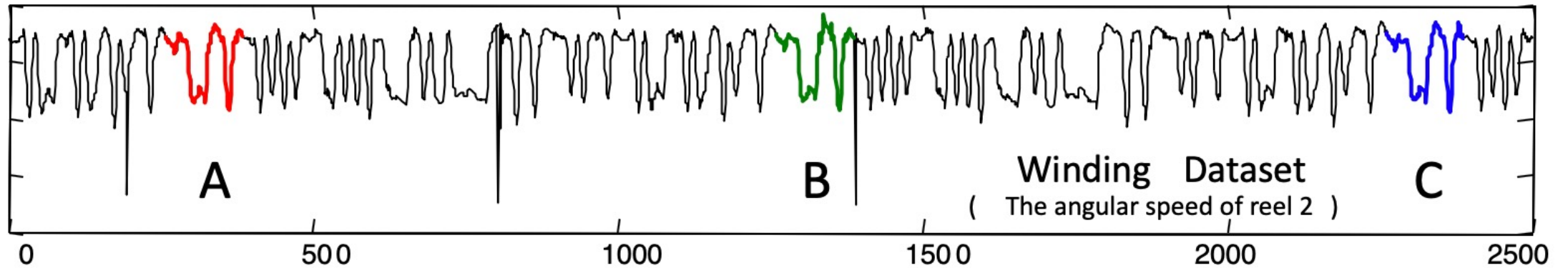


# Motif

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# Time Series Motif Discovery

- Finding repeated patterns, i.e., pattern mining.
- Are there any repeated patterns, of length  $m$  in the TS?



# Why Find Motifs?

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- Mining **association rules** in TS requires the discovery of motifs. These are referred to as primitive shapes and frequent patterns.
- Several **TS classifiers** work by constructing typical prototypes of each class. These prototypes may be considered motifs.
- Many **TS anomaly detection** algorithms consist of modeling normal behavior with a set of typical shapes (which we see as motifs), and detecting future patterns that are dissimilar to all typical shapes.

# Matrix Profile

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- The Matrix Profile (MP) is a data structure that annotates a TS and can be exploited for many purposes: e.g. efficient Motif Discovery.
- Given a time series,  $T$  and a desired subsequence length,  $m$ .



$m$

# Matrix Profile

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We can use sliding window of length  $m$  to extract all subsequences of length  $m$ .

$|T| - m + 1$

...

# Matrix Profile



We can then compute the pairwise distance among these subsequences.

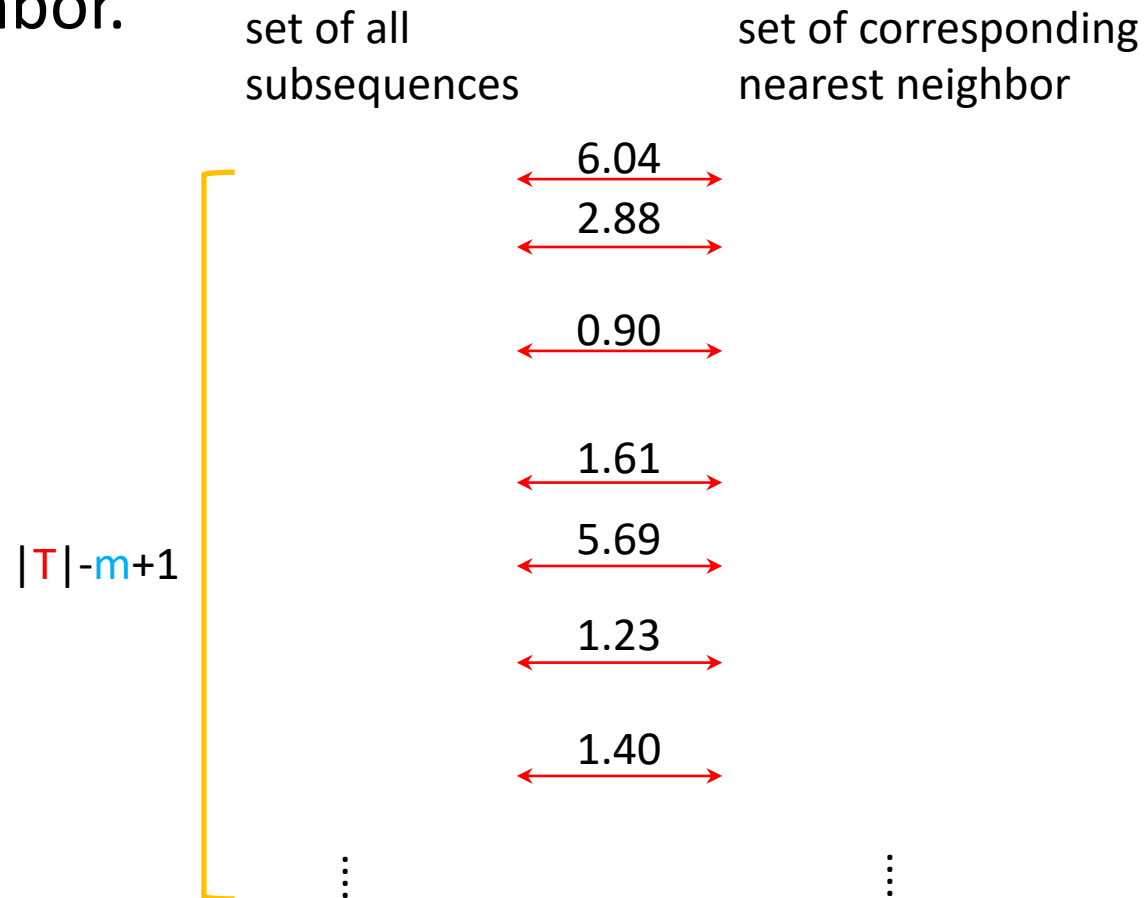
0	7.6952	7.7399	...
7.6952	0	7.7106	...
7.7399	7.7106	0	...
...	...	...	...

...

$|T| - m + 1$

# Matrix Profile

- For each subsequence we keep only the distance with the closest nearest neighbor.





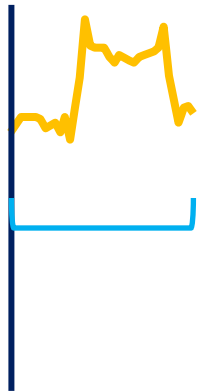
# Matrix Profile

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- The **distance** to the corresponding **nearest neighbor** of each subsequence can be stored in a vector called **matrix profile P**.

time  
series,  $T$

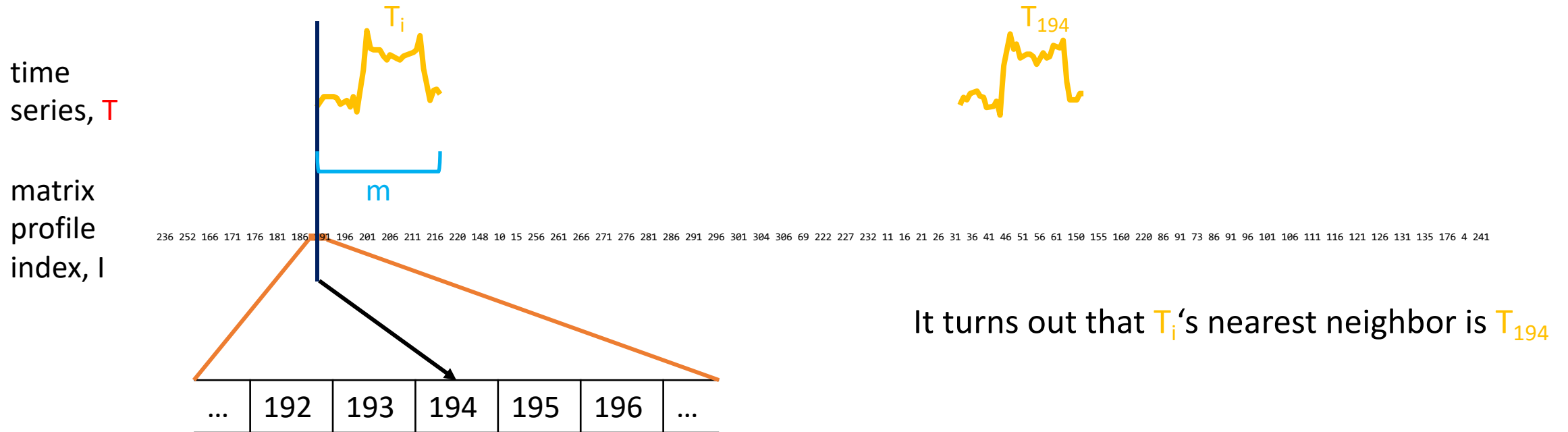
matrix  
profile,  $P$



The matrix profile value at location  $i$  is the distance between  $T_i$  and its nearest neighbor

# Matrix Profile

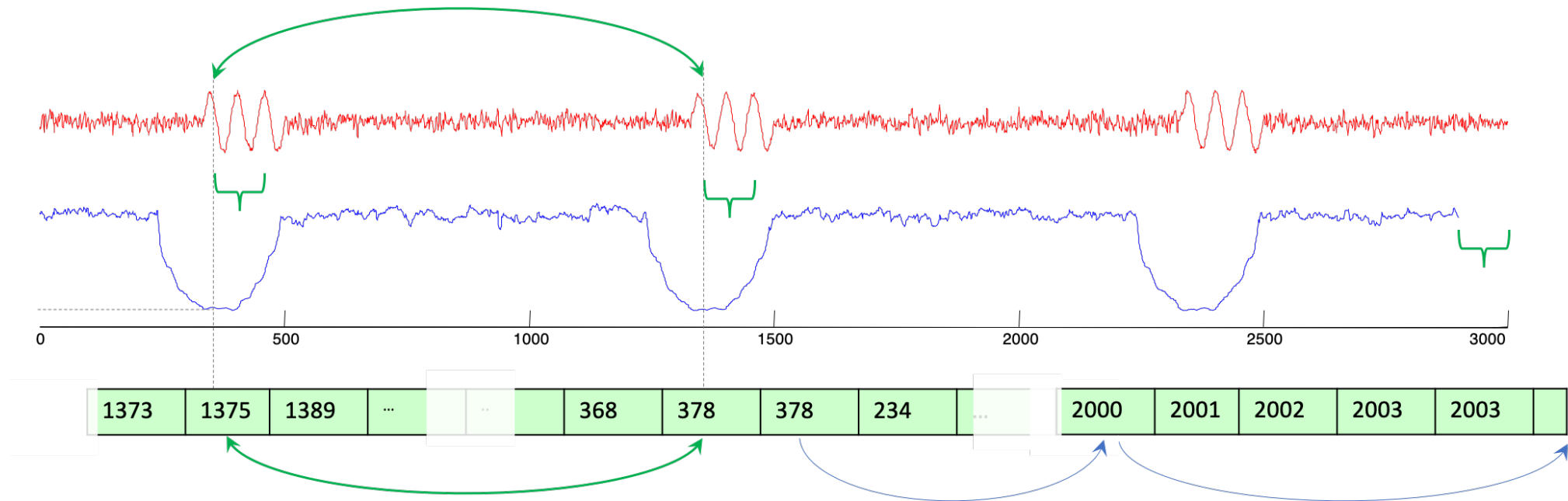
- The **index of corresponding nearest neighbor** of each subsequence is also stored in a vector called **matrix profile index**.



The matrix profile value at location  $i$  is the distance between  $T_i$  and its nearest neighbor

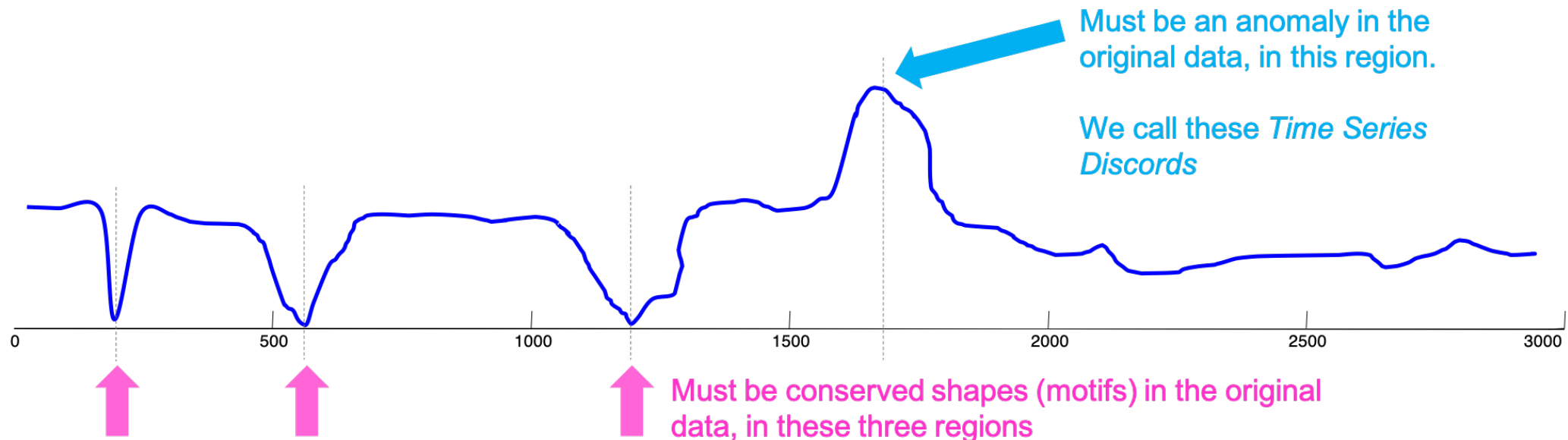
# Matrix Profile

- The MP index allows to find the nearest neighbor to any subsequence in constant time.
- Note that the pointers in the matrix profile index are not necessarily symmetric.
- If A points to B, then B may or may not point to A
- The classic TS motif: the two smallest values in the MP must have the same value, and their pointers must be mutual.



# How to “read” a Matrix Profile

- For relatively low values, you know that the subsequence in the original TS must have (at least one) relatively similar subsequence elsewhere in the data (such regions are “motifs”)
- For relatively high values, you know that the subsequence in the original TS must be unique in its shape (such areas are anomalies).



# How to Compute Matrix Profile?

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- Given a time series,  $T$  and a desired subsequence length,  $m$ .



inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf
-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----

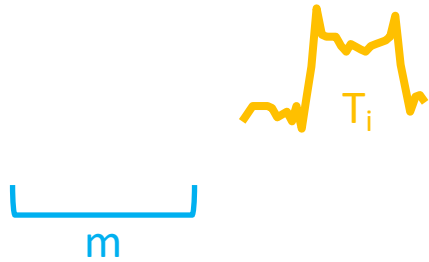
Matrix profile is initialized as inf vector

This is just a toy example, so the values and the vector length does not fit the time series shown above

# How to Compute Matrix Profile?

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- Given a time series,  $T$  and a desired subsequence length,  $m$ .

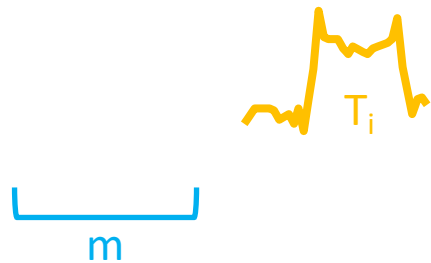


inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf
-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----

At the first iteration, a subsequence  $T_i$  is randomly selected from  $T$

# How to Compute Matrix Profile?

- Given a time series,  $T$  and a desired subsequence length,  $m$ .



inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf
-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----

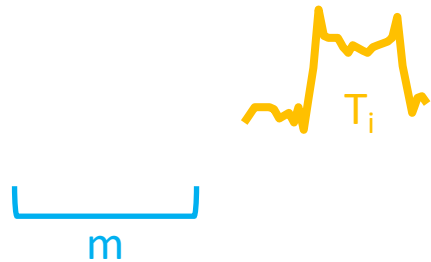
We compute the distances between  $T_i$  and every subsequences from  $T$  (time complexity =  $O(|T| \log(|T|))$ )  
We then put the distances in a vector based on the position of the subsequences

3	2	0	5	3	4	5	1	2	9	8	4	2	3	4	8	6	2	1
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

↖ The distance between  $T_i$  and  $T_1$  (first subsequence) is 3

# How to Compute Matrix Profile?

- Given a time series,  $T$  and a desired subsequence length,  $m$ .



inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf
-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----

We compute the distances between  $T_i$  and every subsequences from  $T$  (time complexity =  $O(|T| \log(|T|))$ )  
We then put the distances in a vector based on the position of the subsequences

3	2	0	5	3	4	5	1	2	9	8	4	2	3	4	8	6	2	1
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---



Let say  $T_i$  happen to be the third subsequences, therefore the third value in the distance vector is 0

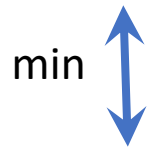


# How to Compute Matrix Profile?

- Given a time series,  $T$  and a desired subsequence length,  $m$ .



inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf
-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----

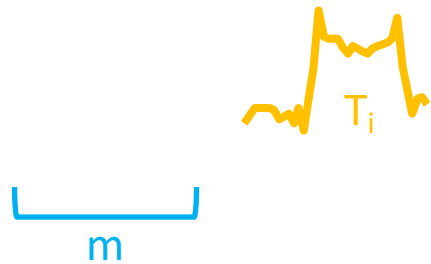


Matrix profile is updated by apply elementwise minimum to these two vectors

3	2	0	5	3	4	5	1	2	9	8	4	2	3	4	8	6	2	1
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

# How to Compute Matrix Profile?

- Given a time series,  $T$  and a desired subsequence length,  $m$ .



3	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf
---	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----

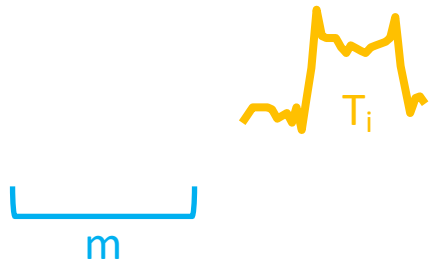
min

Matrix profile is updated by apply elementwise minimum to these two vectors

3	2	0	5	3	4	5	1	2	9	8	4	2	3	4	8	6	2	1
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

# How to Compute Matrix Profile?

- Given a time series,  $T$  and a desired subsequence length,  $m$ .



3	2	inf	5	3	4	5	1	2	9	8	4	2	3	4	8	6	2	1
---	---	-----	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

After we finish to update matrix profile for the first iteration

3	2	0	5	3	4	5	1	2	9	8	4	2	3	4	8	6	2	1
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

# How to Compute Matrix Profile?

- Given a time series,  $T$  and a desired subsequence length,  $m$ .

  
m



3	2	inf	5	3	4	5	1	2	9	8	4	2	3	4	8	6	2	1
---	---	-----	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

In the second iteration, we randomly select another subsequence  $T_j$  and it happens to be the 12<sup>th</sup> subsequence

# How to Compute Matrix Profile?

- Given a time series,  $T$  and a desired subsequence length,  $m$ .

  
m



3	2	inf	5	3	4	5	1	2	9	8	4	2	3	4	8	6	2	1
---	---	-----	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

Once again, we compute the distance between  $T_j$  and every subsequences of  $T$

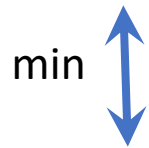
2	3	1	4	4	3	6	2	1	5	8	0	2	3	5	9	4	2	2
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

# How to Compute Matrix Profile?

- Given a time series,  $T$  and a desired subsequence length,  $m$ .



3	2	inf	5	3	4	5	1	2	9	8	4	2	3	4	8	6	2	1
---	---	-----	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---



The same elementwise minimum

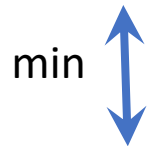
2	3	1	4	4	3	6	2	1	5	8	0	2	3	5	9	4	2	2
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

# How to Compute Matrix Profile?

- Given a time series,  $T$  and a desired subsequence length,  $m$ .



2	2	inf	5	3	4	5	1	2	9	8	4	2	3	4	8	6	2	1
---	---	-----	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---



The same elementwise minimum

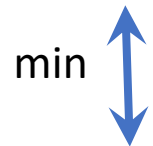
2	3	1	4	4	3	6	2	1	5	8	0	2	3	5	9	4	2	2
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

# How to Compute Matrix Profile?

- Given a time series,  $T$  and a desired subsequence length,  $m$ .



2	2	inf	5	3	4	5	1	2	9	8	4	2	3	4	8	6	2	1
---	---	-----	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---



The same elementwise minimum

2	3	1	4	4	3	6	2	1	5	8	0	2	3	5	9	4	2	2
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---



# How to Compute Matrix Profile?

- Given a time series,  $T$  and a desired subsequence length,  $m$ .



$m$

2	2	1	5	3	4	5	1	2	9	8	4	2	3	4	8	6	2	1
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

min

The same elementwise minimum

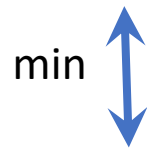
2	3	1	4	4	3	6	2	1	5	8	0	2	3	5	9	4	2	2
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

# How to Compute Matrix Profile?

- Given a time series,  $T$  and a desired subsequence length,  $m$ .



2	2	1	5	3	4	5	1	2	9	8	4	2	3	4	8	6	2	1
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---



2	3	1	4	4	3	6	2	1	5	8	0	2	3	5	9	4	2	2
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

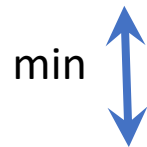
We repeat the two steps (distance computation and update) until we have used every subsequences

# How to Compute Matrix Profile?

- Given a time series,  $T$  and a desired subsequence length,  $m$ .



2	2	1	5	3	4	5	1	2	9	8	4	2	3	4	8	6	2	1
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---



2	3	1	4	4	3	6	2	1	5	8	0	2	3	5	9	4	2	2
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

There are  $|T|$  subsequences and the distance computation is  $O(|T|\log(|T|))$

The overall time complexity is  $O(|T|^2\log(|T|))$

# Motif Discovery From Matrix Profile

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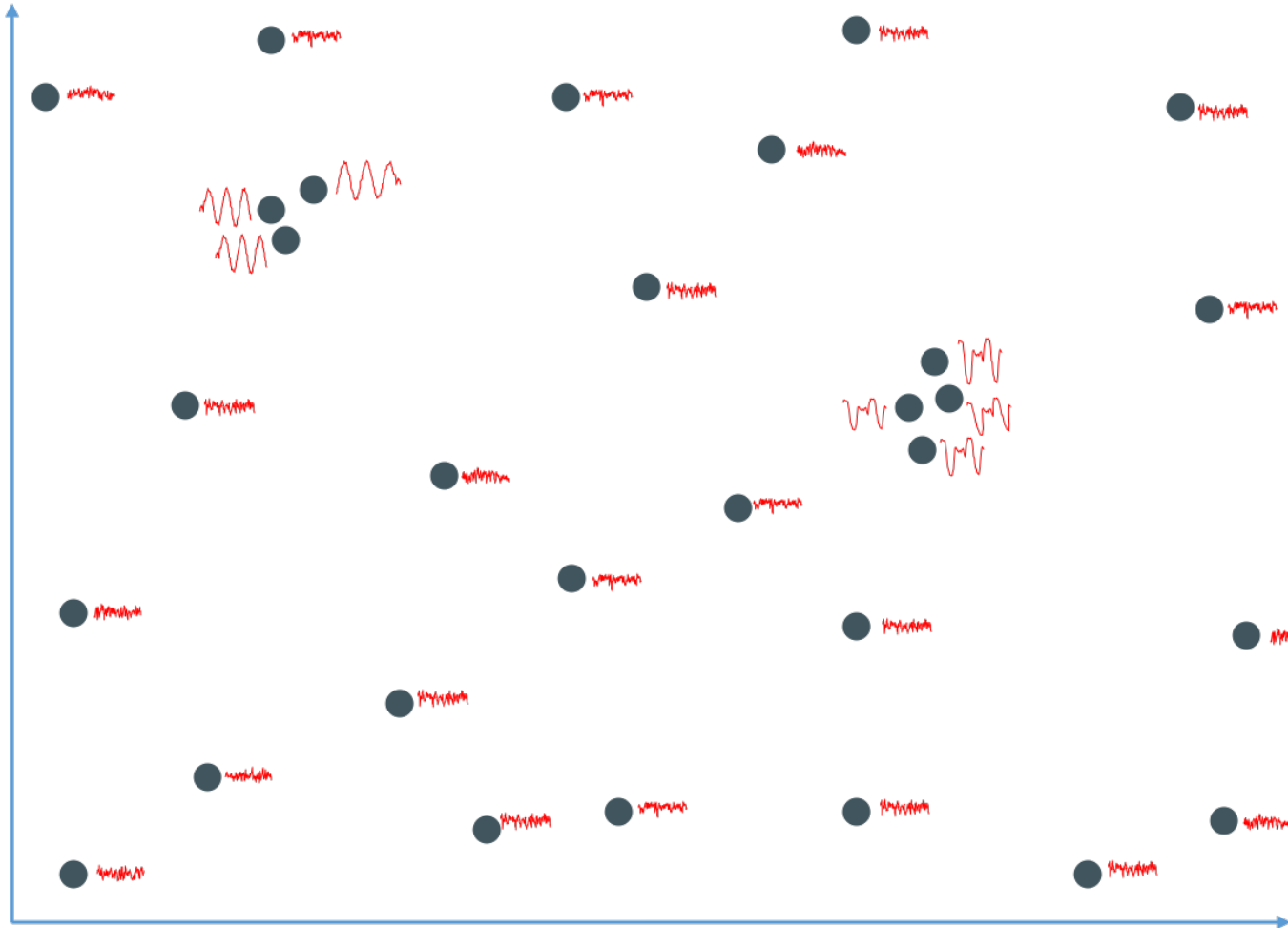
time  
series,  $T$

matrix  
profile,  $P$

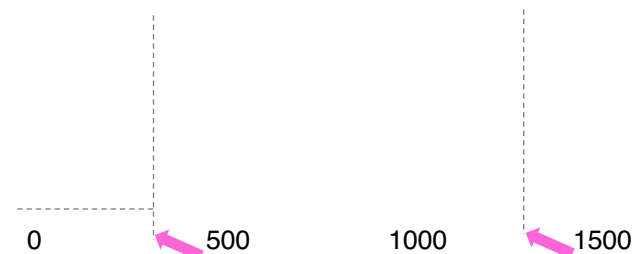


Local minimums are corresponding to motifs

# Motif Discovery From Matrix Profile



- It is sometime useful to think of time series subsequences as points in m-dimensional space.
- In this view, dense regions in the m-dimensional space correspond to regions of the time series that have a low corresponding MP.



# Shapelet

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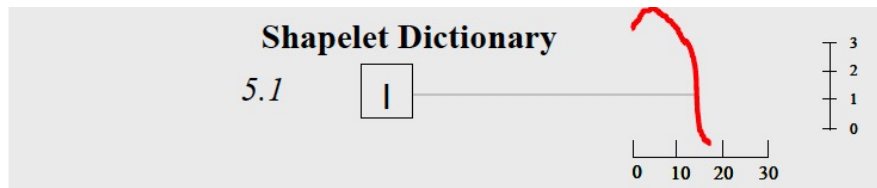
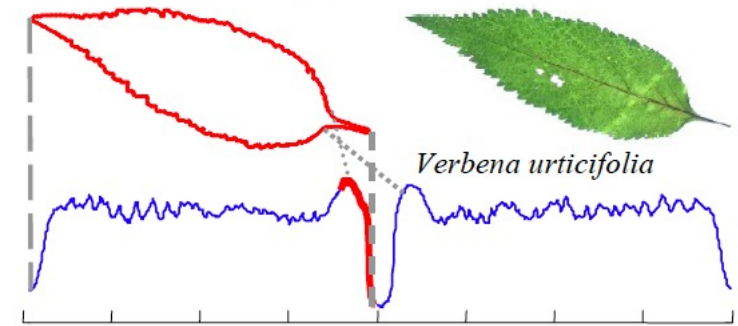
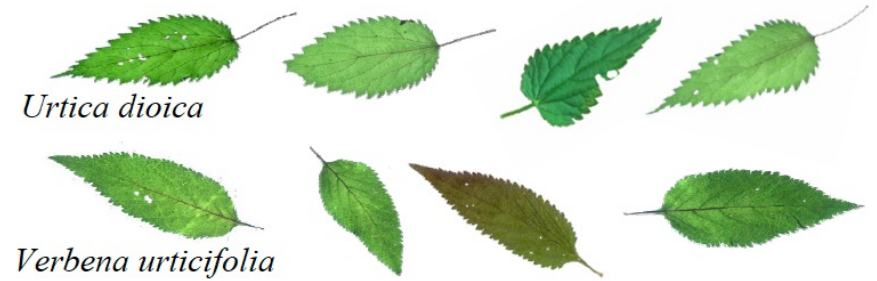
# Time Series Classification

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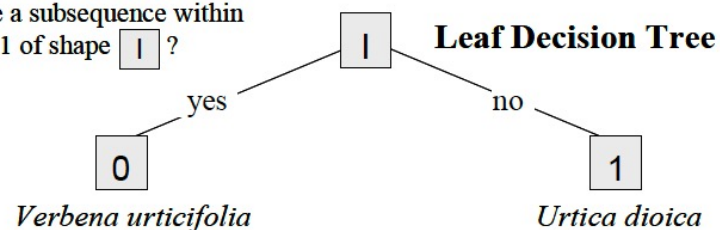
- Given a set  $X$  of  $n$  time series,  $X = \{x_1, x_2, \dots, x_n\}$ , each time series has  $m$  ordered values  $x_i = \langle x_{t1}, x_{t2}, \dots, x_{tm} \rangle$  and a class value  $c_i$ .
- The objective is to find a function  $f$  that maps from the space of possible time series to the space of possible class values.
- Generally, it is assumed that all the TS have the same length  $m$ .

# Shapelet-based Classification

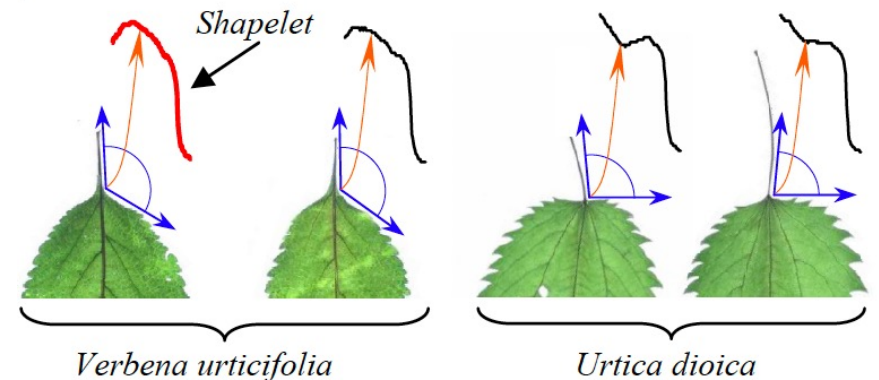
1. Represent a TS as a vector of distances with representative subsequences, namely shapelets.
2. Use it as input for machine learning classifiers.



Does  $Q$  have a subsequence within a distance 5.1 of shape 1?



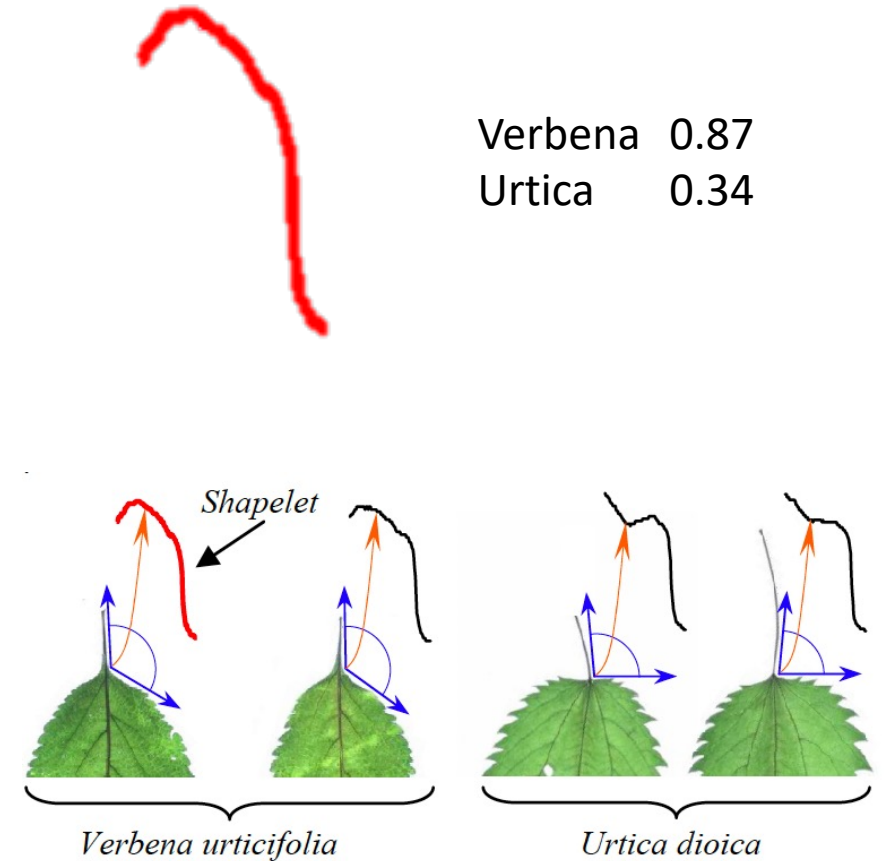
3.2	8.7
1.4	7.9
6.7	4.2
9.2	3.4





# Time Series Shapelets

- Shapelets are TS subsequences which are maximally representative of a class.
- Shapelets can provide interpretable results, which may help domain practitioners better understand their data.
- Shapelets can be significantly more accurate/robust because they are *local features*, whereas most other state-of-the-art TS classifiers consider *global features*.



# Finding Shapelets

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FindingShapeletBF (dataset  $\mathbf{D}$ ,  $MAXLEN$ ,  $MINLEN$ )

```
1  candidates  $\leftarrow$  GenerateCandidates( $\mathbf{D}$ ,  $MAXLEN$ ,  $MINLEN$ )
2  bsf_gain  $\leftarrow$  0
3  For each  $S$  in candidates
4      gain  $\leftarrow$  CheckCandidate( $\mathbf{D}$ ,  $S$ )
5      If gain  $>$  bsf_gain
6          bsf_gain  $\leftarrow$  gain
7          bsf_shapelet  $\leftarrow$   $S$ 
8      EndIf
9  EndFor
10 Return bsf_shapelet
```

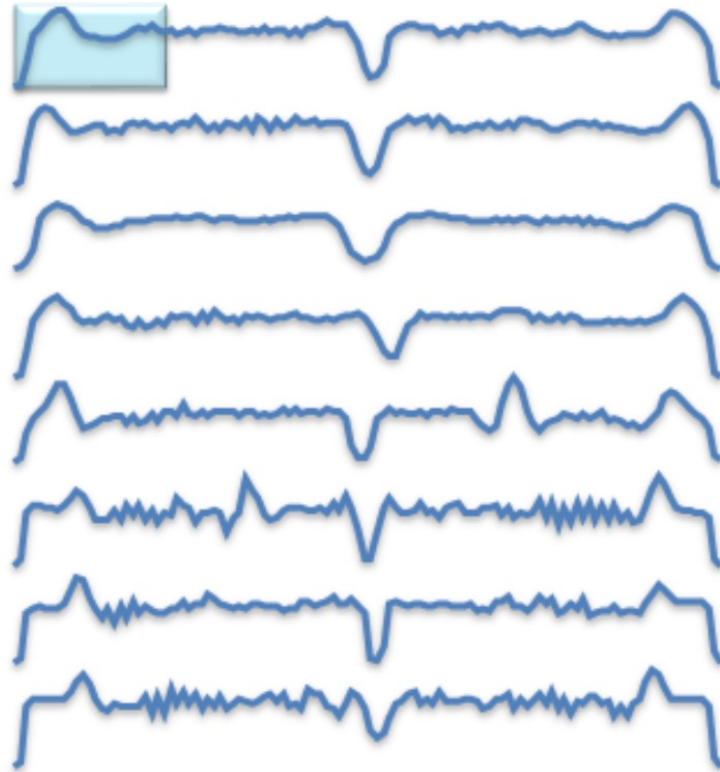
# Generate Candidate

Sliding a **window of size  $l$**  across all of the time series objects in the dataset  $D$ , extracts all of the possible candidates and adds them to the pool

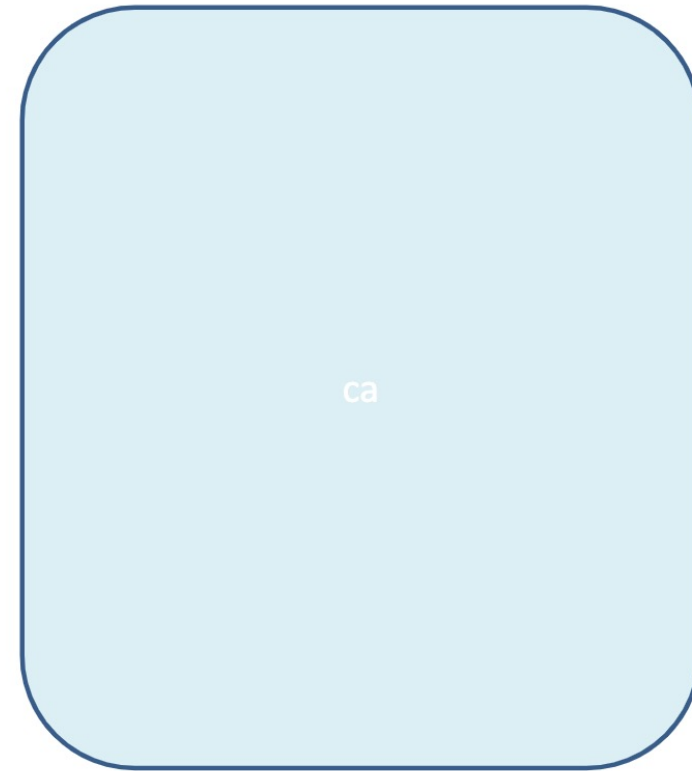
GenerateCandidates (dataset $D$ , $MAXLEN$ , $MINLEN$ )	
1	$pool \leftarrow \emptyset$
2	$l \leftarrow MAXLEN$
3	<b>While</b> $l \geq MINLEN$
4	<b>For</b> $T$ in $D$
5	$pool \leftarrow pool \cup S_T^l$
6	<b>EndFor</b>
7	$l \leftarrow l - 1$
8	<b>EndWhile</b>
9	<b>Return</b> $pool$

# Extract Subsequences of all Possible Lengths

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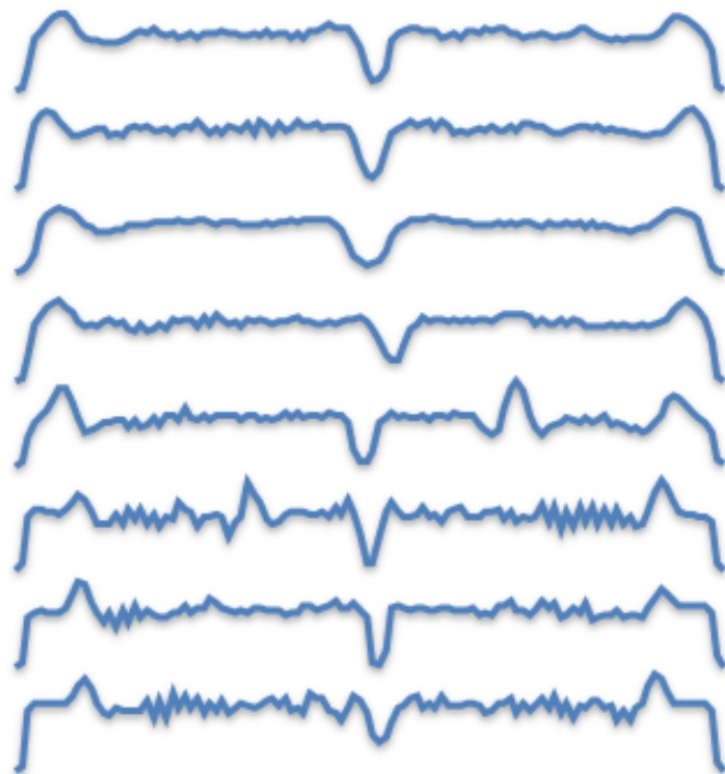


Candidates Pool

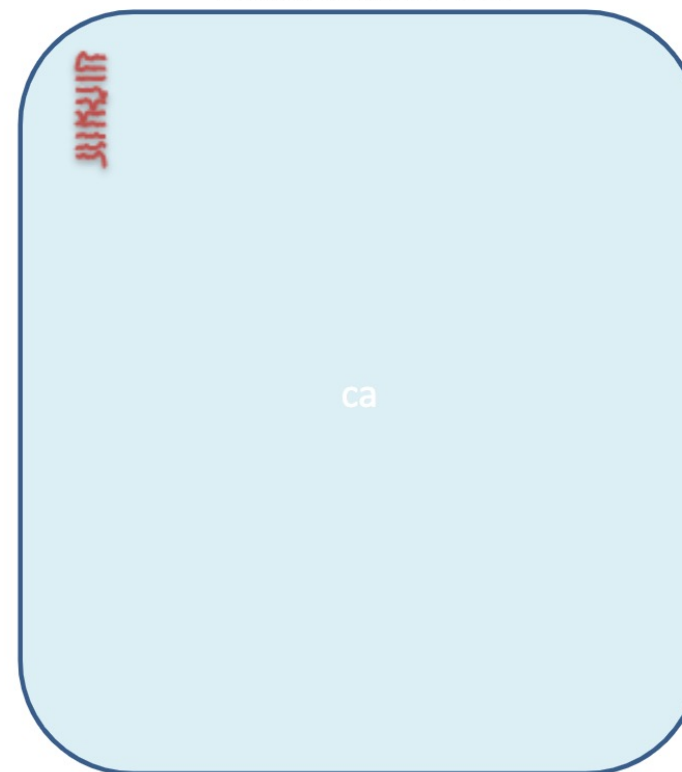


# Extract Subsequences of all Possible Lengths

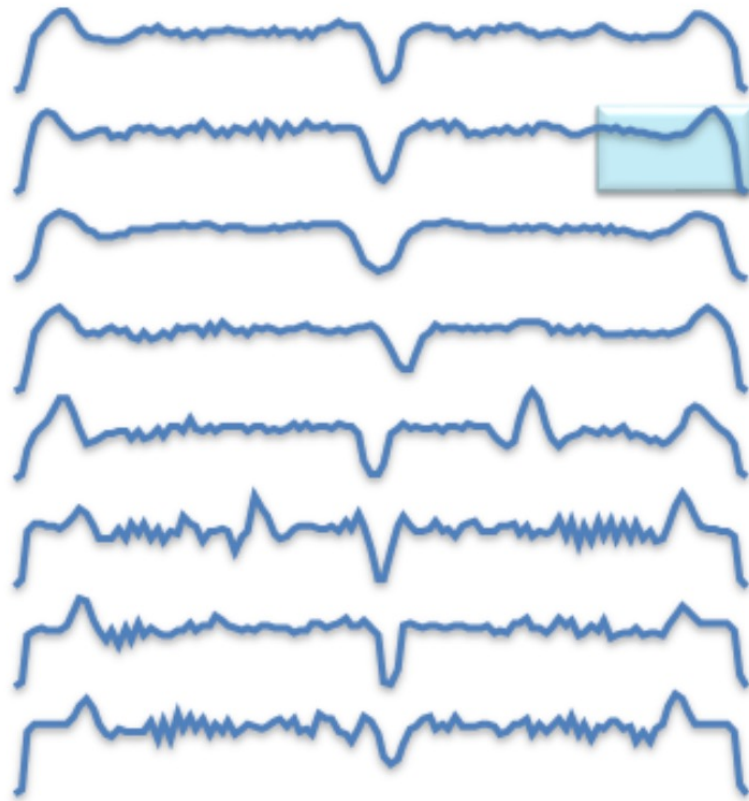
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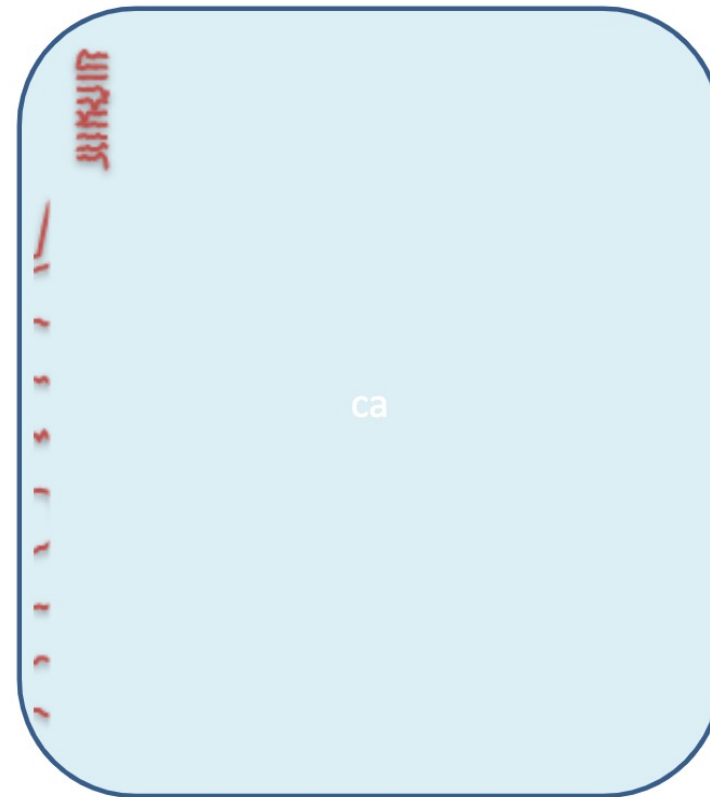
Candidates Pool



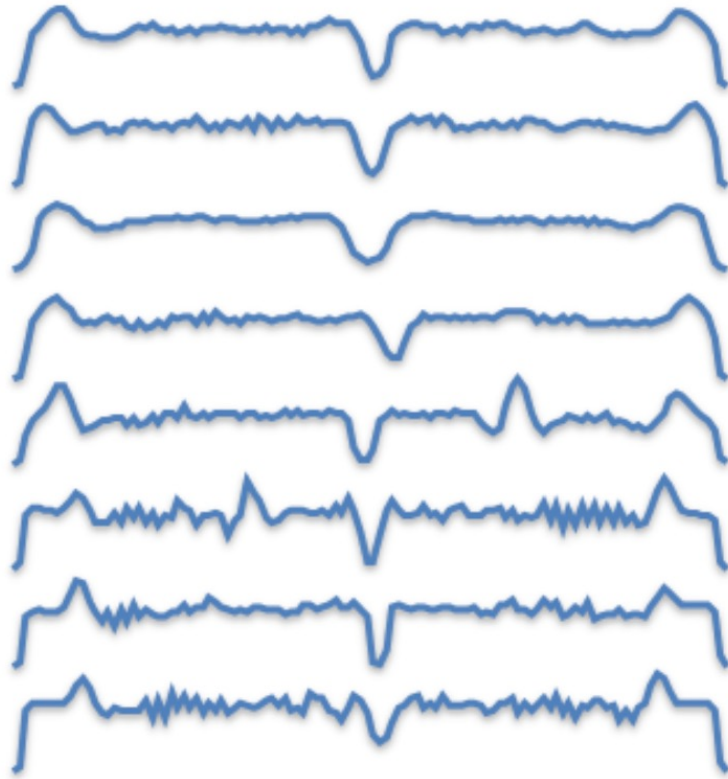
# Extract Subsequences of all Possible Lengths



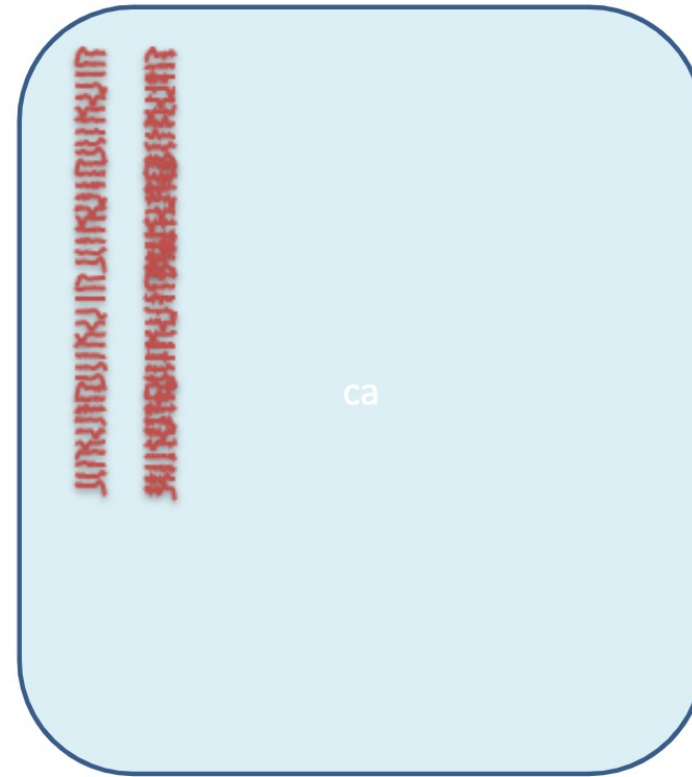
Candidates Pool



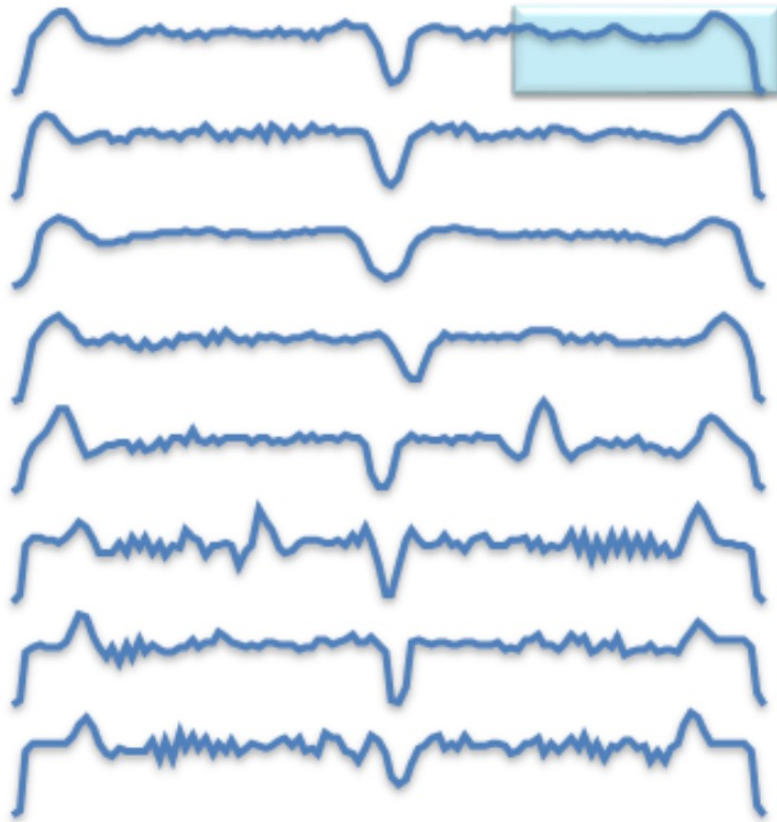
# Extract Subsequences of all Possible Lengths



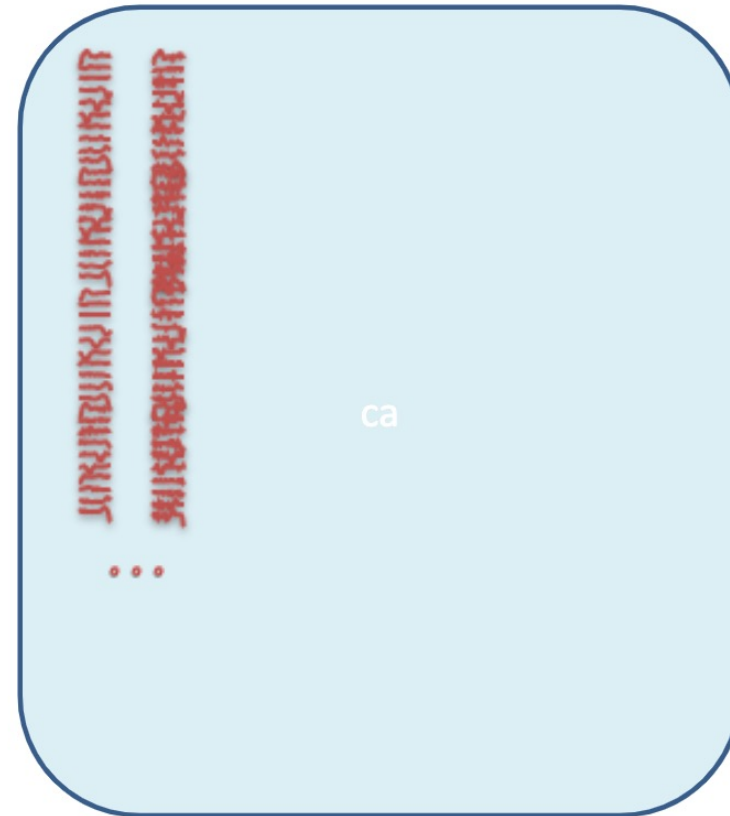
Candidates Pool



# Extract Subsequences of all Possible Lengths



Candidates Pool





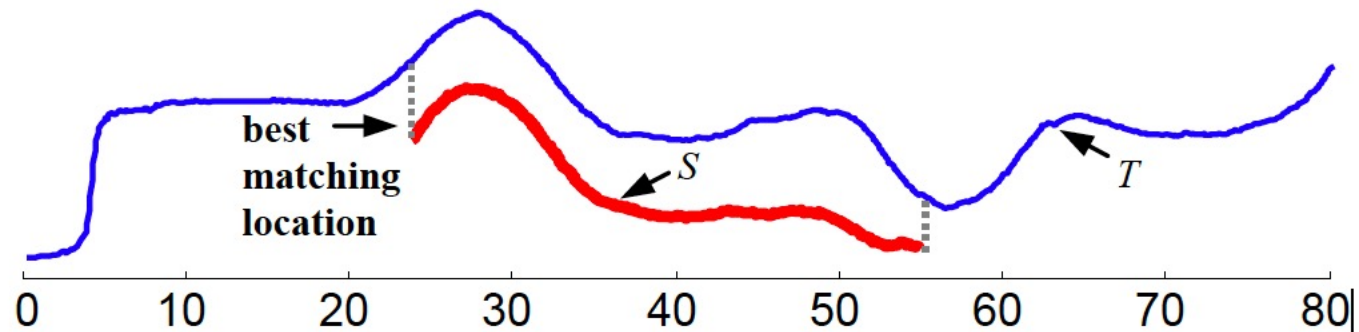
# Check Candidates

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CheckCandidate (dataset $\mathbf{D}$ , shapelet candidate $S$ )	
1	$objects\_histogram \leftarrow \emptyset$
2	<b>For each</b> $T$ <b>in</b> $\mathbf{D}$
3	$dist \leftarrow \text{SubsequenceDist}(T, S)$
4	insert $T$ into $objects\_histogram$ by the key $dist$
5	<b>EndFor</b>
6	<b>Return</b> CalculateInformationGain( $objects\_histogram$ )

# Distance with a Subsequence

- Distance from the TS to the subsequence  $SubsequenceDist(T, S)$  is a distance function that takes time series  $T$  and subsequence  $S$  as inputs and returns a nonnegative value  $d$ , which is the distance from  $T$  to  $S$ .
- $SubsequenceDist(T, S) = \min(Dist(S, S')), \text{ for } S' \in S_T^{|S|}$
- where  $S_T^{|S|}$  is the set of all possible subsequences of  $T$
- Intuitively, it is the distance between  $S$  and its best matching location in  $T$ .



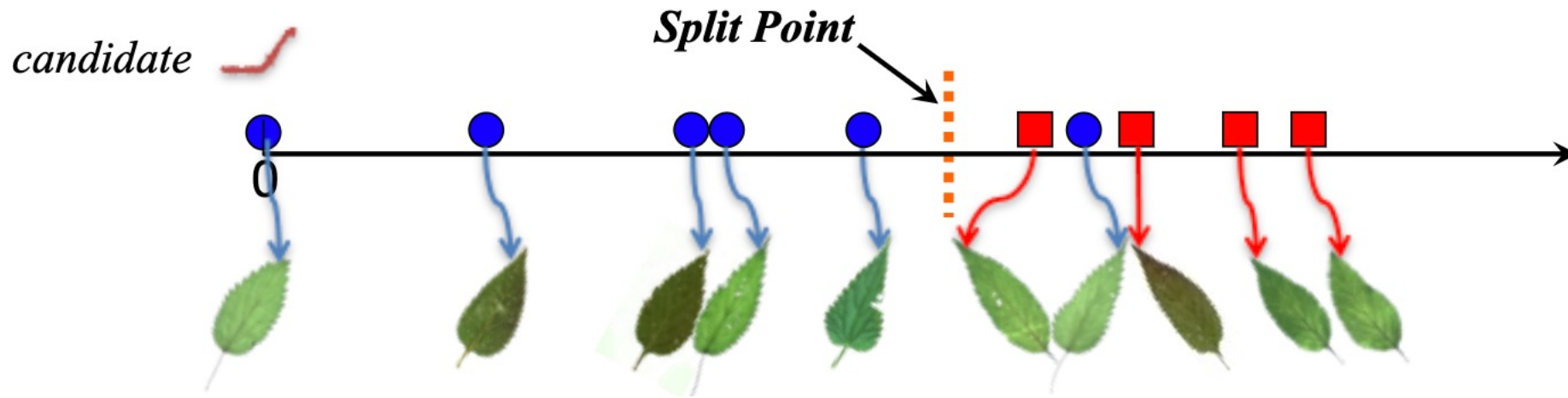
# Check Candidates

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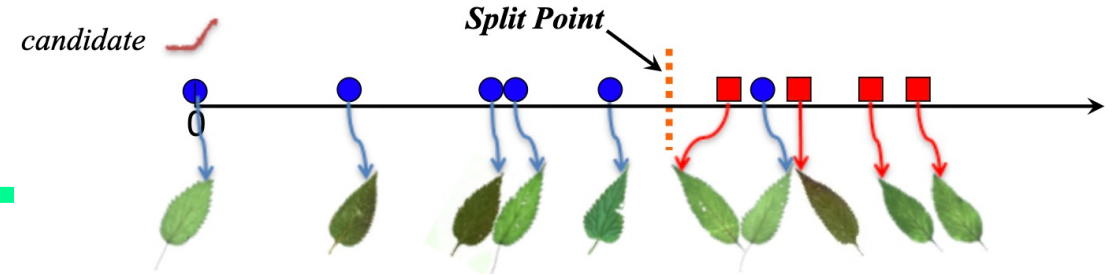
CheckCandidate (dataset $\mathbf{D}$ , shapelet candidate $S$ )	
1	$objects\_histogram \leftarrow \emptyset$
2	<b>For each</b> $T$ <b>in</b> $\mathbf{D}$
3	$dist \leftarrow \text{SubsequenceDist}(T, S)$
4	insert $T$ into $objects\_histogram$ by the key $dist$
5	<b>EndFor</b>
6	<b>Return</b> CalculateInformationGain( $objects\_histogram$ )

# Testing The Utility of a Candidate Shapelet

- Arrange the TSs in the dataset  $D$  based on the distance from the candidate.
- Find the optimal split point that maximizes the information gain (same as for Decision Tree classifiers)
- Pick the candidate achieving best utility as the shapelet

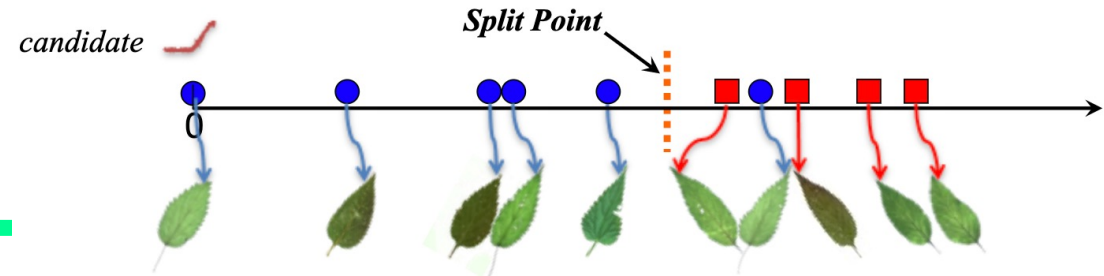


# Entropy



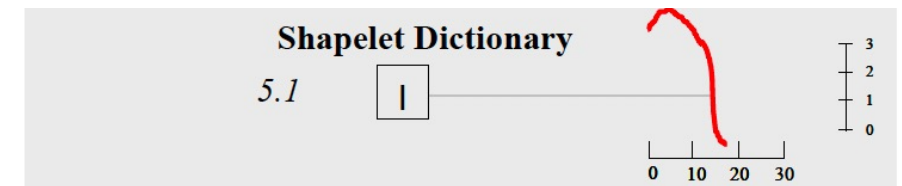
- A TS dataset  $D$  consists of two classes, A and B.
- Given that the proportion of objects in class A is  $p(A)$  and the proportion of objects in class B is  $p(B)$ ,
- The **Entropy** of  $D$  is:  $I(D) = -p(A)\log(p(A)) - p(B)\log(p(B))$ .
- Given a strategy that divides the  $D$  into two subsets  $D_1$  and  $D_2$ , the information remaining in the dataset after splitting is defined by the weighted average entropy of each subset.
- If the fraction of objects in  $D_1$  is  $f(D_1)$  and in  $D_2$  is  $f(D_2)$ ,
- The total entropy of  $D$  after splitting is  $\hat{I}(D) = f(D_1)I(D_1) + f(D_2)I(D_2)$ .

# Information Gain

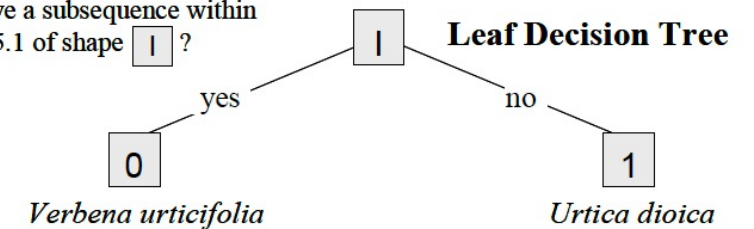


Split point  
distance from  
shapelet = 5.1

- Given a certain split strategy  $sp$  which divides  $D$  into two subsets  $D_1$  and  $D_2$ , the entropy before and after splitting is  $I(D)$  and  $\hat{I}(D)$ .
- The **information gain** for this splitting rule is:
- $Gain(sp) = I(D) - \hat{I}(D) =$
- $= I(D) - f(D_1)I(D_1) + f(D_2)I(D_2).$
- We use the distance from  $T$  to a shapelet  $S$  as the splitting rule  $sp$ .



Does  $Q$  have a subsequence within a distance 5.1 of shape  $I$ ?



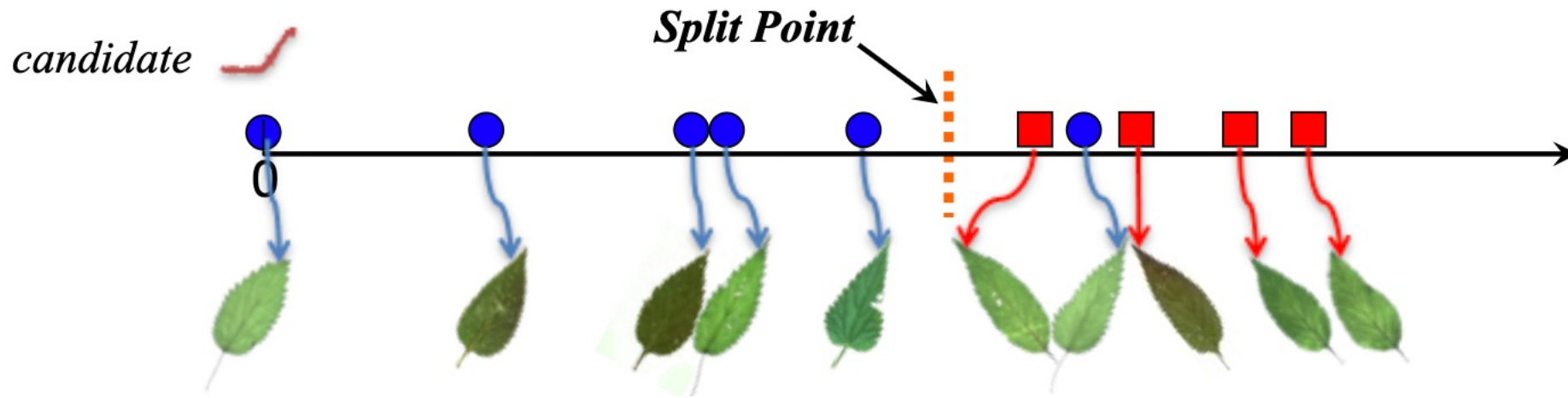
# Problem

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- The total number of candidate is 
$$\sum_{l=MINLEN}^{MAXLEN} \sum_{T_i \in D} (|T_i| - l + 1)$$
- For each candidate you have to compute the distance between this candidate and each training sample (space inefficiency)
- For instance
  - 200 instances with length 275
  - 7,480,200 shapelet candidates

# Speedup

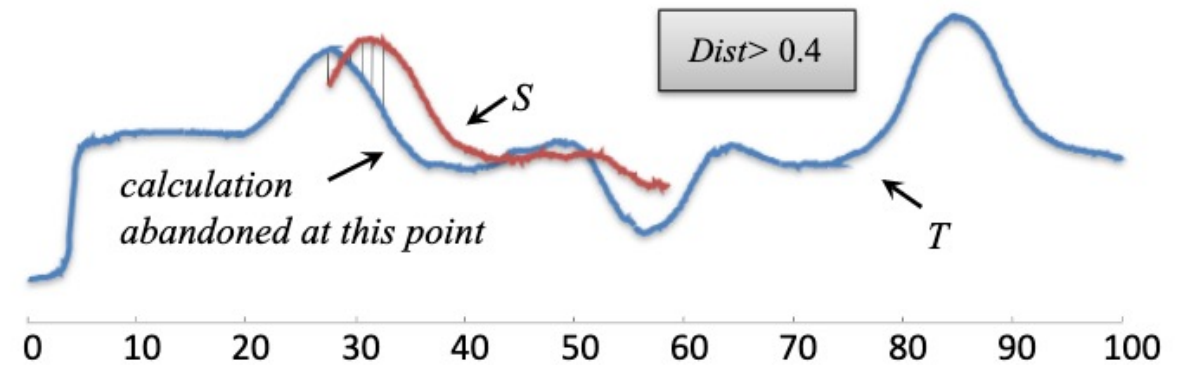
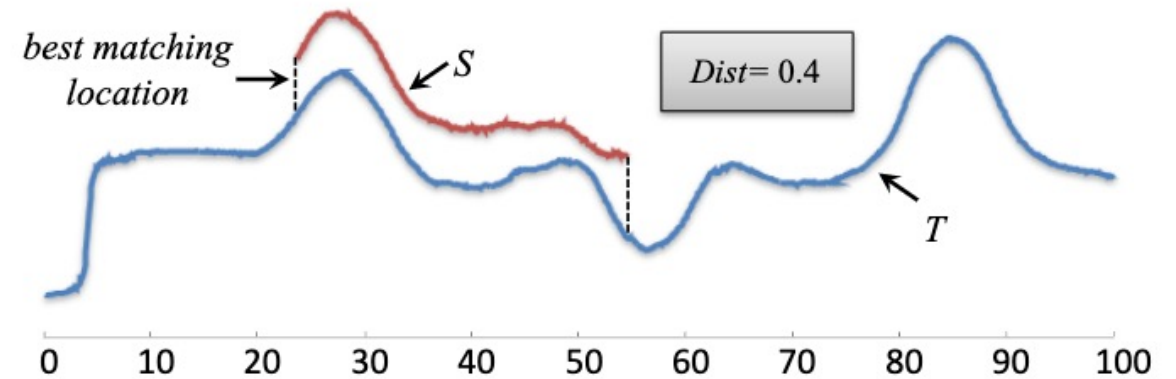
- Distance calculations from TSs to shapelet candidates is expensive.
- Reduce the time in two ways
  - Distance Early Abandon
  - reduce the distance computation time between two TS
- Admissible Entropy Pruning
  - reduce the number of distance calculations





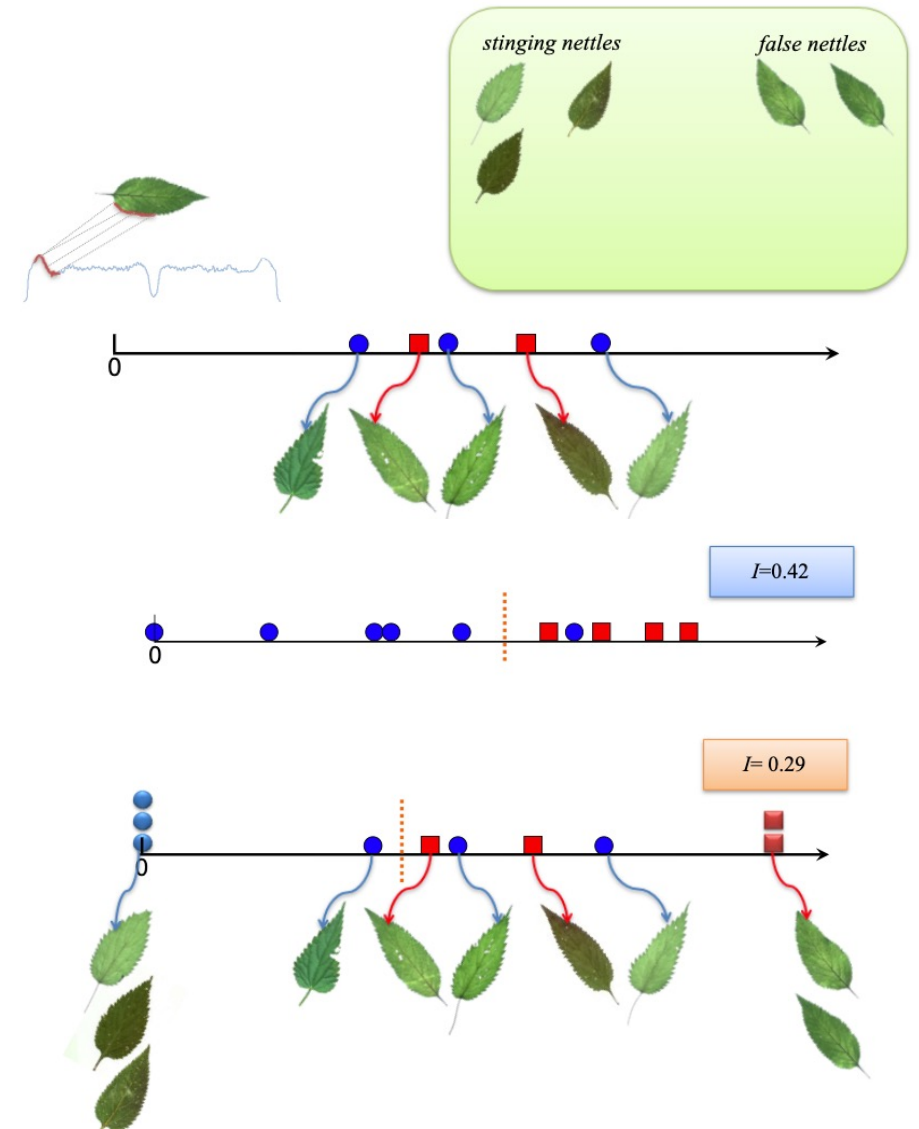
# Distance Early Abandon

- We only need the minimum distance.
- Method
  - Keep the best-so-far distance
  - Abandon the calculation if the partial current distance is larger than best-so-far.
  - We can avoid to compute the full distance for  $S$  if the partial one is greater than the best so far



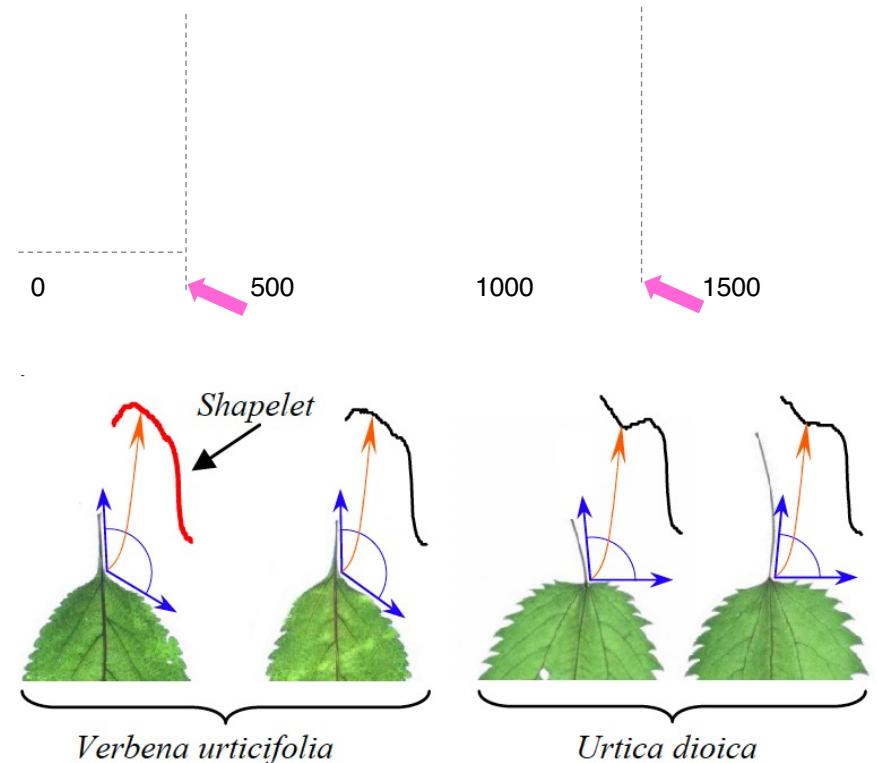
# Admissible Entropy Pruning

- We only need the best shapelet for each class
- For a candidate shapelet
  - We do not need to calculate the distance for each training sample
  - After calculating some training samples, the **upper bound** of information gain (corresponding to the optimistic scenario)  $<$  best candidate shapelet
  - Stop calculation
  - Try next candidate



# Motif/Shapelet Summary

- A **motif** is a repeated pattern/subsequence in a given TS.
- A **shapelet** is a pattern/subsequence which is maximally representative of a class with respect to a given dataset of TSs.



# References

- Matrix Profile I: All Pairs Similarity Joins for Time Series: A Unifying View that Includes Motifs, Discords and Shapelets. Chin-Chia Michael Yeh et al. 1997
- Time Series Shapelets: A New Primitive for Data Mining. Lexiang Ye and Eamonn Keogh. 2016.
- Josif Grabocka, Nicolas Schilling, Martin Wistuba, Lars Schmidt-Thieme (2014): Learning Time-Series Shapelets, in Proceedings of the 20th ACM SIGKDD Conference on Knowledge Discovery and Data Mining, KDD 2014

## Matrix Profile I: All Pairs Similarity Joins for Time Series: A Unifying View that Includes Motifs, Discords and Shapelets

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**Abstract**—The all-pairs-similarity-search (or similarity join) problem has been extensively studied for text and a handful of other datasets. However, surprisingly little progress has been made on similarity joins for time series subsequences. The lack of progress probably stems from the daunting nature of the problem. For even modest sized datasets the obvious nested-loop algorithm can take months, and the typical speed-up techniques in this domain (i.e., indexing, lower-bounding, triangularity pruning and early abandonment) at best produce one or two orders of magnitude speeding. In this work we introduce a novel scalable algorithm for time series subsequence all-pairs-similarity-search. For exceptionally large datasets, the algorithm can be trivially cast as an anytime algorithm and produce high-quality approximate solutions in reasonable time. The exact similarity join algorithm computes the answer to the time series motif and time series discord problem as a side-effect, and our algorithm incidentally provides the fastest known algorithm for both these extensively-studied problems. We demonstrate the utility of our ideas for many time series data mining problems, including motif discovery, anomaly discovery, shapelet discovery, semantic segmentation, density estimation, and contract set mining.

**Keywords**—Time Series, Similarity Joins, Motif Discovery

### 1. INTRODUCTION

The all-pairs-similarity-search (also known as similarity join) problem comes in several variants. The basic task is: Given a collection of data objects, retrieve the nearest neighbor for each object. In the text domain the algorithm has applications in a host of problems, including community discovery, duplicate detection, collaborative filtering, clustering, and query refinement [1]. While virtually all text processing algorithms have analogues in time series data mining, there has been surprisingly little progress on Time Series subsequence All-Pairs-Similarity-Search (TSAPSS).

We believe that this lack of progress stems not from a lack of interest in this useful primitive, but from the daunting nature of the problem. Consider the following example that reflects the needs of an industrial collaborator. A boiler at a chemical refinery reports pressure once a minute. After a year, we have a time series of length 525,600. A plant manager may wish to do a similarity self-join on this data with week-long subsequences (10,080) to discover operating regimes (summer vs. winter or light loading vs. heavy loading, etc). The obvious nested loop algorithm requires 132,800,000,000 Euclidean distance computations. If we assume each one takes 0.0001 seconds, then the join will take 13.28 days. The core contribution of this work is to show that we can reduce this time to 6.3 hours, using an off-the-shelf desktop computer. Moreover, we show that this join can be computed and/or updated incrementally. This we could maintain this join essentially forever on a standard

desktop, even if the data arrival frequency was much faster than one a minute.

Our algorithm uses an ultra-fast similarity search algorithm under  $\epsilon$ -normalized Euclidean distance as a subroutine, exploiting the overlap between subsequences using the classic Fast Fourier Transform (FFT) algorithm.

Our method has the following advantages/features:

- It is exact, providing no false positives or false dismissals.
- It is simple and parameter-free. In contrast, the more general metric space APSS algorithms require building and managing spatial access methods and/or hash functions.
- Our algorithm requires an inconsequential space overhead, just  $O(n)$  with a small constant factor.
- While our exact algorithm is extremely scalable, for extremely large datasets we can compute the results in an anytime fashion, allowing ultra-fast approximate solutions.
- Having computed the similarity join for a dataset, we can incrementally update it very efficiently. In many domains this means we can effectively maintain exact joins on streaming data forever.
- Our method provides full joins, eliminating the need to specify a similarity threshold, which as we will show, is a near impossible task in this domain.
- Our algorithm is embarrassingly parallelizable, both on multicore processors and in distributed systems.

## Time Series Shapelets: A New Primitive for Data Mining

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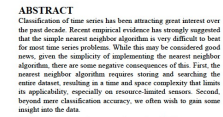


Figure 1 shows some examples of leaves from two classes, *Urcea dioica* (stinging nettle) and *Farberia arctifolia*. These two plants are commonly confused, hence the colloquial name "white nettle" for *Farberia arctifolia*.



Figure 2. A shape can be converted into a one-dimensional "time series" representation. The reason for the highlighted section of the time series will be made apparent shortly.

Suppose we wish to build a classifier to distinguish these two plants, what features should we use? Note the inter-variability of color and size within each class completely defeats the inter-variability between classes, our best hope is based on the shapes of the leaves. However, as we can see in Figure 1, the differences in the global shape are very subtle. Furthermore, it is very common for leaves to have distortions or "irregularities" due to insect damage, and these are likely to confuse any global measures of shape. Instead we attempt the following. We first convert each leaf into a one-dimensional representation as shown in Figure 2.

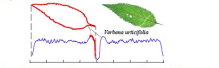


Figure 3. A shape can be converted into a one-dimensional "time series" representation. The reason for the highlighted section of the time series will be made apparent shortly.

Such representations have been successfully used for the classification, clustering and outlier detection of shapes in recent years [8]. However, here we find that using a nearest neighbor classifier with either the (position invariant) Euclidean distance or Dynamic Time Warping (DTW) distance does not significantly outperform random guessing. The reason for this poor performance of these otherwise very competitive classifiers seems to be due to the fact that the data is somewhat noisy (i.e. insect holes, and different stem lengths), and this noise is enough to swamp the subtle differences in the shapes.

In this work we present a novel time series data mining primitive called time series shapelets. Informally, shapelets are time series subsequences which are in some sense maximally representative of a class. While no-better-shapelets can have many uses in data mining, one obvious implication of them is to mitigate the two weaknesses of the nearest neighbor algorithm noted above.

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